

ABSTRACT

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This thesis investigates the existence of a Euclidean algorithm in cubic fields with complex conjugates. This investigation uses the following methods.

The **EUCLID'S ALGORITHM IN CUBIC FIELDS** result of Cassels which states **WITH COMPLEX CONJUGATES** that less than  $-4(12.364...)^2$  possesses a Euclidean Algorithm. By using the modification it is possible to show that some fields of discriminant greater than the above bound, but close to it, also do not possess **ELIZABETH MARY TAYLOR** them.

A second method is to choose an algebraic integer  $\mu$  which is a divisor **ROYAL HOLLOWAY COLLEGE, LONDON** fundamental unit of the field in question and  $n$  is a rational integer. We then determine whether there are any residue classes modulo  $\mu$  which do not contain  $n$  elements of norm of absolute value less than the absolute value of **SUPERVISOR: PROFESSOR H.J. GODWIN**

The next method is an adaptation of a method of Barnes and Swinnerton-Dyer for the real quadratic fields, modified here for the fields in question. This method aims to isolate the points with minimum at least 1.

An indirect method, which is used as the final step of the last method described, is to determine the minimum of numbers of the form  $\frac{x}{1-t}$ , where  $x$  is an integer of the field in question and  $n$  is a positive rational integer.

In addition to existing results, 37 fields have been shown

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ABSTRACT

This thesis investigates the existence of a Euclidean Algorithm in cubic fields with complex conjugates. This investigation is made using the following methods.

The first method is a modification of a result of Cassels which states that no field of discriminant less than  $-\left(\frac{81(5+3\sqrt{3})}{2}\right)^2 = -(412.944 \dots)^2$  possesses a Euclidean Algorithm. By using the modification it is possible to show that some fields of discriminant greater than the above bound, but close to it, also do not possess a Euclidean Algorithm.

A second method is to choose an algebraic integer  $\beta$  which is a divisor of  $1 \pm \epsilon^n$ , where  $\epsilon$  is the fundamental unit of the field in question and  $n$  is a rational integer. We then determine whether there are any residue classes modulo  $\beta$  which do not contain an integer of norm of absolute value less than the absolute value of the norm of  $\beta$ .

The next method is an adaptation of a method of Barnes and Swinnerton-Dyer for the real quadratic fields, modified here for the fields in question. This method aims to isolate the points with minimum at least 1.

An indirect method, which is used as the final step of the last method described, is to determine the minimum of numbers of the form  $\frac{\alpha}{1-\epsilon^n}$ , where  $\alpha$  is an integer of the field in question and  $n$  is a positive rational integer.

In addition to existing results, 37 fields have been shown

to possess a Euclidean Algorithm and it has been established that there is no Euclidean Algorithm in 289 fields. For some fields the inhomogeneous minimum has been determined. Professor R.J. Godwin, The numerical results obtained are given in the last chapter of this work. The listings of the computer programs used for the above methods are in the appendix to this thesis. assistance with some computing details and the adaptations of his routines for use in the methods of this thesis; and who, together with Professor Godwin, provided details of the fields of discriminants -160087 and -163871.

I am grateful to members of the department of Statistics and Computer Science at Royal Holloway College for assistance with the compilation and running of the programs; and as these link to the University of London's CDC 6600 computer the programs were run.

Finally, I wish to acknowledge the financial support of the Science Research Council.



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My thanks are also due to Dr. I.O. Angell for his assistance with some computing details and the adaptations of his routines for use in the methods of this thesis; and who, together with Professor Godwin, provided details of the fields of discriminants -160087 and -169271.

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$\beta \neq 0$  and  
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$$|N(\gamma - \gamma\beta)| \geq |N(\beta)| \tag{1.1}$$

for every algebraic integer  $\gamma$  in  $K$ , where  $N$  represents the norm of the algebraic number, there is said to be no Euclidean algorithm in the field.

Algorithm in the field.

1.2 Let  $L_1, \dots, L_n$  be  $n$  linear forms in the real variables  $x_1, \dots, x_n$  of determinant  $\Delta$ . Let  $(x_1, \dots, x_n)$  be the co-ordinates of the point  $P$  in  $n$ -dimensional space, then, if  $y_1, \dots, y_n$  are real numbers and  $(y_1, \dots, y_n)$  are the co-ordinates of the point  $P_0$ , we say

$$P \equiv P_0 \pmod{1} \text{ if } x_i \equiv y_i \pmod{1}, \dots, x_n \equiv y_n \pmod{1}.$$

We define

$$M(L_1, \dots, L_n; P_0) = \min_{P \equiv P_0} |L_1 \dots L_n|$$

$$\text{and } M(L_1, \dots, L_n) = \max_{P_0} M(L_1, \dots, L_n; P_0).$$

$M(L_1, \dots, L_n)$  is then said to be the inhomogeneous minimum of the linear forms  $L_1, \dots, L_n$ .

1.3 If  $\tilde{\alpha} = (\alpha_1, \dots, \alpha_n)$  where  $\alpha_1, \dots, \alpha_n$  are rational and  $\alpha = \alpha^{(1)}$  is the value of  $L_1$  for  $(x_1, \dots, x_n) = (\alpha_1, \dots, \alpha_n)$ ;  $\alpha$  is an algebraic number and  $\alpha^{(2)}, \dots, \alpha^{(k)}$ , the corresponding values of  $L_2, \dots, L_n$ , are its algebraic conjugates; we also have

INTRODUCTION

and the set of values of  $L_1$  over all rational values of  $x_1, \dots, x_n$ .

I.1 If there are algebraic integers  $\beta, \gamma$  in  $K$  such that  $\beta \neq 0$  and

$$|N(\gamma - \delta\beta)| \geq |N(\beta)| \quad (I.1)$$

for every algebraic integer  $\delta$  in  $K$ , where  $N$  represents the norm of the algebraic number, there is said to be no Euclidean Algorithm in the field.

I.2 Let  $L_1, \dots, L_n$  be  $n$  linear forms in the real variables  $x_1, \dots, x_n$  of determinant  $\Delta$ . Let  $(x_1, \dots, x_n)$  be the co-ordinates of the point  $P$  in  $n$ -dimensional space, then, if  $y_1, \dots, y_n$  are real numbers and  $(y_1, \dots, y_n)$  are the co-ordinates of the point  $P_0$ , we say

$$P \equiv P_0 \pmod{1} \text{ if } x_1 \equiv y_1, \dots, x_n \equiv y_n \pmod{1}.$$

We define

$$M(L_1, \dots, L_n; P_0) = \min_{P \equiv P_0} |L_1 \dots L_n|$$

and  $M(L_1, \dots, L_n) = \max_{P_0} M(L_1, \dots, L_n; P_0)$ .

$M(L_1, \dots, L_n)$  is then said to be the inhomogeneous minimum of the linear forms  $L_1, \dots, L_n$ .

I.3 If  $\tilde{\alpha} = (\alpha_1, \dots, \alpha_n)$  where  $\alpha_1, \dots, \alpha_n$  are rational and  $\alpha = \alpha^{(0)}$  is the value of  $L_1$  for  $(x_1, \dots, x_n) = (\alpha_1, \dots, \alpha_n)$ ,  $\alpha$  is an algebraic number and  $\alpha^{(1)}, \dots, \alpha^{(n)}$ , the corresponding values of  $L_2, \dots, L_n$ , are its algebraic conjugates; we also have

for every algebraic integer  $\alpha$  in  $K$ , since every algebraic integer in  $K$  can be expressed as the quotient of two algebraic integers in  $K$ . The set of values of  $L_1$  over all rational values of  $(x_1, \dots, x_n)$  constitute an algebraic number field  $K$  of discriminant  $D$  where  $D = \Delta^2$ .

We now define

$$M(K, \alpha) = \min_{\gamma} |N(\alpha - \gamma)|$$

where the minimum is over all algebraic integers  $\gamma$  in  $K$ , and

$$M(K) = \max_{\alpha} M(K, \alpha) \tag{I.2}$$

$$= \max_{\alpha} \min_{\gamma} |N(\alpha - \gamma)|$$

where the maximum is over all algebraic numbers  $\alpha$  in  $K$ . We then say  $M(K)$  is the inhomogeneous minimum of the field  $K$ .

We have

$$M(K, \alpha) = M(L_1, \dots, L_n; \tilde{\alpha})$$

and 
$$M(K) \leq M(L_1, \dots, L_n)$$

with equality in the latter relationship when  $M(L_1, \dots, L_n)$  is attained at a rational point.

I.4 Since the quotient of any two numbers of  $K$  is itself a number in  $K$ , for all algebraic integers  $\zeta, \beta$  in  $K$  we have

$$|N(\frac{\zeta}{\beta} - \gamma)| \leq M(K)$$

for some algebraic integer  $\gamma$  in  $K$ ; so that

$$|N(\zeta - \beta\gamma)| \leq M(K) |N(\beta)|.$$

Hence, if  $M(K) < 1$ , the field has a Euclidean Algorithm.

Conversely, if  $M(K) \geq 1$  there exists an algebraic number  $\alpha = \frac{\zeta}{\beta}$  in  $K$ , where  $\zeta, \beta$  are algebraic integers of  $K$ , so that

$$|N(\frac{\zeta}{\beta} - \gamma)| \geq 1$$

for every algebraic integer  $\gamma$  in  $K$ , since every algebraic number in  $K$  may be expressed as the quotient of two algebraic integers in  $K$ ; thus

$$|N(\xi - \beta\gamma)| \geq |N(\beta)|$$

for every algebraic integer  $\gamma$  in  $K$ . Hence, if  $M(K) \geq 1$ , the field has no Euclidean Algorithm.

Thus

$$K \text{ has a Euclidean Algorithm if and only if } M(K) < 1 \quad (I.3)$$

We see that a sufficient condition for the existence of a Euclidean Algorithm is

$$M(L_1, \dots, L_n) < 1.$$

If  $M(K, \alpha) \geq 1$  for some algebraic number  $\alpha$  in  $K$ , immediately  $M(K) \geq 1$ ; thus the non-existence of a Euclidean Algorithm in an algebraic number field may be established by finding a single algebraic number  $\alpha$  in  $K$  for which  $M(K, \alpha) \geq 1$ .

I.5 The question of a Euclidean Algorithm in complex quadratic fields was relatively easily answered, a proof is given in (22) of the fact that the complex quadratic field  $K(\sqrt{m})$  has a Euclidean Algorithm only when  $m = -1, -2, -3, -7, -11$ .

For the real quadratic fields, the fact that the field  $K(\sqrt{d})$  of discriminant  $d > 0$  does not have a Euclidean Algorithm if  $d$  is sufficiently large was first established by a combination of several results, notably those of Berg (3), Behrbohm and Rédei (4), Erdős and Ko (15) and Heilbronn (24). Davenport (8) gave an independent proof of this result; the underlying principles of



the method employed are given in (9). Davenport's method is based on a consideration of the corresponding binary quadratic forms, the more general result given is that: if  $f(x,y) = ax^2 + bxy + cy^2$  is an indefinite binary quadratic form with real coefficients and discriminant  $d = b^2 - 4ac > 0$ , and  $f(x,y)$  does not represent 0 for any integral values of  $x,y$  other than 0,0, there exist real numbers  $\xi, \eta$  with the property that

$$|f(x + \xi, y + \eta)| > K^2 \sqrt{d}.$$

The proof is based on the construction of the point  $(\xi, \eta)$  from an infinite chain of reduced forms all equivalent to the original form. In (8) the reduction due to Hurwitz is used and leads to a specific constant,  $2^{-7}$ , in place of  $K^2$ . Davenport then goes on to show that, when  $a, b, c$  are rational integers,  $\xi$  and  $\eta$ , when constructed as above, are rational, thus there is no Euclidean Algorithm if  $\sqrt{d} > 2^7$ . Since there are only a finite number of real quadratic fields with bounded discriminant there are only a finite number of real quadratic fields with a Euclidean Algorithm.

Davenport used this basic method to show that there is a Euclidean Algorithm in only a finite number of cubic fields with complex conjugate fields (11), and in only a finite number of complex quartic fields whose conjugate fields are also complex (12).

I.6 Cassels (5) establishes the same general results as Davenport but with the following specific constants

For the real quadratic fields  $M(K) > \frac{\Delta}{45.2}$

For the cubic fields with complex conjugate fields  $M(K) > \frac{|\Delta|}{429}$

For the complex quartic fields whose conjugate fields are also complex  $M(K) > \frac{|A|}{5300}$  and to show that for  $(x_1, x_2)$  in

Cassels in fact establishes these results for the minima of the corresponding linear forms and uses the same basic method in each case. A further consideration and extension of the case of the cubic fields with complex conjugate fields is made in chapter 2 of this thesis.

1.7 Chatland (7) gives a summary of the results concerning the real quadratic fields and investigates those fields of discriminant less than  $2^{14}$ , Davenport's bound. He states that there is no Euclidean Algorithm in  $K(\sqrt{m})$  unless

$m = 2, 3, 5, 6, 7, 11, 13, 17, 19, 21, 29, 33, 37, 41, 57, 73, 97$  except possibly when  $m = 193, 241, 313, 337, 457, 601$ . For these last six cases Chatland and Davenport (6) show that there is no Euclidean Algorithm by a modification of Davenport's general method of (8). Inkeri (25) used a method based on that of Erdős and Ko (15) to give an independent proof of these results.

The statement that  $K(\sqrt{97})$  has a Euclidean Algorithm was shown to be false by Barnes and Swinnerton-Dyer (2). In (2) the inhomogeneous minima of the norm forms  $f_m$  of the fields  $K(\sqrt{m})$  for  $m \leq 101$ , except for  $m = 46, 57, 67, 71, 73, 86, 94$ , are obtained; for these last seven forms an upper bound for  $M(f_m)$  is given. Godwin (20) modified the method of Barnes and Swinnerton-Dyer to give  $M(f_m)$  for these seven forms which are the norm forms of those fields with a large fundamental unit.



The basis of the method of (2) is to choose a value  $M'$  slightly less than  $M(f_m)$  and to show that for  $(x_1, y_1)$  in all but a small sub-region  $R$  of a fundamental region modulo 1 in the  $(x, y)$ -plane there is an integer point  $(x_0, y_0)$  such that

$$|f_m(x_1 - x_0, y_1 - y_0)| \leq M'.$$

Having found such a region  $R$ , any point  $P^* = (x^*, y^*)$  at which  $M(f_m, P^*) > M'$  must lie in  $R$  and also in its transform by the fundamental automorph  $E$  of  $f_m(x, y)$ ; hence points at which  $M(f_m)$  is taken must lie in the set  $R_0$  which is common to  $E^n(R)$ , reduced modulo 1, for all integral values of  $n$ . This method is extended to the case of the cubic fields with complex conjugate fields in chapter 4 of this thesis.

I.8 Ennola (14) used a modification of Davenport's method of (8) to obtain the result that for the real quadratic fields

$M(K) > \frac{\Delta}{16 + 6\sqrt{6}} > \frac{\Delta}{30.69}$ . In this paper Ennola supplies a proof of the results concerning the existence of the Euclidean Algorithm in real quadratic fields which are stated above.

I.9 The results outlined so far are concerned with those fields which possess only one fundamental unit; there are no corresponding general results for other fields. Heilbronn (23) proves that Euclid's Algorithm holds in only a finite number of cyclic cubic fields. Godwin (16) obtains an upper bound for the inhomogeneous minimum of some totally real cubic norm forms, which gives rise to the fact that there is a Euclidean Algorithm in those fields of discriminant 49, 81, 148, 169, 229, 257, 316, and 361; an

upper bound is also supplied for the inhomogeneous minimum of the field of discriminant 321 which was shown to possess a Euclidean Algorithm by Smith, the result being quoted in (16). Samet (26) extends the method of Barnes and Swinnerton-Dyer (2) to a special set of linear forms, those corresponding to the field of discriminant 148, and shows that the inhomogeneous minimum is  $\frac{1}{2}$  so that the field has a Euclidean Algorithm. In (27) Samet considers the family of fields defined by

$$\theta^3 - a\theta^2 - 2\theta - a = 0 \quad \text{I.4}$$

where  $a$  is a large positive integer;  $\theta + 1, \theta - 1$  are units with norms  $-1$  and  $+1$  respectively and the discriminant is  $\Delta^1 = 4a^4 + 13a^2 + 32$ . He assumes that  $1, \theta, \theta^2$  is a basis for the field, which is the case if  $\Delta^1$  has no squared factor other than 4 and  $a \equiv 1, 2, 3 \pmod{4}$ . Then, if  $\xi, \eta, \zeta$  are the forms

$$\xi = x + y\theta + z\theta^2$$

$$\eta = x + y\phi + z\phi^2$$

$$\zeta = x + y\psi + z\psi^2$$

where  $\theta, \phi, \psi$  are the roots of I.4, by using a modification of the method of Barnes and Swinnerton-Dyer, Samet shows that for all sufficiently large  $a$

$$(i) \quad M = \frac{1}{8}a^2 \quad \text{if } a \equiv 2 \pmod{4}$$

and this value is taken only at points  $P \equiv (0, 0, \frac{1}{2})$

$$(ii) \quad M = \frac{1}{8}a^2 \left(1 - \frac{1}{a}\right)^3 \quad \text{if } a \equiv 1 \text{ or } 3 \pmod{4}$$

and this value is taken only at points  $P \equiv \left(\frac{1}{a}, 0, \frac{1}{2} - \frac{1}{2a}\right)$ .

Smith (28) uses Heilbronn's method to show that the only cyclic

cubic fields of discriminant  $< 10^8$  which may possess a Euclidean Algorithm are those of discriminants  $7^2, 9^2, 13^2, 19^2, 31^2, 37^2, 43^2, 61^2, 67^2, 73^2, 103^2, 109^2, 127^2$  and  $157^2$ . He then goes on to use the method of Barnes and Swinnerton-Dyer (2), as extended by Samet (26), to give the values of the inhomogeneous minima of the fields of discriminants  $13^2, 19^2, 31^2, 37^2, 43^2$ , and  $73^2$  and to show that the fields of discriminants  $61^2$  and  $67^2$  possess (1) a Euclidean Algorithm. Davenport (10) had previously obtained the values of the inhomogeneous minima of the fields of discriminants  $7^2$  and  $9^2$ , and in this way shown that they possess a Euclidean Algorithm, also, as stated above, Godwin (16) had shown that the fields of discriminants  $13^2 (=169)$  and  $19^2 (=361)$  are Euclidean. Smith shows that the inhomogeneous minima of the fields of discriminants  $13^2, 19^2, 31^2, 37^2$  and  $43^2$  are less than 1, and so the fields possess a Euclidean Algorithm; but that of the field of discriminant  $73^2$  is  $\frac{9}{8}$ , greater than 1, so that this field does not possess a Euclidean Algorithm.

For fields of higher degree, Godwin (17) has shown that the totally real quartic fields with discriminants 725, 1125, 1600, 1957, 2225, 2304, 2624, 2777 and 4205 possess a Euclidean Algorithm, as does the totally real quintic field  $K(2\cos(2\pi/11))$ , which is the totally real quintic field of least discriminant.

I.10 I now turn to the cubic fields with complex conjugate fields, which are the subject of this thesis. Godwin (18) adapted his method of (16) and (17) to show that the fields of discriminants

-23, -31, -44, -59, -76, -83, -87, -104, -107, -108, -116, -135, -139, -140, -152 possess a Euclidean Algorithm. The basis of this method is to show that, if a polynomial  $P(x)$  has zeros which are not too far apart, the set of  $x$  for which  $P(x) < 1$  contains a complete set of residues modulo 1. The fifteen fields shown to possess a Euclidean Algorithm are the first fifteen in the table of cubic fields with complex conjugate fields of Angell (1); the method yields no result for other fields, but by the methods described in this thesis I have shown that the next two fields in the table, those of discriminants -172 and -175, also possess a Euclidean Algorithm. The next, that of discriminant -199, has inhomogeneous minimum 1; thus, it is the first in the table which does not possess a Euclidean Algorithm.

The method of Barnes and Swinnerton-Dyer may be adapted to show that a field possesses I.11 In the following  $(1, \theta, \lambda)$  is an integral basis for the cubic number field  $K = K(\theta)$ , where  $\theta$  is the real zero of the polynomial  $x^3 - ax^2 + bx - c$ . The field has discriminant  $D$  where  $D < 0$ , and, if  $l$  is the index of the polynomial over the field,

$$\lambda = \frac{\theta^2 + t\theta + s}{l} \quad 0 \leq t, s < l$$

where  $t$  and  $s$  are rational integers.

$\epsilon$  is the fundamental unit of  $K$  satisfying  $0 < \epsilon < 1$  and has algebraic conjugates  $\epsilon', \epsilon'' = \bar{\epsilon}'$ .

Let  $\theta$  have algebraic conjugates  $\phi, \bar{\phi}$ , and  $\lambda$  have algebraic conjugates  $\psi, \bar{\psi}$  then  $\psi = \frac{\phi^2 + t\phi + s}{l}$ .

Throughout the thesis, the terms number and integer without



qualification mean cubic number and cubic integer. Only fields of class number 1 will be considered, since this is a necessary condition for the existence of a Euclidean Algorithm.

Part I describes the theory of the methods of investigation. Chapter 1 gives a summary of the properties of relative minima, since these will be required in later chapters. Chapter 2 gives an extension of a result of Cassels. Chapter 3 considers the fact: if there is no integer of  $K$  of norm of absolute value less than  $|N(\beta)|$ , congruent to  $\gamma$  modulo  $\beta$  for two integers  $\beta$  and  $\gamma$  in  $K$ , where  $\beta$  is non-zero, there is no Euclidean Algorithm in  $K$ ; it is shown how this fact may be formulated into a practicable method to show that a field does not possess a Euclidean Algorithm. Chapter 4 shows how the method of Barnes and Swinnerton-Dyer may be adapted to show that a field possesses a Euclidean Algorithm. Chapter 5 shows how the minimum of a point of the form  $\frac{\alpha}{1-\epsilon^n}$ , where  $\alpha$  is an integer of  $K$ , may be computed. Chapter 6 uses one particular field to demonstrate each of the methods employed, and describes the limitations of each of these methods.

Part II gives descriptions of the actual programs used, and contains a table showing the results obtained.

## CHAPTER 1

### RELATIVE MINIMA AND IDEALS WHOSE NORM IS A RATIONAL PRIME.

1.1 We give a short summary of certain properties of relative minima, in particular in the context of fields  $K$  which consist of elements  $x + y\theta + z\theta^2$  where  $x, y, z$  are rational numbers

#### PART I

and  $\theta$  is the real root of a cubic equation of negative discriminant; relative minima will be used in some later chapters.

Consider the space  $B$  with general point  $(x, y + iz)$ ,  
**THE THEORY OF THE METHODS OF INVESTIGATION.**

where  $x, y, z$  are elements of the real number field. We define

the space  $A$  to be the set of points  $\hat{\alpha} = (\delta(\theta), \delta(\phi))$  where  $\alpha = \delta(\theta) = d_1\theta^2 + d_2\theta + d_3$  is in  $\mathcal{U}$ , and  $d_1, d_2, d_3$  are rational real numbers; the isomorphism  $\alpha \leftrightarrow \hat{\alpha}$  then establishes an isomorphism  $K \leftrightarrow A$ . We have  $B \supset A$  and addition, subtraction, multiplication and division are defined component-wise in  $B$  and in  $A$ .

1.2 If  $\hat{\alpha} = (x, y + iz)$  is in  $B$ , the directional parameters  $\rho_1, \rho_2$  are defined by

$$\rho_1(\hat{\alpha}) = |x|$$

$$\rho_2(\hat{\alpha}) = y^2 + z^2$$

From their definitions  $\rho_1$  and  $\rho_2$  are multiplicative.

The normed body of a point  $\hat{\alpha}$  in  $A$  is defined to be the region  $V \subset B$  given by

$$V = \{ \hat{\omega} : \hat{\omega} \in B, \rho_j(\hat{\omega}) < \rho_j(\hat{\alpha}) ; j = 1, 2 \}$$

1.3 If  $\mathfrak{f}$  is an ideal of  $K$  and has typical element  $\alpha$ ,  $\hat{\mathfrak{f}}$  is

the ideal lattice is CHAPTER 1 Multiplicative  
 lattices  $\mathfrak{f}$  in  $K$  are those lattices for which the product of any  
 RELATIVE MINIMA AND IDEALS WHOSE NORM IS A RATIONAL PRIME.  
 two points of  $\mathfrak{f}$  also belongs to  $\mathfrak{f}$ . The ideal lattices are

1.1 We give a short summary of certain properties of relative  
 minima, in particular in the context of fields  $K$  which consist  
 of elements  $x + y\theta + z\theta^2$  where  $x, y, z$  are rational numbers  
 and  $\theta$  is the real root of a cubic equation of negative discriminant;  
 relative minima will be used in some later chapters.

Consider the space  $B$  with general point  $\hat{\omega} = (x, y + iz)$ ,  
 where  $x, y, z$  are elements of the real number field. We define  
 the space  $A$  to be the set of points  $\hat{\alpha} = (\delta(\theta), \delta(\phi))$  where  
 $\alpha = \delta(\theta) = d_1\theta^2 + d_2\theta + d_3$  is in  $K$ , and  $d_1, d_2, d_3$  are rational  
 real numbers; the isomorphism  $\alpha \leftrightarrow \hat{\alpha}$  then establishes an  
 isomorphism  $K \leftrightarrow A$ . We have  $B \supset A$  and addition, subtraction,  
 multiplication and division are defined component-wise in  $B$  and  
 in  $A$ .

1.2 If  $\hat{\omega} = (x, y + iz)$  is in  $B$ , the directional parameters  
 $\rho_1, \rho_2$  are defined by

$$\rho_1(\hat{\omega}) = |x|$$

$$\rho_2(\hat{\omega}) = y^2 + z^2.$$

From their definitions  $\rho_1$  and  $\rho_2$  are multiplicative.

The normed body of a point  $\hat{\alpha}$  in  $A$  is defined to be the  
 region  $V \subset B$  given by

$$V = \{ \hat{\omega} : \hat{\omega} \in B, \rho_j(\hat{\omega}) < \rho_j(\hat{\alpha}); j = 1, 2 \}.$$

1.3 If  $\mathfrak{f}$  is an ideal of  $K$  and has typical element  $\alpha$ ,  $\hat{\mathfrak{f}}$  is

the ideal lattice in  $A$  with typical point  $\hat{\alpha}$ . Multiplicative lattices  $\hat{\mathfrak{f}}$  in  $A$  are those lattices for which the product of any two points of  $\hat{\mathfrak{f}}$  also belongs to  $\hat{\mathfrak{f}}$ . The ideal lattices are multiplicative.

If  $\hat{\alpha}$  is in  $\hat{\mathfrak{f}}$ ,  $\hat{\mathfrak{f}}' = \{\hat{\alpha}\hat{\delta} : \hat{\delta} \in \hat{\mathfrak{f}}\}$  is also in  $\hat{\mathfrak{f}}$  and is a multiplicative lattice; point-lattice multiplication is defined by

$$\hat{\mathfrak{f}}' = \hat{\alpha}\hat{\mathfrak{f}}.$$

If there exists a multiplicative lattice  $\hat{\mathfrak{f}}''$  such that  $\hat{\mathfrak{f}} = \hat{\alpha}\hat{\mathfrak{f}}''$ , point-lattice division is defined by

$$\hat{\mathfrak{f}}'' = \frac{1}{\hat{\alpha}}\hat{\mathfrak{f}}.$$

1.4 A point of an ideal lattice  $\hat{\mathfrak{f}}$  in  $A$  is called a relative minimum of  $\hat{\mathfrak{f}}$  if its normed body contains no other point of  $\hat{\mathfrak{f}}$  except the origin.

For a relative minimum  $\hat{\Omega}$  of  $\hat{\mathfrak{f}}$  we define the region  $V(\hat{\Omega}, d)$  by

$$V(\hat{\Omega}, d) = \left\{ \begin{array}{l} \hat{\alpha} : \rho_{j_1}(\hat{\alpha}) < \rho_{j_1}(\hat{\Omega}) \quad j_1 = 1, 2, \dots \\ \rho_{j_2}(\hat{\alpha}) \leq \rho_{j_2}(\hat{\Omega}) + d \end{array} \right.$$

From Minkowski's convex body theorem there is a lattice point  $\hat{\Omega}_1$  which lies in  $V(\hat{\Omega}, d)$  for some  $d > 0$ ; and if  $\hat{\Omega}_1$  is the first such point obtained by increasing  $d$  from 0,  $\hat{\Omega}_1$  is also a relative minimum of  $\hat{\mathfrak{f}}$ . If the point  $\hat{\delta}$  satisfies the conditions stated above for  $\hat{\Omega}_1$ ,  $-\hat{\delta}$  also satisfies them.  $\hat{\delta}' = (x, y + iz)$  is chosen to be  $\hat{\delta}$  or  $-\hat{\delta}$  so that



We suppose,  $x > 0$  if  $j_1 = 2, j_2 = 1$   
 $-\frac{\pi}{2} < \arg(y + iz) \leq \frac{\pi}{2}$  if  $j_1 = 1, j_2 = 2$ ;

in this way  $\hat{\Omega}_1 = \hat{\delta}'$  is uniquely defined, and will be said to be the relative minimum adjacent to  $\hat{\Omega}$  in the  $\rho_{j_1}$  direction.

By this means a two-way chain of relative minima  $\dots \hat{\Omega}_{-1}, \hat{\Omega}_0, \hat{\Omega}_1, \hat{\Omega}_2, \dots$  may be defined. If  $\hat{\Omega}_1, \hat{\Omega}_2, \dots$  are relative minima of  $\hat{\mathcal{J}}, \Omega_1, \Omega_2, \dots$  are said to be relative minima of  $\mathcal{J}$ .

A more detailed account of relative minima may be found in (13).

1.5 From now on, in this chapter, we restrict the consideration to the case when  $K$  is a cubic field of negative discriminant. In this case, if  $\alpha$  is in  $K$ , so that  $\hat{\alpha}$  is in  $A$ , and  $N(\hat{\alpha})$  is defined by

$$N(\hat{\alpha}) = \rho_1(\hat{\alpha}) \rho_2(\hat{\alpha}),$$

we have  $N(\hat{\alpha}) = |N(\alpha)|$ .

#### LEMMA 1.1

Given an ideal  $\mathcal{J}$ , if  $\Omega_1, \Omega_2, \dots, \Omega_j, \dots$  are successive relative minima of  $\mathcal{J}$ , and  $\mathcal{J}$  is divided by these relative minima to give the ideals

$$\mathcal{J}_1 = \frac{\mathcal{J}}{\Omega_1}; \quad \mathcal{J}_2 = \frac{\mathcal{J}}{\Omega_2}; \quad \dots; \quad \mathcal{J}_j = \frac{\mathcal{J}}{\Omega_j}; \quad \dots,$$

for some  $k$ ,  $\mathcal{J}_k$  is the unit ideal  $(1, \theta, \lambda)$  and  $\Omega_k$  produces  $\mathcal{J}$ .

#### PROOF OF LEMMA 1.1

As stated in the introduction, this thesis considers only cubic fields of negative discriminant with class number 1; thus there exists an integer  $\alpha_1$  for which  $\mathcal{J} = (\alpha_1)$ .

for finding the ideals and corresponding relative minima for

We suppose, first of all, that  $\alpha_1$  is not a relative minimum of  $\mathfrak{f}$ , where  $\mathfrak{f}$  is an ideal of norm  $p$ .

of  $\mathfrak{f}$ , we may find  $\alpha_2$ , a relative minimum of  $\mathfrak{f}$ , such that

$$\rho_{j_1}(\hat{\alpha}_2) < \rho_{j_1}(\hat{\alpha}_1)$$

and

$$\rho_{j_2}(\hat{\alpha}_2) < \rho_{j_2}(\hat{\alpha}_1)$$

which implies

$$|N(\alpha_2)| = N(\hat{\alpha}_2) = \rho_{j_1}(\hat{\alpha}_2)\rho_{j_2}(\hat{\alpha}_2) < \rho_{j_1}(\hat{\alpha}_1)\rho_{j_2}(\hat{\alpha}_1) = N(\hat{\alpha}_1) = |N(\alpha_1)|,$$

this contradicts the choice of  $\alpha_1$  as an integer which produces  $\mathfrak{f}$ .

From Voronoi, as described in (13), there are only a finite

number of distinct ideals among  $\mathfrak{f}_1, \mathfrak{f}_2, \dots, \mathfrak{f}_j, \dots$ ; so

we have, for some rational integer  $f$ ,

$$\mathfrak{f}_i = \mathfrak{f}_{i+f}, \quad \mathfrak{f}_i \neq \mathfrak{f}_{i+j} \quad \text{for } 0 < j \leq f-1$$

so that  $\frac{\mathfrak{f}}{\Omega_i} = \frac{\mathfrak{f}}{\Omega_{i+f}}$

which gives  $\frac{\Omega_{i+f}}{\Omega_i} = \epsilon$ .

Since  $\alpha_1$  is a relative minimum of  $\mathfrak{f}$ , for some rational

integers  $k, k_1$  we have

$$\Omega_k = \epsilon^{k_1} \alpha_1$$

so that  $\Omega_k$  produces  $\mathfrak{f}$ .

Also, if  $\mathfrak{f}_k = \frac{\mathfrak{f}}{\Omega_k}$

$$N(\mathfrak{f}_k) = N\left(\frac{\mathfrak{f}}{\Omega_k}\right) = 1,$$

where the norm of an ideal is the norm of any integer which produces it, thus  $\mathfrak{f}_k$  is the unit ideal.

1.6 We now turn to a description of the method used for computing, for any rational prime  $p$ , the ideals of norm  $p$ , the relative minima of each of these ideals and the integers which produce the ideals. The method described follows that used by Angell (1)

for finding the ideals and corresponding relative minima for of the algorithm we consider the ideal  $\mathfrak{f} = (1, \frac{\theta + i_2}{p}, \frac{\lambda + i_4}{p})$  rational integers which are not necessarily prime.

Any ideal of norm  $p$  may be represented in the form (for some rational integers  $q$  and  $a$ , because

$$\mathfrak{f} = (i_1, i_2\theta + i_3, i_4\lambda + i_5\theta + i_6)$$

where  $i_1, i_2, i_3, i_4, i_5, i_6$  are positive rational integers,  $i_1 \cdot i_2 \cdot i_4 = p$  and from the additive property of ideals  $0 \leq i_2 < i_1$ ,  $0 \leq i_4 < i_1$ ,  $0 \leq i_5 < i_2$ ,  $0 \leq i_6 < i_1$ . Since  $p$  is prime we have  $i_1 = p$ ,  $i_2 = 1$ ,  $i_4 = 1$ ,  $i_5 = 0$ , so the ideal may be represented in the form

$$\mathfrak{f} = (p, \theta + i_3, \lambda + i_6) \quad \text{where } 0 \leq i_3 < p, 0 \leq i_6 < p.$$

Thus, given  $p$ , for each pair of values,  $i_3, i_6$ , in the given range, a test is made for

$$\mathfrak{f} = (p, \theta + i_3, \lambda + i_6)$$

being an ideal. Automatically the additive property of ideals is held by  $\mathfrak{f}$ , now it remains to show that all the linear combinations of  $p, \theta + i_3, \lambda + i_6$ , when multiplied by an integer of the field, are in  $\mathfrak{f}$ . This is done by checking that  $p\theta, p\lambda, \theta^2 + i_3\theta, \theta\lambda + i_3\lambda, \lambda\theta + i_6\theta, \lambda^2 + i_6\lambda$  all belong to  $\mathfrak{f}$ . If all the conditions are satisfied then  $\mathfrak{f}$  is an ideal of norm  $p$ .

We note that, since  $N(p) = p^3$ , there are at most three distinct ideal factors of  $(p)$ , other than the unit ideal, and so at most three distinct ideals of norm  $p$ .

1.7 The method of determination of the relative minima of the ideals follows the steps of Voronoi's algorithm as enumerated in (13). We have the ideal  $\mathfrak{f} = (p, \theta + i_3, \lambda + i_6)$ , for application

of the algorithm we consider the ideal  $\frac{f}{p} = (1, \frac{\theta + i_3}{p}, \frac{\lambda + i_6}{p})$  which under the transformation  $\xi = 3\theta - a$ , so that  $\xi^3 = q\xi + n$  for some rational integers  $q$  and  $n$ , becomes

$$(1, \frac{m_1 + m_2\xi + m_3\xi^2}{i_9}, \frac{n_1 + n_2\xi + n_3\xi^2}{i_9})$$

for some rational integers  $m_1, m_2, m_3, n_1, n_2, n_3, i_9$ . Following the steps of the algorithm we obtain a basis  $(1, \theta_3, \theta_h)$  for  $\frac{f}{p}$ , where  $\theta_3$  is the relative minimum of  $\frac{f}{p}$  adjacent to 1 in the  $\rho_2$  direction; this basis is said to be the reduced basis of the lattice. The ideal  $\frac{f}{p}$  is divided by  $\theta_3$  to give the ideal with basis  $(1, \frac{\theta_h}{\theta_3}, \frac{1}{\theta_3})$  and the process is repeated to obtain a reduced basis for this ideal.

In this way we obtain a sequence of ideals  $\mathfrak{f}_1 = \frac{f}{p}, \mathfrak{f}_2 = \frac{\mathfrak{f}_1}{\theta_3^{(1)}}, \dots, \mathfrak{f}_{k+1} = \frac{\mathfrak{f}_k}{\theta_3^{(k)}}$  with reduced bases  $(1, \theta_3^{(1)}, \theta_h^{(1)}) = (1, \theta_3, \theta_h), (1, \theta_3^{(2)}, \theta_h^{(2)}), \dots, (1, \theta_3^{(k+1)}, \theta_h^{(k+1)})$  respectively, where  $\mathfrak{f}_{k+1}$  is the unit ideal. The reduced basis of the unit ideal is previously found by applying one cycle of Voronoi's algorithm to the ideal  $(1, \theta, \lambda)$ .

From this we see that

$$\mathfrak{f} = (p \prod_{i=1}^k \theta_3^{(i)})$$

and, when  $p = 1$ , the relative minima of the field are given by

$$1, \theta_3^{(1)}, \theta_3^{(1)}\theta_3^{(2)}, \theta_3^{(1)}\theta_3^{(2)}\theta_3^{(3)}, \dots$$

so that  $\mathfrak{f}$  is the integral lattice of the field; in this case

$$\prod_{i=1}^k \theta_3^{(i)} = \epsilon \quad \text{the fundamental unit of the field.}$$

## AN ADAPTATION OF A RESULT OF CASSELS.

In the more general case when  $l$  is not equal to 1 we may

2.1 Let  $K(\theta)$  be the cubic field of discriminant  $D = -\Delta^2$ , establish that there exists a point  $(\xi, \eta, \zeta)$  of  $\mathfrak{a}$  such that if  $\theta$  is a zero of the polynomial  $P(x)$  let  $l$  be the index of  $P$  over the field. Define the space  $\mathfrak{a}$  to consist of points  $(\xi, \eta, \zeta)$  where  $\zeta \in K(\theta)$  and  $\xi \pm i\eta$  are the algebraic conjugates of  $\zeta$ ; let  $(\xi, \eta, \zeta)$  be said to be an integer point of  $\mathfrak{a}$  when  $\zeta$  is an integer of  $K(\theta)$ , then the integer points of  $\mathfrak{a}$  form a lattice  $\mathfrak{K}_0$ .

Cassels (5) shows that there exists a point  $(\xi, \eta, \zeta)$  of  $\mathfrak{a}$  such that

$$\min |\zeta_0 (\xi_0^2 + \eta_0^2)| > \frac{2|\Delta|}{81(5 + 3\sqrt{3})} = \frac{|\Delta|}{412.944\dots} \quad 2.1$$

where the minimum is over all points  $(\xi_0, \eta_0, \zeta_0) \equiv (\xi, \eta, \zeta) \pmod{1}$ .

In the following it will be shown that, for any field for which  $l = 1$ , a value  $M$ , which depends on the relative minima of the field, may be found such that there exists a point  $(\xi, \eta, \zeta)$  of  $\mathfrak{a}$  such that

$$\min |\zeta_0 (\xi_0^2 + \eta_0^2)| > \frac{2|D|}{81(9 + 5\sqrt{3})M} = \frac{|D|}{715.24\dots M} \quad 2.2$$

where the minimum is over all points  $(\xi_0, \eta_0, \zeta_0) \equiv (\xi, \eta, \zeta) \pmod{1}$ .

For any such field

$$\frac{2|D|}{81(9 + 5\sqrt{3})M} \geq \frac{2|\Delta|}{81(5 + 3\sqrt{3})} \quad 2.3$$

we note that

$$81(9 + 5\sqrt{3}) = \sqrt{3} \cdot 81(5 + 3\sqrt{3}).$$

Using this result I have investigated two fields of discriminant such that no result concerning the Euclidean Algorithm can be

$$|\Delta| > 24 \geq |\Delta_{\text{next}}|$$



obtained using Cassels' result; both of these fields were shown to possess no Euclidean Algorithm.

In the more general case when  $l$  is not equal to 1 we may establish that there exists a point  $(\xi, \eta, \zeta)$  of  $\mathcal{A}$  such that

$$\min |\zeta(\xi^2 + \eta^2)| > \frac{2|D|l^2}{81(9 + 5\sqrt{3})M^*}, \quad 2.4$$

where the minimum is over all points  $(\xi_0, \eta_0, \zeta_0) \equiv (\xi, \eta, \zeta) \pmod{1}$ , for a particular value  $M^*$  which depends on the field in question. However, we cannot establish a relationship with Cassels' result similar to 2.3.

2.2 Suppose that  $\mathcal{K}$  is any three dimensional lattice which contains an infinite sequence of points  $P_n = (\beta_n, \gamma_n, \alpha_n)$  such that

$$\begin{aligned} |\alpha_n| &> |\alpha_{n+1}| \\ |\beta_n^2 + \gamma_n^2| &< |\beta_{n+1}^2 + \gamma_{n+1}^2| \end{aligned}$$

LEMMA 2.1

Define  $M = \max_n |\alpha_n(\beta_{n+1}^2 + \gamma_{n+1}^2)|$

then, given  $p, q$  such that  $pq^2 \geq M$ , there is a point of the lattice,  $P = (\beta, \gamma, \alpha)$ , not equal to  $(0,0,0)$ , such that

$$|\beta^2 + \gamma^2| \leq q^2$$

$$|\alpha| \leq p$$

This result is best possible for the value of the lower bound of  $pq^2$  when the sequence  $\{\beta_n, \gamma_n, \alpha_n\}$  is a sequence of relative minima of the lattice and is periodic.

PROOF OF LEMMA 2.1

Given  $p$ , we define  $n$  by

$$|\alpha_n| > p \geq |\alpha_{n+1}|$$

which is possible since the sequence  $\{\alpha_n\}$  is decreasing. If

$$q^2 < |\beta_{n+1}^2 + \gamma_{n+1}^2|$$

then  $pq^2 < |\alpha_n| |\beta_{n+1}^2 + \gamma_{n+1}^2| \leq M$

which is false. Hence

$$q^2 \geq |\beta_{n+1}^2 + \gamma_{n+1}^2|$$

and  $(\beta_{n+1}, \gamma_{n+1}, \alpha_{n+1})$  is the required point.

If the sequence  $\{\beta_n, \gamma_n, \alpha_n\}$  is a sequence of relative minima of the lattice, for any rational integer  $n$  there is no point  $(\beta, \gamma, \alpha)$  for which

$$|\alpha| < |\alpha_n|$$

and  $|\beta^2 + \gamma^2| < |\beta_{n+1}^2 + \gamma_{n+1}^2|$ .

If the sequence is periodic, there exists a rational integer  $m$  for which

$$M = \alpha_m (\beta_{m+1}^2 + \gamma_{m+1}^2).$$

Let  $M' = M - \delta$  for some positive real number  $\delta$  and let  $p, q$  be such that

$$M' \leq pq^2 < M$$

$$p < |\alpha_m|$$

$$q < |\beta_{m+1}^2 + \gamma_{m+1}^2|^{\frac{1}{2}}$$

then any point  $(\beta, \gamma, \alpha)$  satisfying

$$|\beta^2 + \gamma^2| \leq q^2 \quad \text{and} \quad |\alpha| \leq p$$

must also satisfy

$$|\beta^2 + \gamma^2| < |\beta_{m+1}^2 + \gamma_{m+1}^2| \quad \text{and} \quad |\alpha| < |\alpha_m|$$

which contradicts the definition of  $\{\beta_n, \gamma_n, \alpha_n\}$  as a sequence of relative minima. Hence  $M$  is best possible in this case.

2.3 Let the linear forms  $L_1, L_2, L_3$  be defined by,  $t_1, t_2, t_3$

$$\begin{aligned} L_1 &= x + y\theta + z\lambda \\ L_2 &= x + y\phi + z\psi \\ L_3 &= x + y\bar{\phi} + z\bar{\psi} \end{aligned} \tag{2.5}$$

where  $x, y, z$  are rational, so that  $L_2$  and  $L_3 = \bar{L}_2$  are the algebraic conjugates of  $L_1$ .

The matrix corresponding to 2.5 holds for  $\phi, \psi$  and for  $\bar{\phi}, \bar{\psi}$ .

$$\begin{pmatrix} 1 & \theta & \lambda \\ 1 & \phi & \psi \\ 1 & \bar{\phi} & \bar{\psi} \end{pmatrix}$$

has adjoint

$$M_1 = \begin{pmatrix} \phi\bar{\psi} - \bar{\phi}\psi & \bar{\phi}\lambda - \theta\bar{\psi} & \theta\psi - \phi\lambda \\ \psi - \bar{\psi} & \bar{\psi} - \lambda & \lambda - \psi \\ \bar{\phi} - \phi & \theta - \bar{\phi} & \phi - \theta \end{pmatrix}$$

Define the linear forms  $M_1, M_2, M_3$  in the rational variables

$u, v, w$  by  $M_1 = \frac{1}{i} \{ (\phi\bar{\psi} - \bar{\phi}\psi)u + (\psi - \bar{\psi})v + (\bar{\phi} - \phi)w \}$  2.10

$$M_2 = \frac{1}{i} \{ (\bar{\phi}\lambda - \theta\bar{\psi})u + (\bar{\psi} - \lambda)v + (\theta - \bar{\phi})w \}$$

$$M_3 = \frac{1}{i} \{ (\theta\psi - \phi\lambda)u + (\lambda - \psi)v + (\phi - \theta)w \}$$

then  $M_1$  is real,  $M_2, M_3$  are conjugate complex and are the algebraic

conjugates of  $M_1$ ;  $M_1, M_2, M_3$  have determinant  $\Delta^2$ . We also have

$$M_1L_1 + M_2L_2 + M_3L_3 = (xu + yv + zw)\Delta \tag{2.7}$$

These definitions of  $M_1, M_2, M_3, L_1, L_2, L_3$  are then the same as those given for the ternary cubic case in Cassels (5).



2.4 There are rational integers  $a_\lambda, b_\lambda, c_\lambda, l_1, l_2, l_3, t_1, t_2, t_3, p_1, p_2, p_3$  such that  $\lambda, \psi, \bar{\psi}$  are the zeros of

$$x^3 - a_\lambda x^2 + b_\lambda x - c_\lambda = 0, \quad 2.8$$

and  $\lambda^2 = l_1 \lambda + l_2 \theta + l_3$  and  $\theta^2 = t_1 \lambda + t_2 \theta + t_3$  and only if there is a point

$$\theta \lambda = p_1 \lambda + p_2 \theta + p_3;$$

also, equalities corresponding to 2.8 hold for  $\phi, \psi$  and for  $\bar{\phi}, \bar{\psi}$ .

We have  $\phi \bar{\psi} - \bar{\phi} \psi = (\phi - \bar{\phi})(a_\lambda - \lambda) + p_1(\bar{\psi} - \psi) + p_2(\bar{\phi} - \phi)$

and  $(\phi + \bar{\phi})(\phi - \bar{\phi}) = t_1(\psi - \bar{\psi}) + t_2(\phi - \bar{\phi})$ ,

$$\text{thus } \psi - \bar{\psi} = (\bar{\phi} - \phi) \left\{ \frac{\theta - a_\lambda + t_1}{t_1} \right\}$$

and so

$$M_1 = \frac{(\bar{\phi} - \phi)}{i} \left\{ u(\lambda - a_\lambda) + (v - up_1) \frac{(\theta - a_\lambda)}{t_1} + w + up_2 + \frac{(v - up_1)}{t_1} t_2 \right\}$$

with similar expressions for  $M_2$  and  $M_3$ .

Thus the linear forms  $M_1, M_2, M_3$  may be given by

$$\begin{aligned} M_1 &= \frac{(\bar{\phi} - \phi)}{it_1} \{ u\lambda t_1 + v\theta + w \} \\ M_2 &= \frac{(\theta - \bar{\phi})}{it_1} \{ u\psi t_1 + v\phi + w \} \\ M_3 &= \frac{(\phi - \theta)}{it_1} \{ u\bar{\psi} t_1 + v\bar{\phi} + w \} \end{aligned} \quad 2.10$$

and the set of points  $(\frac{M_1 + M_2}{2}, \frac{M_1 - M_2}{2i}, M_1)$ , for integer values of  $u, v, w$ , forms a three dimensional lattice  $\mathcal{K}_T$ . We also note that  $t_1 = l$ , the index of  $P$  over  $K$ .

2.5 The general case when  $l$  may take any value is discussed

in section 2.10.

We now restrict the consideration to the case when  $l = 1$ .

We have  $\lambda = \theta^2, \psi = \phi^2$  and  $\bar{\psi} = \bar{\phi}^2$  thus

$$M_1 = \frac{(\bar{\phi} - \phi)}{i} \{ u\theta^2 + v\theta + w \}$$

$$M_2 = \frac{(\theta - \bar{\phi})}{i} \{ u\phi^2 + v\phi + w \}$$

which implies  $M_3 = \frac{(\phi - \theta)}{i} \{ u\bar{\phi}^2 + v\bar{\phi} + w \}$ ,  $\gamma = 1$  of  $\mathcal{K}_0$  for which

then  $(\beta, \gamma, \alpha)$ , where  $\alpha, \beta \pm i\gamma$  are values of  $M_1, M_2, M_3$

respectively, is a point of  $\mathcal{K}_T$  if and only if there is a point

$(\xi, \eta, \zeta)$  of  $\mathcal{K}_0$  such that

$$\alpha = \frac{(\bar{\phi} - \phi)}{i} \zeta; \quad \beta + i\gamma = \frac{(\theta - \bar{\phi})}{i} (\xi + i\eta); \quad \beta - i\gamma = \frac{(\phi - \theta)}{i} (\xi - i\eta).$$

In the field  $K$  we may find an infinite sequence of relative minima  $\{\zeta_n\}$  which corresponds to an infinite sequence of points

of  $\mathcal{K}_0, \{\xi_n, \eta_n, \zeta_n\}$  for which

$$|\zeta_{n-1}| > |\zeta_n|$$

$$|\xi_{n-1}^2 + \eta_{n-1}^2| < |\xi_n^2 + \eta_n^2|.$$

Thus there is a sequence of points  $\{\beta_n, \gamma_n, \alpha_n\}$  in  $\mathcal{K}_T$ , for which

$$\alpha_n = \frac{(\bar{\phi} - \phi)}{i} \zeta_n$$

so that  $|\alpha_n| = \left| \frac{\bar{\phi} - \phi}{i} \right| |\zeta_n| > \left| \frac{\bar{\phi} - \phi}{i} \right| |\zeta_{n+1}| = |\alpha_{n+1}|;$

that is

$$|\alpha_n| > |\alpha_{n+1}|.$$

Also

$$\beta_n + i\gamma_n = \frac{(\theta - \bar{\phi})}{i} (\xi_n + i\eta_n)$$

$$\beta_n - i\gamma_n = \frac{(\phi - \theta)}{i} (\xi_n - i\eta_n)$$

so that

$$|\beta_n^2 + \gamma_n^2| = |(\theta - \bar{\phi})(\phi - \theta)| |\xi_n^2 + \eta_n^2| < |(\theta - \bar{\phi})(\phi - \theta)| |\xi_{n+1}^2 + \eta_{n+1}^2| = |\beta_{n+1}^2 + \gamma_{n+1}^2|;$$

that is

$$|\beta_n^2 + \gamma_n^2| < |\beta_{n+1}^2 + \gamma_{n+1}^2|.$$

We also have the sequence  $\{\beta_n, \gamma_n, \alpha_n\}$  a sequence of relative minima of the lattice  $\mathcal{K}_T$ ; if not, for some integer  $n$  there is

a point  $(\beta, \gamma, \alpha)$  of  $\mathcal{K}_T$  such that

hence  $|\beta^2 + \gamma^2| \leq |\beta_{n+1}^2 + \gamma_{n+1}^2| + |\eta_{n+1}^2|$  E.11

Therefore, in calculating  $|\alpha| < |\alpha_n|$ , one loop of relative minima

which implies that there is a point  $(\xi, \eta, \zeta)$  of  $\mathcal{K}_0$  for which

$(\xi_n, \eta_n, \zeta_n)$  is periodic  $|\xi^2 + \eta^2| \leq |\xi_{n+1}^2 + \eta_{n+1}^2|, \gamma_n, \alpha_n$  is also periodic.

We now see that  $(\beta, \gamma, \alpha)$  satisfies the which contradicts the definition of  $\{\xi_n, \eta_n, \zeta_n\}$  as a sequence of relative minima of  $\mathcal{K}_0$ . lattice,  $(\beta, \gamma, \alpha)$ , not equal to  $(0, 0, 0)$ .

Now let

$$M = \max_n |\zeta_n (\xi_{n+1}^2 + \eta_{n+1}^2)|$$

and

$$M_T = \max_n |\alpha_n (\beta_{n+1}^2 + \gamma_{n+1}^2)|$$

thus, by definition  $= \max_n |(\bar{\phi} - \phi)(\theta - \bar{\phi})(\phi - \theta)| |\zeta_n (\xi_{n+1}^2 + \eta_{n+1}^2)|$

of  $u, v, w$ , not  $= |\Delta| \max_n |\zeta_n (\xi_{n+1}^2 + \eta_{n+1}^2)|$

$$|M_T| \leq |\Delta| M, \quad |M_T| \leq q.$$

2.6 Define  $E$  by the following locus corresponding to locus 12

$$E(\xi, \eta, \zeta) = (\xi', \eta', \zeta')$$

when  $\zeta = \zeta'$

$$\text{Let } \epsilon > 1 \quad e'(\xi + i\eta) = (\xi' + i\eta')$$

$$e''(\xi - i\eta) = (\xi' - i\eta'), \text{ giving integer values}$$

then for some rational integer  $f > 0$  and every rational integer  $n$

$$(1) \quad E(\xi_n, \eta_n, \zeta_n) = (\xi_{n+f}, \eta_{n+f}, \zeta_{n+f}) \quad (\text{see chapter 1})$$

so that  $(11) \quad \zeta_n = \zeta_{n+f}$

and if  $(112) \quad \rho_\epsilon^2 = e'e''$  E.12

$$(12) \quad \rho_\epsilon^2 (\xi_n^2 + \eta_n^2) = (\xi_{n+f}^2 + \eta_{n+f}^2)$$

for every rational integer  $n$ .

Thus we have

$$\begin{aligned} |\zeta_{n+f} (\xi_{n+f}^2 + \eta_{n+f}^2)| &= |\zeta_n \rho_\epsilon^2 (\xi_{n+1}^2 + \eta_{n+1}^2)| \\ &= |\zeta_n (\xi_{n+1}^2 + \eta_{n+1}^2)| \end{aligned}$$

hence  $M = \max_{0 \leq n \leq n_0 + f - 1} |\xi_n (\xi_{n+1}^2 + \eta_{n+1}^2)| \cdot \quad 2.11$

Therefore, in calculating  $M$  only one loop of relative minima of  $K$  need be considered; we note that, since the sequence  $\{\xi_n, \eta_n, \zeta_n\}$  is periodic, the sequence  $\{\beta_n, \gamma_n, \alpha_n\}$  is also periodic.

We now see that the sequence  $\{\beta_n, \gamma_n, \alpha_n\}$  satisfies the conditions of lemma 2.1, hence, given  $p, q$  such that  $pq^2 \geq M_T$ , there is a point of the lattice,  $(\beta, \gamma, \alpha)$ , not equal to  $(0, 0, 0)$ , such that

$$|\beta^2 + \gamma^2| \leq q^2$$

$$|\alpha| \leq p;$$

thus, by definition of  $(\beta, \gamma, \alpha)$ , there exist integer values

of  $u, v, w$ , not all zero, such that

$$|M_1| \leq p, \quad |M_2| = |M_3| \leq q.$$

2.7 We now have the following lemma corresponding to lemma 12 of (5). In the definitions 2.5, such that

LEMMA 2.2  $\|x_n, y_n, z_n\| \geq \frac{k-2}{2(k-1)}$

for Let  $k > 1$  be given. There is an infinite set of values  $\alpha_n, \mu_n, \nu_n = \bar{\mu}_n$  of  $M_1, M_2, M_3$  corresponding to integer values  $u_n, v_n, w_n$  of  $u, v, w$ , as used in the expressions 2.6, such that

(i)  $|\alpha_n \mu_n^2| \leq \Delta |M|$

(ii)  $k |\alpha_n| \leq |\alpha_{n-1}|$

(iii)  $|\mu_n^2| |\alpha_{n-1}| \leq k \Delta |M|$

(iv)  $|\mu_n| \geq |\mu_{n-1}|$

2.12

so that  $qr^2 = k, q + 2q^2 = 2r + r^2 = 2^2 \frac{(1+k^2+k)}{(k+k^2)^2} = \lambda$  (asy).

PROOF OF LEMMA 2.2

By construction of  $M_T$  there is a set of values  $\alpha_0, \mu_0, \nu_0$  satisfying (i). Now define  $\alpha_n, \mu_n, \nu_n$  by induction, given

$\alpha_{n-1}, \mu_{n-1}, \nu_{n-1}$

$$|\alpha_n| \leq k^{-1} |\alpha_{n-1}| \quad \text{hence (ii) is satisfied}$$

$$|\mu_n| \leq \frac{k|\Delta|M}{|\alpha_{n-1}|} \quad \text{hence (iii) is satisfied.}$$

$\alpha_n, \mu_n, \nu_n$  exist since  $k^{-1} |\alpha_{n-1}| \frac{k|\Delta|M}{|\alpha_{n-1}|} = |\Delta|M$ .

At each stage the smallest  $\mu_n$  satisfying (iv) is chosen, otherwise  $\mu_n$  instead of  $\mu_{n-1}$  would have been chosen at the previous stage.

We now have the following two lemmas, reproduced from lemmas 13 and 14 of (5).

LEMMA 2.3

If  $k > 2$ , we may find values  $x_0, y_0, z_0$  of  $x, y, z$ , as used in the definitions 2.5, such that

$$\|x_0 u_n + y_0 v_n + z_0 w_n\| \geq \frac{k-2}{2(k-1)}$$

for all  $n$ ; where  $u_n, v_n, w_n$  are the values of  $u, v, w$ , in the definitions 2.6, corresponding to the values  $\alpha_n, \mu_n, \nu_n$  of  $M_1, M_2, M_3$ , and

$$\|x\| = \min |m - x|$$

where the minimum is over rational integer values of  $m$ .

LEMMA 2.4

Suppose that  $k > 2$

$r = \left(\frac{k+k^i}{2}\right)^{\frac{1}{i}}$ ,  $q = kr^{-2}(\sqrt{3+1})^2$ , so we have

$$\text{so that } qr^2 = k, \quad q + 2q^{\frac{1}{i}} = 2r + r^{-2} = 2^{\frac{1}{i}} \frac{(1+k^{\frac{1}{i}}+k)}{(k+k^{\frac{1}{i}})^{\frac{1}{i}}} = \lambda \quad (\text{say}).$$



If  $l, m, p$  are positive numbers such that

$$lm^2 \leq p^3, \quad l \leq qp, \quad m \leq rp$$

then  $l + 2m \leq \lambda p$ .

2.8 Now let  $\zeta, \xi + i\eta, \xi - i\eta$  be values taken by  $L_1, L_2, L_3$  for some

$$(x, y, z) \equiv (x_0, y_0, z_0) \pmod{1},$$

and let  $r$  and  $q$  be defined in terms of  $k$  as for lemma 2.4.

Choose  $n$  such that

$$|\mu_n^2| \leq r^2 |\zeta|^{\frac{2}{3}} |\xi^2 + \eta^2|^{\frac{1}{3}} (|\Delta|M)^{\frac{1}{3}} \leq |\mu_{n+1}^2|,$$

then by (iii) of lemma 2.2

$$|\alpha_n| \leq \frac{k|\Delta|M}{|\mu_{n+1}^2|} \leq kr^{-2} |\zeta|^{\frac{2}{3}} |\xi^2 + \eta^2|^{\frac{1}{3}} (|\Delta|M)^{\frac{1}{3}}$$

$$|\alpha_n| \leq q |\zeta|^{\frac{2}{3}} |\xi^2 + \eta^2|^{\frac{1}{3}} (|\Delta|M)^{\frac{1}{3}}.$$

By applying lemma 2.4 with

$$l = |\zeta \alpha_n|, \quad m = |(\xi + i\eta)\mu_n| \quad \text{and} \quad p = |\zeta|^{\frac{2}{3}} |\xi^2 + \eta^2|^{\frac{1}{3}} (|\Delta|M)^{\frac{1}{3}}$$

we have

$$|\zeta \alpha_n| + 2|(\xi + i\eta)\mu_n| \leq \lambda |\zeta|^{\frac{2}{3}} |\xi^2 + \eta^2|^{\frac{1}{3}} (|\Delta|M)^{\frac{1}{3}},$$

but

$$|\Delta| |x_n u_n + y_n v_n + z_n w_n| = |\zeta \alpha_n + (\xi + i\eta)\mu_n + (\xi - i\eta)\nu_n|$$

$$\leq |\zeta \alpha_n| + 2|(\xi + i\eta)\mu_n|;$$

thus, since

$$(1) \quad \|x_0 u_n + y_0 v_n + z_0 w_n\| \geq \frac{k-2}{2(k-1)} \quad \text{for all } n,$$

$$|\zeta(\xi^2 + \eta^2)| \geq \frac{|\Delta|^3}{|\Delta|M} \left( \frac{(k-2)}{2(k-1)\lambda} \right)^3 = \frac{|\Delta|}{M} \frac{(k-2)^3 (k+k^i)^2}{2^5 (k-1)^3 (1+k^i+k)^3}$$

$$(11) \quad \alpha = \frac{(\phi-\phi)}{12} \zeta; \quad \beta + i\delta = \frac{(\phi-\phi)}{12} (\xi + i\eta) = \frac{(\phi-\phi)}{12} (\xi - i\eta) \quad \Rightarrow \quad = \frac{|\Delta| f(k)}{M 2^5}.$$

$f(k)$  attains a maximum,  $F$ , when  $k = (\sqrt{3} + 1)^2$ , so we have

$$|\zeta(\xi^2 + \eta^2)| \geq \frac{|\Delta|}{M} \frac{(2 + 2\sqrt{3})^3 (5 + 3\sqrt{3})^2}{2^5 (3 + 2\sqrt{3})^3 (6 + 3\sqrt{3})^3} = \frac{2|\Delta|}{81(9 + 5\sqrt{3})M} \geq \frac{|\Delta|}{M \cdot 715.24 \dots}$$

that is the result 2.2.

2.9 For the second part of the result we note that in lemma 11 of (5) it is stated that:

if  $p > 0, q > 0$  are given and  $pq^2 \geq |\Delta|^2 3^{-\frac{1}{2}}$ , then there exist integer values  $u, v, w$  (as in 2.10), not all zero, such that

$$|M_1| \leq p, \quad |M_2| = |M_3| \leq q;$$

but from lemma 2.1, if  $p > 0, q > 0$  are given and  $pq^2 \geq M_T = |\Delta|M$ , there exist integer values  $u, v, w$ , not all zero, such that

$$|M_1| \leq p, \quad |M_2| = |M_3| \leq q.$$

But lemma 2.1 also states that  $M_T = |\Delta|M$  is best possible as a lower bound of  $pq^2$ , thus we have

$$|\Delta|M \leq |\Delta|^2 3^{-\frac{1}{2}}$$

hence

$$\frac{2|D|}{81(9 + 5\sqrt{3})M} = \frac{|D|F}{M \cdot 2^5} = \frac{|\Delta|^3 F}{M \cdot |\Delta| \cdot 2^5} \geq \frac{|\Delta|^3 F}{|\Delta|^2 3^{-\frac{1}{2}} 2^5} = \frac{|\Delta| 3^{\frac{1}{2}} F}{2^5} = \frac{2|\Delta|}{81(5 + 3\sqrt{3})}$$

that is

$$\frac{2|D|}{81(9 + 5\sqrt{3})M} \geq \frac{2|\Delta|}{81(5 + 3\sqrt{3})}$$

which is the result 2.3.

2.10 I return now to the case when  $\ell$  may have any value. We note, from 2.10, that  $(\beta, \delta, \alpha)$  is a point of  $K_\tau$  if and only if there is a point  $(\xi, \eta, \zeta)$  in  $K$ , such that

$$(i) \quad \zeta = x + y\theta + z\theta^2 = x' + y'\theta + zt_1\lambda$$

for some rational integers  $x, y, x', y', z$

$$(ii) \quad \alpha = \frac{(\bar{\phi} - \phi)}{i\ell} \zeta; \quad \beta + i\delta = \frac{(\theta - \bar{\phi})}{i\ell} (\xi + i\eta); \quad \beta - i\delta = \frac{(\phi - \theta)}{i\ell} (\xi - i\eta).$$

We now define  $K^*$  to be that subset of  $K$  which has basis  $(1, \theta, \theta^2)$ , that is the subset with basis  $(1, \theta, \ell\lambda)$ .  $K^*$  is

also a field, and the sublattice  $\mathcal{K}_0^*$  of  $\mathcal{K}_0$  which consists of those points  $(\xi, \eta, \zeta)$ , where  $\zeta \in K^*$ , is a multiplicative lattice.  $\zeta = x + y\theta + z\lambda$  is said to be an integer of  $K^*$  when  $x, y, z$  are all rational integers; and  $(\xi, \eta, \zeta)$  is said to be an integer point of  $\mathcal{K}_0^*$  when  $\zeta$  is an integer of  $K^*$ . Since  $\mathcal{K}_0^*$  is a multiplicative lattice, we may find an infinite sequence of relative minima  $\{\zeta_n^*\}$  in  $K^*$  which corresponds to an infinite sequence of integer points of  $\mathcal{K}_0^*$ ,  $\{\xi_n^*, \eta_n^*, \zeta_n^*\}$ , for which

$$|\zeta_{n-1}^*| > |\zeta_n^*|$$

Suppose

$$|\xi_{n-1}^{*2} + \eta_{n-1}^{*2}| < |\xi_n^{*2} + \eta_n^{*2}|,$$

and a corresponding sequence,  $\{\beta_n, \gamma_n, \alpha_n\}$ , in  $\mathcal{K}_T$  such that

$$|\alpha_{n-1}| > |\alpha_n|$$

$$|\beta_{n-1}^2 + \gamma_{n-1}^2| < |\beta_n^2 + \gamma_n^2|.$$

However, this sequence in  $\mathcal{K}_T$  is not necessarily a sequence of relative minima; nor can it be shown to be periodic, since the field  $K^*$  cannot be shown to have a fundamental automorphism.

Following the arguments of 2.5, 2.7 and 2.8 with  $M$  replaced by  $M^*$ , where

$$M^* = \max_n |\zeta_n^* (\xi_{n+1}^{*2} + \eta_{n+1}^{*2})|$$

gives by

so that

$$M_T = \max_n |\alpha_n (\beta_{n+1}^2 + \gamma_{n+1}^2)|$$

$$= \frac{|\Delta| M^*}{\ell^3} = \frac{|\Delta| M^*}{\ell^2},$$

since  $\prod \theta_j$  is the fundamental unit of the field

and eventually we obtain

$$|\zeta(\xi^2 + \eta^2)| \geq \frac{2|D|\ell^2}{M^* 81(\theta + 5\sqrt{3})}, \quad \text{the result 2.4.}$$

is also a complete loop of relative minima of the field when  $j$

2.11 For computing this method, the sequence of numbers  $\theta_9^{(1)}, \dots, \theta_9^{(k)}$



corresponding to the ideal of integers of the field, are calculated as described in chapter 1. We now observe that, in obtaining the products of these numbers to give the relative minima of the field, if  $r_1\theta^2 + r_2\theta + r_3$ , where  $r_1, r_2, r_3$  are rational integers, is such a relative minimum, the coefficients  $r_1, r_2, r_3$  become large in absolute value as the relative minima become closer to  $\epsilon$ . This fact causes loss of accuracy in the calculation of  $\alpha_n, \beta_n, \gamma_n$  and so in the calculation of  $M$ , even with double precision arithmetic.

Suppose

$$\epsilon = i_\epsilon\theta^2 + j_\epsilon\theta + k_\epsilon$$

then

$$\begin{aligned} \epsilon^{-1} &= \epsilon'\epsilon'' = (i_\epsilon\theta^2 + j_\epsilon\theta + k_\epsilon)(i_\epsilon\bar{\theta}^2 + j_\epsilon\bar{\theta} + k_\epsilon) \\ &= i_\epsilon^2(\theta^2 - a\theta + b)^2 + j_\epsilon^2(\theta^2 - a\theta + b) + k_\epsilon^2 \\ &\quad + i_\epsilon j_\epsilon(\theta^2 - a\theta + b)(a - \theta) \\ &\quad + i_\epsilon k_\epsilon(\theta^2 - a\theta + b) + j_\epsilon k_\epsilon(a - \theta). \end{aligned}$$

From this expression we see that if the coefficients of  $1, \theta, \theta^2$  in the expression for  $\epsilon$  are  $O(n)$  then those in the expression for  $\epsilon^{-1}$  are  $O(n^2)$ .

2.12 We have a complete loop of relative minima of the field given by

$$1, \theta_9^{(1)}, \theta_9^{(2)}\theta_9^{(1)}, \dots, \prod_{i=1}^k \theta_9^{(i)}$$

since  $\prod_{i=1}^k \theta_9^{(i)}$  is the fundamental unit of the field

$$\left(\prod_{i=j+1}^k \theta_9^{(i)}\right)^{-1}, \left(\prod_{i=j+2}^k \theta_9^{(i)}\right)^{-1}, \dots, \left(\theta_9^{(k)}\right)^{-1}, 1, \theta_9^{(1)}, \dots, \prod_{i=1}^j \theta_9^{(i)}$$

is also a complete loop of relative minima of the field when  $j$  is an integer satisfying  $0 \leq j \leq k$ .

We choose  $j$  so that the order of magnitude of the coefficients of  $1, \theta, \theta^2$  in the expression for  $\prod_{i=1}^j \theta_g^{(i)}$  and in that for  $(\prod_{i=1}^k \theta_g^{(i)})^{-1}$  will be approximately the same. In view of the result stated above for the coefficients in the expressions for  $\epsilon$  and  $\epsilon^{-1}$ , we choose  $j = \lfloor \frac{2k}{3} \rfloor$  where  $\lfloor i \rfloor$  represents the greatest integer less than  $i$ .

When this particular loop of relative minima has been calculated, we may calculate the  $\alpha_n$  and  $\beta_n$  and so  $M$ ; the loss of accuracy when using double precision arithmetic is not now significant. The multiplication and division routines used in calculating the relative minima check for the possibility of overflow, which in this case is interpreted as the result yielding an integer with more digits than double precision arithmetic allows significant figures.

In this way we find a lower bound on the inhomogeneous minimum of the field; for simplicity, in the program the constant 715.24... is replaced by 720.

$$\Lambda = \left\{ \xi_1 + \frac{1}{3} \alpha \xi_2^n \mid \xi_1 \in \mathbb{Z}, \xi_2 \in \mathbb{Z}, \xi_1 + \xi_2 = 1, \xi_1 \neq 0 \right\}$$

If  $\xi = \Lambda$  then no conclusion is reached about the existence of a Euclidean algorithm in  $K_\alpha$ . If there is a

A METHOD WHICH USES CONGRUENCES TO FIND A POINT  $\alpha$  IN  $K$  FOR WHICH  $N(K, \alpha) \geq 1$ .

3.1 I.1 is equivalent to the statement that there is no integer of  $K$  with norm of absolute value less than  $|N(\beta)|$  congruent to  $\zeta$  modulo  $\beta$ .

The method of congruences which is used to prove that there is no Euclidean Algorithm in a particular field is based on this fact, and may be summarized in the following steps:

(I) An integer  $\beta$  is chosen and representatives of each of its non-zero residue classes are calculated; suppose these representatives are denoted by  $\zeta_1, \zeta_2, \dots, \zeta_{r_\beta}$  where  $r_\beta = |N(\beta)| - 1$ . Let  $Z = \{\zeta_j : 1 \leq j \leq r_\beta\}$  be the set of representatives.

(II) For each ideal  $\mathfrak{f}$  of norm of absolute value less than  $|N(\beta)|$ , an integer  $\alpha$  which produces it is calculated; we then have  $\mathfrak{f} = (\alpha)$ . Let the set of these integers be  $\{\alpha_j : 1 \leq j \leq m\}$ , then any integer of norm of absolute value less than  $|N(\beta)|$  is of the form  $\alpha_j \epsilon^w$  for some  $j$  satisfying  $1 \leq j \leq m$  and some rational integer  $w$ .

(III) If  $h$  is the smallest positive rational integer for which  $\epsilon^h \equiv 1 (\beta)$ , we find the set  $\Lambda$  defined by

$$\Lambda = \{\zeta_j : \zeta_j \equiv \alpha_j \epsilon^w (\beta), 1 \leq j \leq m, 0 \leq w \leq h - 1, \zeta_j \in Z\}.$$

If  $Z = \Lambda$  then no conclusion is reached about the existence of a Euclidean Algorithm in  $K$ . If there is a

representative  $\zeta = \zeta_L$  where  $1 \leq L \leq r_p$  and  $\zeta$  is in  $(Z - \Lambda)$ , there is no Euclidean Algorithm in  $K$ .

The remainder of this chapter describes the method of computing this algorithm.

3.2 An underlying problem of this method is to find, for any given positive rational integer  $n$ , the distinct ideals of norm  $n$  and the integers which produce them.

If  $n = p$ , a rational prime, the ideals and corresponding integers are found as described in chapter 1.

We now consider the case when  $n = p^r$ , where  $p$  is a rational prime and  $r$  is a positive rational integer not equal to 1.

It has been noted in chapter 1 that  $(p)$  has at most three distinct ideal factors; we also note that  $(p)$  has a square factor if and only if  $p$  divides  $D$ , the discriminant of the field. The possible factorizations of  $(p)$  and consequently the ideals of norm  $n$  are as follows

(i)  $(p)$  is a prime ideal in  $K$  hence there are no ideals of norm  $p$ ; this may occur only if  $p$  does not divide  $D$ .

There will only be an ideal of norm  $n$  if  $r$  is a multiple of 3, when the only such ideal is  $(p^{\frac{r}{3}})$ ; if  $r$  is not a multiple of 3,  $(n)$  is a prime ideal in  $K$ .

(ii)  $(p) = \mathfrak{f} \cdot \mathfrak{q}$  where  $\mathfrak{f}$  is an ideal of norm  $p$  and  $\mathfrak{q}$  is a prime ideal of norm  $p^2$ ; this may occur only if  $p$  does not divide  $D$ . There is one ideal  $(\alpha)$  of norm  $p$  and one ideal  $(\beta)$  of norm  $p^2$ ; the ideals of norm  $p^r$

are  $(\alpha^u \beta^{\frac{r-u}{2}})$  where 2 divides  $r - u$  and  $r > u \geq 0$ .

(iii)  $(p) = \mathfrak{f}^3$  where  $\mathfrak{f}$  is an ideal of norm  $p$ ; we only have this situation if  $p$  divides  $D$ . There is just one ideal  $(\alpha)$  of norm  $p$  and, consequently, just one ideal  $(\alpha^r)$  of norm  $p^r$ .

(iv)  $(p) = \mathfrak{f}_1 \mathfrak{f}_2^2$  where  $\mathfrak{f}_1$  and  $\mathfrak{f}_2$  are ideals of norm  $p$ ; this happens only if  $p$  divides  $D$ . There are two ideals  $(\alpha_1)$  and  $(\alpha_2)$  of norm  $p$  and the ideals of norm  $p^r$  are  $(\alpha_1^u \alpha_2^{r-u})$  where  $r \geq u \geq 0$ .

(v)  $(p) = \mathfrak{f}_1 \mathfrak{f}_2 \mathfrak{f}_3$  where  $\mathfrak{f}_1, \mathfrak{f}_2, \mathfrak{f}_3$  are ideals of norm  $p$ ; this occurs only if  $p \nmid D$ . There are three ideals  $(\alpha_1), (\alpha_2)$  and  $(\alpha_3)$  of norm  $p$ , the ideals of norm  $p^r$  are  $(\alpha_1^u \alpha_2^v \alpha_3^{r-u-v})$  where  $r \geq u \geq 0, r \geq v \geq 0, r \geq u + v \geq 0$ .

3.3 Finally we consider the case  $n = p_1^{r_1} p_2^{r_2} \dots p_u^{r_u}$  where  $p_1, \dots, p_u$  are rational primes and  $r_1, \dots, r_u$  are positive rational integers. The ideals of norm  $n$  are of the form  $(\alpha_1 \alpha_2 \dots \alpha_u)$  where  $(\alpha_i)$  is an ideal of norm  $p_i^{r_i}$  for  $i = 1, \dots, u$ ; all ideals of norm  $n$  are found by considering all products  $\alpha_1 \dots \alpha_u$  for all possible values of  $\alpha_1, \dots, \alpha_u$ .

For both of the cases where  $n$  is composite, having found the ideals of norm a prime  $p$ , and the integers which produce them, for all  $p$  less than  $n$ ; by using the expressions enumerated in section 3.2, it is possible to calculate all ideals of norm  $n$  and the integers which produce them.

3.4 Having chosen a value for  $\beta$  the next problem is to find



a complete set of residues modulo  $\beta$ .

LEMMA 3.1

Suppose that  $r$  is the smallest positive rational integer

for which  $r + 1 \equiv 0 \pmod{\beta}$ ,

where  $q$  is the smallest positive rational integer

for which  $(q + 1)\theta + w_1 \equiv 0 \pmod{\beta}$ , for

some  $w_1$  satisfying  $0 \leq w_1 \leq r$

and  $p$  is the smallest positive rational integer

for which  $(p + 1)\lambda + v_2\theta + w_2 \equiv 0 \pmod{\beta}$ ,

for some  $w_2$  satisfying  $0 \leq w_2 \leq r$

and some  $v_2$  satisfying  $0 \leq v_2 \leq q$ ,

then  $Z = \{\rho : \rho = u\lambda + v\theta + w; 0 \leq u \leq p, 0 \leq v \leq q, 0 \leq w \leq r\}$

is a complete set of residues modulo  $\beta$ .

PROOF OF LEMMA 3.1

Suppose that  $d_1\lambda + d_2\theta + d_3$  is some integer of the field, then

$$d_1\lambda + d_2\theta + d_3 \equiv d_1^*\lambda + d_2^*\theta + d_3^* \pmod{\beta}$$

for some rational integers  $d_2^*, d_3^*$  and some rational integer  $d_1^*$

satisfying  $0 \leq d_1^* \leq p$  since  $(p + 1)\lambda \equiv -v_2\theta - w_2 \pmod{\beta}$ ,

$$d_1^*\lambda + d_2^*\theta + d_3^* \equiv d_1^*\lambda + d_2^*\theta + d_3'' \pmod{\beta}$$

for some rational integer  $d_3''$  and some rational integer  $d_2^*$

satisfying  $0 \leq d_2^* \leq q$  since  $(q + 1)\theta \equiv -w_1 \pmod{\beta}$ ,

$$d_1^*\lambda + d_2^*\theta + d_3'' \equiv d_1^*\lambda + d_2^*\theta + d_3^* \pmod{\beta}$$

for some rational integer  $d_3^*$  which satisfies  $0 \leq d_3^* \leq r$  since

$r + 1 \equiv 0 \pmod{\beta}$ .

Thus  $d_1\lambda + d_2\theta + d_3 \equiv d_1^*\lambda + d_2^*\theta + d_3^*$  where  $0 \leq d_1^* \leq p$ ,

$0 \leq d_2^* \leq q, 0 \leq d_3^* \leq r$ ; that is every integer of the field is congruent to an integer of  $Z$ , thus  $Z$  contains a complete set of residues modulo  $\beta$ .

It remains to show that no two elements of  $Z$  are congruent to each other. Suppose

$$d_1^\lambda + d_2^\theta + d_3 \equiv e_1^\lambda + e_2^\theta + e_3 \pmod{\beta}$$

where  $0 \leq e_1 \leq d_1 \leq p$

$$0 \leq d_2 \leq q, 0 \leq e_2 \leq q$$

$$0 \leq d_3 \leq r, 0 \leq e_3 \leq r,$$

then  $(d_1 - e_1)\lambda + (d_2 - e_2)\theta + d_3 - e_3 \equiv 0 \pmod{\beta}$  3.1

and  $h$  is the smallest positive rational integer for which such a congruence is true. Hence

$$0 \leq d_1 - e_1 \leq p.$$

If  $d_2 \geq e_2$  and  $d_3 \geq e_3$  we have

$$0 \leq d_2 - e_2 \leq q$$

$$0 \leq d_3 - e_3 \leq r,$$

thus 3.1 implies

Now suppose that  $(\alpha_1), \dots, (\alpha_n)$  are the distinct ideals

$$d_1 = e_1, d_2 = e_2, d_3 = e_3$$

of norm  $\alpha$ , then every integer of norm  $\alpha$  is congruent to  $e_1^\lambda d_2^\theta d_3$  as a result of the definitions of  $p, q,$  and  $r$ .

If  $d_2 \geq e_2$  and  $d_3 < e_3$  we have

$$0 < r + 1 + (d_3 - e_3) \leq r$$

$$0 \leq d_2 - e_2 \leq q,$$

but  $(d_1 - e_1)\lambda + (d_2 - e_2)\theta + r + 1 + (d_3 - e_3) \equiv 0 \pmod{\beta}$

as a result of the definition of  $r$ , this is a contradiction.

If  $d_2 < e_2$  then

$$0 < q + 1 + (d_2 - e_2) \leq q$$

and  $w'$  may be chosen such that

(3) and (31),  $0 \leq w' \leq r$  for which  $M(K) = M(K, \alpha)$  satisfy

and  $w' \equiv w_1 + d_3 - e_3 \pmod{\beta}$  from the definition of  $r$ ,

but  $(d_1 - e_1)\lambda + (q + 1 + d_2 - e_2)\theta + w' \equiv 0 \pmod{\beta}$

as a result of the definitions of  $q$  and  $w_1$ . This again is a contradiction. Hence  $d_1 = e_1, d_2 = e_2, d_3 = e_3$ , thus,  $Z$  is a complete set of residue classes modulo  $\beta$ . of norm  $n$  are

$(\alpha_1), \dots, (\alpha_n)$ , any integer of norm  $n$  will be of the form  $\epsilon^k \alpha_i$  for some rational integers  $k$  and  $i$  such that  $1 \leq i \leq n$ . may be most efficient. For every number  $\alpha$  of  $K$  there is in considering the residue classes which contain integers of norm a positive rational integer  $h$  such that

$$\epsilon^h \alpha \equiv \alpha \pmod{\beta},$$

for  $0 \leq k \leq h-1$  and  $1 \leq i \leq n$ . However, if we restrict and  $h$  is the smallest positive rational integer for which such the values of  $\beta$  to factors of  $\epsilon - 1$  or  $\epsilon + 1$ , we only need a congruence is true. Hence

$$(\epsilon^h - 1)\alpha \equiv 0 \pmod{\beta};$$

the number of integers  $\beta$  considered for any one field, only thus, if  $\alpha = \frac{\xi}{\beta}$ ,  $\beta$  divides  $\epsilon^h - 1$ , that is

$$\epsilon^h \equiv 1 \pmod{\beta}.$$

the cube of a rational prime are used; the latter are included Now suppose that  $(\alpha_1), \dots, (\alpha_{m_n})$  are the distinct ideals so that any rational prime factors of  $\epsilon - 1$  or of  $\epsilon + 1$  are of norm  $n$ , then every integer of norm  $n$  is congruent to  $\epsilon^w \alpha_j$

considered. for some rational integers  $w$  and  $j$  satisfying  $0 \leq w \leq h-1$ ,

To prove the time spent in searching for uncovered residue  $1 \leq j \leq m_n$ . Thus, for given  $n$ , the smaller  $h$  the fewer classes becoming too large,  $\beta$  was further restricted so that incongruent integers there will be of norm  $n$ ; hence, the smaller  $|N(\beta)|$  was less than 500; this was found, by experiment, to be a reasonable limit, in particular since most fields for which  $h$  the fewer incongruent integers there will be with norm of absolute value less than  $|N(\beta)|$ . Thus the smaller  $h$  the there was a possibility of considering integers  $\beta$  with  $|N(\beta)|$  lower the probability will be of all the residue classes modulo greater than 500 gave overflow warnings. All routines involving  $\beta$  being covered.

3.6 For the real quadratic fields for which results are known ( (2) and (21) ), the numbers  $\alpha$  for which  $M(K) = M(K, \alpha)$  satisfy

significant figure  $\epsilon \alpha \equiv \alpha \pmod{1}$  (1)

or overflow, and  $\epsilon \alpha \equiv -\alpha \pmod{1}$  possibilities.

or, in one case,  $\epsilon^2 \alpha \equiv -\alpha \pmod{1}$ .

Thus the possible values of  $\beta$  are the factors of  $\epsilon - 1$ ,  $\epsilon + 1$

or of  $\epsilon^2 + 1$ . We note that if the ideals of norm  $n$  are

$(\alpha_1), \dots, (\alpha_u)$ , any integer of norm  $n$  will be of the form

$\pm \epsilon^k \alpha_i$  for some rational integers  $k$  and  $i$  such that  $1 \leq i \leq u$ .

In considering the residue classes which contain integers of norm

$n$  we must, therefore, consider every integer of the form  $\pm \epsilon^k \alpha_i$

for  $0 \leq k \leq h - 1$  and  $1 \leq i \leq u$ . However, if we restrict

the values of  $\beta$  to factors of  $\epsilon - 1$  or  $\epsilon + 1$ , we only need

consider the integers  $\pm \alpha_i$  for  $1 \leq i \leq u$ . In order to restrict

the number of integers  $\beta$  considered for any one field, only

those values of  $\beta$  the norm of which is a rational prime or

the cube of a rational prime are used; the latter are included

so that any rational prime factors of  $\epsilon - 1$  or of  $\epsilon + 1$  are

considered.

To prevent the time spent in searching for uncovered residue

classes becoming too large,  $\beta$  was further restricted so that

$|N(\beta)|$  was less than 500; this was found, by experiment, to

be a reasonable limit, in particular since most fields for which

there was a possibility of considering integers  $\beta$  with  $|N(\beta)|$

greater than 500 gave overflow warnings. All routines involving

multiplication or division where the numbers involved were likely

to cause overflow, in the sense that the integer answer could have

coefficients which consist of more digits than the number of

significant figures allowed, included tests for the possibility of overflow, and monitored such possibilities.

AN ADAPTATION OF A METHOD OF BARNES AND SWINERTON-DYER FOR CUBIC FIELDS WITH COMPLEX CONJUGATES.

4.1 This chapter describes a method of isolating those points  $\alpha$  in  $K$  for which  $N(K, \alpha)$  is greater than a chosen bound. By this means we may either show that  $N(K, \alpha) < 1$  for all  $\alpha$  in  $K$ , so that  $K$  has a Euclidean Algorithm; or we may find a point  $\alpha$  such that  $N(K) = N(K, \alpha)$ , in this way we determine the inhomogeneous minimum of the field and consequently whether or not it possesses a Euclidean Algorithm. The method described is an adaptation of that used by Barnes and Swinerton-Dyer in (2).

4.2 If a number  $\alpha$  in  $K$  is given by  $\alpha = d(\theta) = d_1\theta^2 + d_2\theta + d_3$ , where  $d_1, d_2, d_3$  are rational numbers, let  $\alpha$  be represented in the space  $\mathbb{R}$  by

$$\begin{aligned} \tilde{\alpha} &= ( \operatorname{Re} d(\phi), \operatorname{Im} d(\phi), d(\theta) ) && \text{(cartesian co-ordinates)} \\ &= ( |d(\phi)|, \arg d(\phi), d(\theta) ) && \text{(cylindrical polar} \\ &&& \text{co-ordinates)} \end{aligned}$$

then the isomorphism  $\alpha \leftrightarrow \tilde{\alpha}$  establishes an isomorphism  $K \leftrightarrow \mathbb{R}$ .

Now define

$$\begin{aligned} N(\tilde{\alpha}) &= N(\alpha) = d(\theta)d(\phi)d(\bar{\phi}) \\ &= d(\theta) \{ (\operatorname{Re} d(\phi))^2 + (\operatorname{Im} d(\phi))^2 \}, \end{aligned}$$

thus, generally, if  $\tilde{\alpha}$  is the point  $(\xi, \eta, \zeta)$  then

$$N(\tilde{\alpha}) = \zeta(\xi^2 + \eta^2).$$

If  $\alpha$  is an integer in  $K$ ,  $\tilde{\alpha}$  is said to be an integer point of  $\mathbb{R}$ . Let the transformation  $E$  on points of  $\mathbb{R}$  be defined by

$$E(\tilde{\alpha}) = \tilde{\beta} \quad \text{when } \epsilon\alpha = \beta \text{ where } \alpha, \beta \text{ are in } K;$$



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$$\begin{aligned} \tilde{\alpha} &= (\operatorname{Re} d(\phi), \operatorname{Im} d(\phi), d(\theta)) \quad (\text{cartesian co-ordinates}) \\ &= (|d(\phi)|, \arg d(\phi), d(\theta)) \quad (\text{cylindrical polar co-ordinates}) \end{aligned}$$

then the isomorphism  $\alpha \leftrightarrow \tilde{\alpha}$  establishes an isomorphism  $K \leftrightarrow \mathcal{A}$ .

Now define

$$\begin{aligned} N(\tilde{\alpha}) &= N(\alpha) = d(\theta)d(\phi)d(\bar{\phi}) \\ &= d(\theta)((\operatorname{Re} d(\phi))^2 + (\operatorname{Im} d(\phi))^2), \end{aligned}$$

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$$E(\tilde{\alpha}) = \tilde{\beta} \quad \text{when } \epsilon\alpha = \beta \quad \text{where } \alpha, \beta \text{ are in } K;$$

expressed in cylindrical polar co-ordinates,  $E$  is then the transformation

$$(\rho, \omega, \zeta) \rightarrow (\rho\rho_\epsilon, \omega + \omega_\epsilon, \zeta/\zeta_\epsilon)$$

where  $0 < \zeta_\epsilon = \epsilon < 1$ ,  $\rho_\epsilon = \sqrt{\epsilon'\epsilon'}$ , so that  $\rho_\epsilon^2 \zeta_\epsilon = 1$  and  $\omega_\epsilon = \arg \epsilon'$ .

4.3 If  $R$  is a set of points of  $\mathcal{A}$ , the statement  $\tilde{\alpha} \in R \pmod{1}$  means that, for some integer  $\tilde{\gamma}$  in  $\mathcal{A}$ ,  $\tilde{\alpha} - \tilde{\gamma}$  is in  $R$ .

**THEOREM 4.1**

Let  $R$  be a bounded point set in the space  $\mathcal{A}$  such that, for some given set  $R^*$  in  $\mathcal{A}$  and some given integer point  $\tilde{\gamma}$  in  $\mathcal{A}$ , any point  $\tilde{\alpha}$  in  $R$  has the property that either  $E(\tilde{\alpha}) \in R^* \pmod{1}$  or  $E(\tilde{\alpha}) - \tilde{\gamma} \in R$ ; and further that  $E^{-1}(\tilde{\alpha})$  is congruent to a point of  $R^*$  or of  $R$ . If  $\tilde{\alpha} \in R$  and  $E^n(\tilde{\alpha})$  is not congruent to a point of  $R^*$  for any  $n \leq 0$ ,  $\tilde{\alpha}$  is the fixed point  $\tilde{\beta}$  of  $E$  defined by

$$E(\tilde{\beta}) = \tilde{\beta} + \tilde{\gamma} \quad \tilde{\beta} = (\xi_\beta, \eta_\beta, \zeta_\beta) = (\rho_\beta, \omega_\beta, \zeta_\beta).$$

The generalized form of this theorem, for algebraic fields of any given degree, is due to Cassels (quoted in (2)); since a proof was not readily available I provide one for this particular case, analogous to that given for the quadratic case in (2).

**LEMMA 4.2**

Let  $S$  be the transformation

$$(\rho, \omega, \zeta) \rightarrow (\sigma\rho, \omega + \omega_\sigma, \zeta/\sigma^2)$$

where  $\sigma > 1$  and  $-\pi < \omega_\sigma \leq \pi$ ; and let  $R$  be a bounded point set. Suppose  $S^n(\tilde{\alpha}_1) \in R$  for all  $n \geq 0$ , then  $\tilde{\alpha}_1$  lies on the line  $\rho = 0$ , the origin  $0$  belongs to the closure  $\bar{R}$  of  $R$ ,

and  $S^n(\tilde{\alpha}_1) \rightarrow 0$  as  $n \rightarrow \infty$ . If  $S^n(\tilde{\alpha}_1) \in R$  for all  $n \leq 0$  then  $\tilde{\alpha}_1$  lies on the plane  $\xi = 0$  and  $S^n(\tilde{\alpha}_1) \rightarrow 0$  as  $n \rightarrow -\infty$ .

PROOF OF LEMMA 4.2

Let  $\tilde{\alpha}_1$  be the point  $(\rho_1, \omega_1, \xi_1)$  then

$$S^n(\tilde{\alpha}_1) = (\sigma^n \rho_1, \omega_1 + n\omega_\sigma, \xi_1/\sigma^{2n}).$$

Thus, since  $R$  is bounded,  $\sigma^n \rho_1$  is bounded as  $n \rightarrow \infty$ ; hence, since  $\sigma > 1$ , it follows that  $\rho_1 = 0$ . Also

$$S^n(\tilde{\alpha}_1) = (0, \omega_1 + n\omega_\sigma, \xi_1/\sigma^{2n}) \rightarrow 0 \quad \text{as } n \rightarrow \infty,$$

hence  $0 \in \bar{R}$ .

Similarly, if  $\tilde{\alpha}_2 = (\rho_2, \omega_2, \xi_2)$  then

$$S^n(\tilde{\alpha}_2) = (\sigma^n \rho_2, \omega_2 + n\omega_\sigma, \xi_2/\sigma^{2n})$$

and  $\xi_2/\sigma^{2n}$  is bounded as  $n \rightarrow -\infty$ , hence  $\xi_2 = 0$  and

$$S^n(\tilde{\alpha}_2) = (\sigma^n \rho_2, \omega_2 + n\omega_\sigma, 0) \rightarrow 0 \quad \text{as } n \rightarrow -\infty.$$

PROOF OF THEOREM 4.1

Suppose that  $\tilde{\alpha}_0 \in R$  and that  $E^n(\tilde{\alpha}_0)$  is not congruent to a point of  $R^*$  for any  $n \geq 0$ .

Let  $\tilde{\gamma}$  be the point  $(\xi_\gamma, \eta_\gamma, \zeta_\gamma) = (\rho_\gamma, \omega_\gamma, \xi_\gamma)$  and define the transformation  $F$  by

$$F(\tilde{\alpha}) = E(\tilde{\alpha}) - \tilde{\gamma}.$$

Then, since  $F(\tilde{\alpha}) \equiv E(\tilde{\alpha}) \pmod{1}$  for all  $\tilde{\alpha}$ ,  $F^n(\tilde{\alpha}_0)$  is not congruent to a point of  $R^*$  for any  $n \geq 0$ .

If the origin is now changed to  $\tilde{\beta}$ , that is  $\tilde{\alpha}' = \tilde{\alpha} - \tilde{\beta}$ , the transformation  $F$  becomes

$$(\rho', \omega', \xi') \rightarrow (\rho'_\epsilon, \omega' + \omega_\epsilon, \xi'_\epsilon)$$

$-\pi < \omega_\epsilon \leq \pi$ ,  $\rho_\epsilon^2 \xi_\epsilon = 1$  and  $0 < \xi_\epsilon < 1$  so that  $\rho_\epsilon > 1$  and  $\xi_\epsilon = 1/\rho_\epsilon^2$

By hypothesis

Let  $R_0, R_1, \dots, R_{n-1}$  be a finite number of bounded point sets.  $\tilde{\alpha}_0 \in R_0$  and some integer points  $\tilde{\delta}_1 \in R_1$  and  $\tilde{\delta}_2 \in R_2$  such that  $F(\tilde{\alpha}_0) = E(\tilde{\alpha}_0) - \tilde{\gamma} = \tilde{\delta}_1 \in R_1$  and  $F^2(\tilde{\alpha}_0) = F(\tilde{\delta}_1) = E(\tilde{\delta}_1) - \tilde{\gamma} = \tilde{\delta}_2 \in R_2$  and ultimately

$$\tilde{\delta}_n = F^n(\tilde{\alpha}_0) \in R_n \text{ for all } n \geq 0.$$

It now follows from lemma 4.2 that  $\tilde{\beta} \in \bar{R}$ ,  $\tilde{\delta}_n \rightarrow \tilde{\beta}$  as  $n \rightarrow \infty$  and  $\tilde{\alpha}_0$  lies on the line  $\rho' = 0$ , that is the line  $\xi - \xi_\beta = \eta - \eta_\beta = 0$ .

If  $\tilde{\alpha}$  is a point in  $R$  such that  $E^{-1}(\tilde{\alpha}) \in R^* \pmod{1}$  then there is an integer point  $\tilde{\mu}$  such that

$$E^{-1}(\tilde{\alpha}) = \tilde{\mu} + \tilde{\delta} \quad \text{where } \tilde{\delta} \in R,$$

thus

$$\tilde{\alpha} = E(\tilde{\mu}) + E(\tilde{\delta}).$$

Since  $\tilde{\alpha}$  is a point of  $R$  and, by hypothesis,  $E(\tilde{\delta}) - \tilde{\gamma} \in R$ , we have  $E(\tilde{\mu}) = -\tilde{\gamma}$  so that  $\tilde{\mu} = -E^{-1}(\tilde{\gamma})$  which is independent of  $\tilde{\alpha}$ . Hence there is an integer point  $\tilde{\mu}$  such that, for any given  $\tilde{\alpha}$  in  $R$ , either  $E^{-1}(\tilde{\alpha}) \in R^* \pmod{1}$  or  $E^{-1}(\tilde{\alpha}) - \tilde{\mu} \in R$ .

Thus, by similar reasoning to that used for  $n \geq 0$ , in which for  $n \leq 0$  we replace  $E$  by  $E^{-1}$  and  $\tilde{\delta}$  by  $\tilde{\mu}$ , we find that  $\tilde{\alpha}_0$  lies on the plane  $\xi - \xi_\beta = 0$ . Hence  $\tilde{\alpha}_0$  lies at the intersection of the line  $\xi - \xi_\beta = \eta - \eta_\beta = 0$  with the plane  $\xi - \xi_\beta = 0$ , thus  $\tilde{\alpha}_0 = \tilde{\beta}$ .

4.4 Theorem 4.1 may be generalized to the case of a finite number of bounded point sets as follows.

$$E^n(\tilde{\alpha}) = \tilde{\mu}_n + \tilde{\gamma}_n = E(\tilde{\gamma}_{n-1}) + \dots + E^{n-1}(\tilde{\gamma}_1)$$

$$E^n(\tilde{\alpha}) - \tilde{\gamma} = \tilde{\alpha}_n \in R_n$$

THEOREM 4.3 Suppose that  $\tilde{\alpha}$  is in  $R_0$  and  $E^n(\tilde{\alpha})$  is not congruent

to a point of  $R^*$ . Let  $R_0, R_1, \dots, R_{m-1}$  be a finite number of bounded point sets. Suppose that for some  $R^*$  and some integer points  $\tilde{\gamma}_1, \tilde{\gamma}_2, \dots, \tilde{\gamma}_m$  every point  $\tilde{\alpha}_i$  in  $R_i$  ( $i = 0, 1, \dots, m-1$ ) has the property that either

$$E(\tilde{\alpha}_i) \in R^* \pmod{1}$$

$$\text{or} \quad E(\tilde{\alpha}_i) - \tilde{\gamma}_{i+1} \in R_{i+1} \quad (\text{where } R_m \text{ is } R_0).$$

Let

$$\tilde{\delta} = \tilde{\gamma}_m + E(\tilde{\gamma}_{m-1}) + E^2(\tilde{\gamma}_{m-2}) + \dots + E^{m-1}(\tilde{\gamma}_1)$$

and let  $\tilde{\beta}$  be the fixed point of  $E^m$  defined by

$$E^m(\tilde{\beta}) = \tilde{\beta} + \tilde{\delta}.$$

Further suppose that any  $\tilde{\alpha}_i$  in  $R_i$  has the property that either  $E^{-1}(\tilde{\alpha}_i) \in R^* \pmod{1}$  or  $E^{-1}(\tilde{\alpha}_i) \in R_j \pmod{1}$  for some  $j$  satisfying  $0 \leq j \leq m-1$ . If  $\tilde{\alpha} \in R_0$  and  $E^n(\tilde{\alpha})$  is not congruent to a point of  $R^*$  for any  $n \geq 0$  then  $\tilde{\alpha}$  is the fixed point  $\tilde{\beta}$  of  $E^m$ .

#### PROOF OF THEOREM 4.3

First suppose that  $\tilde{\alpha}$  is in  $R_0$  and  $E^n(\tilde{\alpha})$  is not congruent to a point of  $R^*$  for any  $n \geq 0$ . By hypothesis

$$E(\tilde{\alpha}) = \tilde{\alpha}_1 + \tilde{\gamma}_1 \quad \tilde{\alpha}_1 \text{ in } R_1$$

$$E(\tilde{\alpha}_1) = \tilde{\alpha}_2 + \tilde{\gamma}_2 \quad \tilde{\alpha}_2 \text{ in } R_2$$

.....

$$E(\tilde{\alpha}_{m-1}) = \tilde{\alpha}_m + \tilde{\gamma}_m \quad \tilde{\alpha}_m \text{ in } R_m = R_0,$$

thus

$$E^m(\tilde{\alpha}) = \tilde{\alpha}_m + \tilde{\gamma}_m + E(\tilde{\gamma}_{m-1}) + \dots + E^{m-1}(\tilde{\gamma}_1)$$

$$E^m(\tilde{\alpha}) - \tilde{\delta} = \tilde{\alpha}_m \text{ in } R_0.$$



4.3 Now suppose that  $\tilde{\alpha}$  is in  $R_0$  and  $E^m(\tilde{\alpha})$  is not congruent to a point of  $R^*$  for any  $n \leq 0$ . By hypothesis for some  $i$ ,  $0 \leq i \leq m-1$ , there is an integer point  $\tilde{\mu}_i$  such that and for any  $E^{-1}(\tilde{\alpha}) = \tilde{\delta}_i + \tilde{\mu}_i$ , both in  $\tilde{\delta}_i$  in  $R_i$ , thus and only if  $\tilde{\alpha} = E(\tilde{\delta}_i) + E(\tilde{\mu}_i)$ .

Also, by hypothesis, covering may now be uncurved as follows

I. A fundamental  $E(\tilde{\delta}_i) = \tilde{\delta}_{i+1} \in R_{i+1}$  is chosen,

II. Either a) For every number  $\alpha$  in  $E$ , such that  $\tilde{\alpha}$  thus we have  $R_{i+1} = R_0$ , hence  $i = m-1$ ,

$$E(\tilde{\mu}_{m-1}) = -\tilde{\delta}_m,$$

and so

$$\tilde{\mu}_{m-1} = -E^{-1}(\tilde{\delta}_m)$$

which is independent of  $\tilde{\alpha}$ . Thus there is an integer  $\tilde{\mu}_{m-1}$  such that If this is possible,  $E$  has a Euclidean Algorithm,

$$E^{-1}(\tilde{\alpha}) = \tilde{\delta}_{m-1} + \tilde{\mu}_{m-1} \quad \tilde{\delta}_{m-1} \text{ in } R_{m-1}.$$

Similarly there are integers  $\tilde{\mu}_{m-2}, \dots, \tilde{\mu}_0$  such that

$$E^{-1}(\tilde{\delta}_{m-1}) = \tilde{\delta}_{m-2} + \tilde{\mu}_{m-2} \quad \tilde{\delta}_{m-2} \text{ in } R_{m-2}$$

.....

$$E^{-1}(\tilde{\delta}_1) = \tilde{\delta}_0 + \tilde{\mu}_0 \quad \tilde{\delta}_0 \text{ in } R_0$$

and we have

$$E^{-m}(\tilde{\alpha}) = E^{-(m-1)}(\tilde{\mu}_{m-1}) + E^{-(m-2)}(\tilde{\mu}_{m-2}) + \dots + \tilde{\mu}_0 + \tilde{\delta}_0.$$

Hence

$$E^{-m}(\tilde{\alpha}) \in R_0 \pmod{1}.$$

Thus, by applying theorem 4.1 with  $E$  replaced by  $E^m$ , theorem 4.3 follows.

and  $\min |N(\tilde{\beta} - \tilde{\delta})| = 0$  where the minimum is over integers  $\tilde{\delta}$  in  $\tilde{A}$ ,

4.5 A fundamental region  $\mathcal{J}$  of  $\mathcal{A}$  is one such that for every point  $\tilde{\alpha}_1$  of  $\mathcal{A}$  there is a point  $\tilde{\alpha}_2$  in  $\mathcal{J}$  such that

$$\tilde{\alpha}_1 \equiv \tilde{\alpha}_2 \pmod{1}, \quad < \epsilon$$

and for any two points  $\tilde{\alpha}_1, \tilde{\alpha}_2$  both in  $\mathcal{J}$  the congruence holds if and only if  $\tilde{\alpha}_1 = \tilde{\alpha}_2$ .

The method of covering may now be summarized as follows

- I. A fundamental region  $\mathcal{J}$  of  $\mathcal{A}$  is chosen.
- II. Either a). For every number  $\alpha$  in  $K$ , such that  $\tilde{\alpha}$  is in  $\mathcal{J}$ , it is shown that there exists an integer point  $\tilde{\gamma}$  in  $\mathcal{A}$  such that, for a given real

number  $C < 1$ ,

$$|N(\tilde{\alpha} - \tilde{\gamma})| \leq C \quad 4.1$$

If this is possible,  $K$  has a Euclidean Algorithm.

4.6 In order to do this it is sufficient to show that 4.1

holds for every point of  $\mathcal{J}$ .

or b). Regions  $R_0, R_1, \dots, R_{m-1}$  and integer points  $\tilde{\gamma}_1, \tilde{\gamma}_2, \dots, \tilde{\gamma}_m$  in  $\mathcal{A}$  are found such that

(i) the conditions of theorem 4.3 are satisfied

with

$$R^* = \mathcal{J} - \bigcup_{i=0}^{m-1} R_i.$$

(ii) if  $\tilde{\beta}$  is the fixed point of  $E^m$ ,

satisfying

$$E^m(\tilde{\beta}) = \tilde{\beta} + \tilde{\gamma}$$

and  $\min |N(\tilde{\beta} - \tilde{\delta})| = C$

where the minimum is over integers  $\tilde{\delta}$  in  $\mathcal{A}$ ,

(i) There are for every point  $\tilde{\alpha}$  in  $R^*$  for which  $\alpha$  is for  $i = 1, \dots, n$  in  $K$  such that, if  $\tilde{\alpha}$  is in  $R_{\mu_i}$  for some rational integer  $\mu_i$ ,  $\min |N(\tilde{\alpha} - \tilde{\delta})| < C$  (mod 1).

In which case, where the minimum is over integers  $\tilde{\delta}$  in  $\mathfrak{a}$ .

If II(b) is followed then  $C = \max \min |N(\tilde{\alpha} - \tilde{\delta})|$  where the minimum is over all integers  $\tilde{\delta}$  in  $\mathfrak{a}$  and the maximum is over points  $\tilde{\alpha}$  in  $\mathfrak{a}$  for which  $\alpha$  is in  $K$ ; thus  $C$  is the inhomogeneous minimum of  $K$ , and so the field has a Euclidean Algorithm if and only if  $C < 1$ .

The alternatives II(a) and II(b) are provided since if it is possible to follow II(a) the method is shorter, as no consideration of transformation by the fundamental unit is required.

4.6 In order to satisfy the conditions of II(a) or of II(b), the fundamental region is divided into sub-regions  $\mathfrak{F}_1, \dots, \mathfrak{F}_n$ . A finite set  $\mathfrak{J}$  of integer points of  $\mathfrak{a}$  is chosen. Then, for given  $C$  we establish for every  $\mathfrak{F}_i, i = 1, \dots, n$ , whether there is an integer point  $\tilde{\delta}$  in  $\mathfrak{J}$  such that

$$\text{for every point } \tilde{\alpha} \text{ in } \mathfrak{F}_i, |N(\tilde{\alpha} - \tilde{\delta})| < C. \quad 4.2$$

Now each region  $\mathfrak{F}_i$ , for which there is no such  $\tilde{\delta}$ , is itself subdivided and its sub-regions tested as were the regions  $\mathfrak{F}_i$ .

4.7 The fundamental region  $\mathfrak{F}$  is chosen in order to allow a simple and efficient algorithm for selecting the integer points in either II(a) will be satisfied or there will be regions  $R_1, \dots, R_m$  so that, for a given integer  $\tilde{\gamma}$  in  $\mathfrak{J}$  there is a high probability of the existence of a subregion  $\mathfrak{F}_i$  of  $\mathfrak{F}$  which satisfies Having calculated the transformations of points of  $R_1, \dots, R_m$  by  $E$ , we will have one or both of the following situations:

(i) There are regions  $R_{\mu_1}, \dots, R_{\mu_n}$  where  $1 \leq \mu_i \leq m$  for  $i = 1, \dots, n$  such that, if  $\tilde{\alpha}$  is in  $R_{\mu_i}$ , for some rational integer  $j_i$ ,  $E^{j_i}(\tilde{\alpha}) \equiv \tilde{\gamma} - \sum_{i=1}^m R_i \pmod{1}$ .

In which case, for  $i = 1, \dots, n$  there is an integer point  $\tilde{\delta}_{\mu_i}$ , not in  $\mathcal{J}$ , which satisfies 4.2 for  $\tilde{\alpha}$  in  $R_{\mu_i}$ .

(ii) For  $\nu = 1, 2, \dots, p$  there are regions  $R_{\nu,0}, \dots, R_{\nu,(m_\nu-1)}$  and integer points  $\tilde{\delta}_{\nu,1}, \dots, \delta_{\nu,m_\nu}$  where  $R_{\nu,i}$  belongs to

the set  $\{R_1, \dots, R_m\}$  such that if  $\tilde{\alpha}$  is in  $R_{\nu,i}$ ,  $E(\tilde{\alpha}) - \tilde{\delta}_{\nu,(i+1)}$  is in  $R_{\nu,(i+1)}$  where  $R_{\nu,m_\nu} = R_{\nu,0}$ ,

and if there are regions  $R_{\mu_j}$  satisfying (i),  $R_{\mu_j} \neq R_{\nu,i}$  for  $\nu = 1, \dots, p$ ,  $i = 0, \dots, m_\nu$ ,  $j = 1, \dots, n$ .

In the second case each of the sets of regions  $R_{\nu,0}, \dots, R_{\nu,(m_\nu-1)}$  satisfies theorem 4.3. Let  $\tilde{\beta}_\nu$  be the corresponding fixed point of  $E^{m_\nu}$ ,  $\tilde{\beta}_\nu$  in  $R_{\nu,0}$ ,  $\beta_\nu$  in  $K$ ; then, if we have

$$C = \max_{\nu} \min_{\tilde{\delta}} |N(\tilde{\beta}_\nu - \tilde{\delta})|$$

where  $1 \leq \nu \leq p$  and  $\tilde{\delta}$  is an integer in  $\mathcal{A}$ , where the maximum is attained for  $\nu = \nu_0$ , II(b) is satisfied with the regions

$R_{\nu_0,0}, \dots, R_{\nu_0,(m_{\nu_0}-1)}$ . It should be noted that the sets  $\{R_{\nu,0}, \dots, R_{\nu,(m_\nu-1)}\}$  for  $\nu = 1, \dots, p$  are not necessarily disjoint, and the value of  $\nu_0$  is not necessarily unique.

4.7 The fundamental region  $\mathcal{J}$  is chosen in order to allow a simple and efficient algorithm for selecting the integer points in  $\mathcal{J}$  so that, for a given integer  $\tilde{\gamma}$  in  $\mathcal{J}$  there is a high probability of the existence of a subregion  $\mathcal{J}_\nu$  of  $\mathcal{J}$  which satisfies 4.2. With this in mind,  $\mathcal{J}$  is chosen so that, if  $\tilde{\alpha}$  in  $\mathcal{J}$  is



given by  $\tilde{\alpha} = (\xi_\alpha, \eta_\alpha, \zeta_\alpha)$  and we define

$$D_\alpha = \max_{\alpha \in \mathbb{Z}^3} |(\xi_\alpha^2 + \eta_\alpha^2 + \zeta_\alpha^2)|,$$

then  $D_\alpha$  is as small as possible. The method used to do this is now described.

4.8 If  $(1, M(\theta), Q(\theta))$  is a basis for  $K$  where  $M$  is linear and monic in  $\theta$  and  $Q$  is a quadratic with coefficient  $\frac{1}{t}$  for  $\theta^2$ ,

$$[(0,0,0), (1,0,1), \begin{cases} (\operatorname{Re} M(\phi), \operatorname{Im} M(\phi), M(\theta)), \\ (\operatorname{Re} Q(\phi), \operatorname{Im} Q(\phi), Q(\theta)) \end{cases}]$$

is a basis for the lattice of integers of  $\mathcal{A}$ . The actual choice of  $M$  and  $Q$  is made so that the scalar product

$$(1,0,1) \cdot (\operatorname{Re} M(\phi), \operatorname{Im} M(\phi), M(\theta)) = \operatorname{Re} M(\phi) + M(\theta)$$

is a minimum over all possible choices of  $M$ , say for  $M'$ ; and the scalar products

$$(1,0,1) \cdot (\operatorname{Re} Q(\phi), \operatorname{Im} Q(\phi), Q(\theta)) = \operatorname{Re} Q(\phi) + Q(\theta)$$

and

$$\begin{aligned} & (\operatorname{Re} M'(\phi), \operatorname{Im} M'(\phi), M'(\theta)) \cdot (\operatorname{Re} Q(\phi), \operatorname{Im} Q(\phi), Q(\theta)) \\ &= \operatorname{Re} M'(\phi) \cdot \operatorname{Re} Q(\phi) + \operatorname{Im} M'(\phi) \cdot \operatorname{Im} Q(\phi) + M'(\theta) \cdot Q(\theta) \end{aligned}$$

are each a minimum over all possible choices of  $Q$ , say for  $Q'$ .

$\mathcal{F}$  is then chosen to be the set of points  $(\xi, \eta, \zeta)$  where

$$\xi = x + y \cdot \operatorname{Re} M'(\phi) + z \cdot \operatorname{Re} Q'(\phi)$$

$$\eta = y \cdot \operatorname{Im} M'(\phi) + z \cdot \operatorname{Im} Q'(\phi)$$

$$\zeta = x + y \cdot M'(\theta) + z \cdot Q'(\theta)$$

4.3

and  $-\frac{1}{2} < x \leq \frac{1}{2}, -\frac{1}{2} < y \leq \frac{1}{2}, -\frac{1}{2} < z \leq \frac{1}{2}.$



and We determine  $M'$  and  $Q'$  as follows. Suppose

$$M(\theta) = \theta - k$$

and  $Q(\theta) = \lambda - p\theta - q$ , the conditions stated above

then we require

$\text{Re}\phi + \theta - 2k$  a minimum over all possible rational integer values of  $k$ .

Let  $k_1 = [\text{Re}\phi + \theta]$

and  $k_2 = \begin{cases} k_1 & \text{if } k_1 \text{ is even} \\ k_1 + 1 & \text{if } k_1 \text{ is odd and } \text{Re}\phi + \theta > 0 \\ k_1 - 1 & \text{if } k_1 \text{ is odd and } \text{Re}\phi + \theta < 0, \end{cases}$

then  $k_2$  is even and is that even integer closest to  $\text{Re}\phi + \theta$ .

If  $k' = \frac{1}{2}k_2$ ,  $M'(\theta) = \theta - k'$  satisfies the conditions stated above

for  $M'(\theta)$ . To find the rational integer values of  $p$  and  $q$  such that  $Q(\theta)$  satisfies the conditions for  $Q'(\theta)$ , we require

$$\text{Re}\psi + \lambda - p(\text{Re}\phi + \theta) - 2q$$

and  $\text{Re}\psi(\text{Re}\phi - k') + \text{Im}\psi\text{Im}\phi + \lambda(\theta - k')$

$$- p(\text{Re}\phi(\text{Re}\phi - k') + \text{Im}\phi\text{Im}\phi + \theta(\theta - k'))$$

$$- q(\text{Re}\phi + \theta - 2k')$$

to be a minimum over all rational integers  $p$  and  $q$ . Suppose

$p_r$  and  $q_r$  are the real numbers which when substituted for  $p$

and  $q$  respectively cause the above two expressions to become

zero. Let

$$p_\epsilon = \begin{cases} \frac{1}{2} & \text{if } p_r > 0 \\ -\frac{1}{2} & \text{if } p_r < 0 \end{cases}$$

$$q_\epsilon = \begin{cases} \frac{1}{2} & \text{if } q_r > 0 \\ -\frac{1}{2} & \text{if } q_r < 0, \end{cases}$$

and  $|x| \leq \frac{1}{2}$ ,  $p' = [p_r + p_e] \leq \frac{1}{2}$

If  $\tilde{y}$  is the  $q' = [q_r + q_e]$ ;  $(\xi, \eta)$ , it is an integer point for  $Q(\theta) = \lambda - p'\theta - q'$  then satisfies the conditions stated above for  $Q'(\theta)$ .

4.9 The method of subdivision of  $\mathcal{F}$  is chosen to simplify the testing for an integer for which a sub-region satisfies 4.2. The smallest cuboid  $\mathcal{C}$  which contains all points of  $\mathcal{F}$ , and which has

faces parallel to the co-ordinate planes, is chosen.  $\mathcal{C}$  is then divided into  $10 \times 10 \times 10$  similar cuboids  $\mathcal{C}_{ijk}$ ,  $i = 1, \dots, 10$ ;  $j = 1, \dots, 10$ ;  $k = 1, \dots, 10$ . Each  $\mathcal{C}_{ijk}$  has a flag  $\mathcal{L}_{ijk}$  which is set to 2 for those regions which do not contain any point of  $\mathcal{F}$  and to 1 otherwise.  $\mathcal{L}_{ijk}$  is subsequently set to 0 when an integer point  $\tilde{y}$  is found for which  $\mathcal{C}_{ijk}$  satisfies

4.2. The choice of  $\mathcal{F}$  as a fundamental region such that if  $\tilde{y}$

is in  $\mathcal{F}$ . If the region  $\mathcal{C}^*$  to be tested is the set of points  $(\xi, \eta, \zeta)$

for which when considering the set of sub-regions of the original

region  $\mathcal{C}$ , we need only consider those points  $(\xi, \eta, \zeta)$ ,

given by 4.3,  $\xi_1 \leq \xi \leq \xi_2$ ,  $\eta_1 \leq \eta \leq \eta_2$ ,  $|x| \leq \frac{1}{2}$ ,  $|y| \leq \frac{1}{2}$ ,  $0 \leq z \leq \frac{1}{2}$ . When

subsequently  $\xi_1 \leq \xi \leq \xi_2$  sub-regions all points must be

then  $\mathcal{C}^*$  intersects with  $\mathcal{F}$  only if at least one of  $(\xi_i, \eta_j, \zeta_k)$

( $i = 1, 2$ ;  $j = 1, 2$ ;  $k = 1, 2$ ) lies in  $\mathcal{F}$ ; which is so if there

is a set  $i, j, k$  with  $i = 1$  or  $2$ ,  $j = 1$  or  $2$ ,  $k = 1$  or  $2$

such that

as on the  $\xi_i = x + y \cdot \text{Re } M'(\phi) + z \cdot \text{Re } Q'(\phi)$  not cause inaccuracies

in the result  $\eta_j = y \cdot \text{Im } M'(\phi) + z \cdot \text{Im } Q'(\phi)$  for decimal places;

as in the  $\zeta_k = x + y \cdot M'(\theta) + z \cdot Q'(\theta)$

where  $|x| \leq \frac{1}{2}$ ,  $|y| \leq \frac{1}{2}$ ,  $|z| \leq \frac{1}{2}$ .

If  $\tilde{\gamma}$  is the point  $(\xi_2, \eta_2, \zeta_2)$ , it is an integer point for which  $\mathcal{C}^*$  satisfies 4.2 only if, when we define

$$\xi_d = \max(|\xi_2 - \xi_1|, |\xi_1 - \xi_2|)$$

$$\eta_d = \max(|\eta_2 - \eta_1|, |\eta_1 - \eta_2|)$$

and that the analogous expressions hold for  $\zeta$ , we have

$$\xi_d(\xi_d^2 + \eta_d^2) < C.$$

After testing using all points of  $\mathcal{J}$ , those  $\mathcal{C}_{ijk}$  for which  $\mathcal{L}_{ijk}$  is equal to 1 are subdivided and tested as was  $\mathcal{C}$ . The process is continued, using further subdivisions, until those regions for which an integer point cannot be found to satisfy 4.2 can be combined to give regions  $R_0, \dots, R_{m-1}$  to satisfy II(b), or until II(a) is shown to be satisfied.

4.10 The choice of  $\mathcal{J}$  as a fundamental region such that if  $\tilde{\alpha}$  is in  $\mathcal{J}$ ,  $-\tilde{\alpha}$  is in  $\mathcal{J}$  has the following advantages

(i) When considering the set of sub-regions of the original region  $\mathcal{C}$ , we need only consider those points  $(\xi, \eta, \zeta)$ , given by 4.3, for which  $|x| \leq \frac{1}{2}$ ,  $|y| \leq \frac{1}{2}$ ,  $0 \leq z \leq \frac{1}{2}$ . When subsequently testing these sub-regions all points must be considered.

(ii) The quantities involved in the above expressions are all of the same order of magnitude, usually of absolute value less than 10; so with an accuracy of fifteen significant figures, as on the computer used, rounding error will not cause inaccuracies in the results when  $C$  is quoted to just four decimal places, as in the program used.

4.11 I now turn to the choice of the set  $\mathcal{J}$  of integer points.

Suppose that  $\mathcal{C}$  is the set of  $(\xi, \eta, \zeta)$  for which

$$|(\xi_u - \xi_l)((\xi_u - \xi_l)^2 + (\eta_u - \eta_l)^2)| < C \quad \text{for all } \tilde{\alpha} \text{ in } \mathcal{C}^{(i)}$$

$$|(\xi_u - \xi_l)| < \frac{h_u \leq \eta \leq h_l}{((\xi_u - \xi_l)^2 + (\eta_u - \eta_l)^2)} \quad \text{for all } \tilde{\alpha} \text{ in } \mathcal{C}^{(i)} \quad 4.5$$

$$|(\xi_u - \xi_l)^2 + (\eta_u - \eta_l)^2| < \frac{\zeta_u - \zeta_l}{(\xi_u - \xi_l)} \quad \text{for all } \tilde{\alpha} \text{ in } \mathcal{C}^{(i)} \quad 4.6$$

and that the cuboids  $\mathcal{C}^{(i)}$ ,  $i = 1, \dots, n$ , for which we are now trying to find integer points to satisfy 4.2, are the sets of  $(\xi, \eta, \zeta)$  for which

$$\xi_l^{(i)} \leq \xi \leq \xi_u^{(i)}$$

$$h_l^{(i)} \leq \eta \leq h_u^{(i)}$$

$$\zeta_l^{(i)} \leq \zeta \leq \zeta_u^{(i)}$$

where

$$\xi_u^{(i)} - \xi_l^{(i)} = 10^{-h} (\xi_u - \xi_l) = \xi_h \quad \text{say}$$

$$h_u^{(i)} - h_l^{(i)} = 10^{-h} (h_u - h_l) = h_h \quad \text{say}$$

$$\zeta_u^{(i)} - \zeta_l^{(i)} = 10^{-h} (\zeta_u - \zeta_l) = \zeta_h \quad \text{say}$$

$$\xi_l^{(i)} = m_{1i} \xi_h + \xi_l; \quad h_l^{(i)} = m_{2i} h_h + h_l; \quad \zeta_l^{(i)} = m_{3i} \zeta_h + \zeta_l$$

for some positive rational integer  $h$  and non-negative rational integers  $m_{1i}, m_{2i}, m_{3i}$ .

It is required that the set  $\mathcal{J}$  should contain those integer

points  $\tilde{\gamma} = (\xi_r, \eta_r, \zeta_r)$  such that  $\tilde{\beta} = (\xi, \eta, \zeta)$  where  $\tilde{\beta}$  is in  $\mathcal{C}^{(i)}$

and for some  $\mathcal{C}^{(i)}$ , all  $\tilde{\alpha} = (\xi_\alpha, \eta_\alpha, \zeta_\alpha)$  in  $\mathcal{C}^{(i)}$  satisfy  $|N(\tilde{\alpha} - \tilde{\gamma})| < C$

and as few other points as possible. Now to see when progressing,

$$|N(\tilde{\alpha} - \tilde{\gamma})| = |(\zeta_\alpha - \zeta_r)((\xi_\alpha - \xi_r)^2 + (\eta_\alpha - \eta_r)^2)|$$

thus bounds may be found for  $\xi_r, \eta_r, \zeta_r$ , for those  $\tilde{\gamma}$  for which

there is a high probability that 4.4 holds, in terms of the bounds

on  $\xi_\alpha, \eta_\alpha, \zeta_\alpha$  for  $\tilde{\alpha}$  in  $\mathcal{C}$ .

In order that  $M(\theta) = \epsilon, \theta(\theta)$ , with the result that

$$|(\xi_\alpha - \xi_\gamma)((\xi_\alpha - \xi_\gamma)^2 + (\eta_\alpha - \eta_\gamma)^2)| < C, \quad \text{for all } \tilde{\alpha} \text{ in } \mathcal{C}^{(i)}$$

$$|(\xi_\alpha - \xi_\gamma)| < \frac{C}{((\xi_\alpha - \xi_\gamma)^2 + (\eta_\alpha - \eta_\gamma)^2)} \quad \text{for all } \tilde{\alpha} \text{ in } \mathcal{C}^{(i)} \quad 4.5$$

$$|(\xi_\alpha - \xi_\gamma)^2 + (\eta_\alpha - \eta_\gamma)^2| < \frac{C}{|(\xi_\alpha - \xi_\gamma)|} \quad \text{for all } \tilde{\alpha} \text{ in } \mathcal{C}^{(i)} \quad 4.6$$

Thus we may either restrict the values of  $\xi_\gamma$  and  $\eta_\gamma$  and so find bounds for  $\xi_\gamma$ , or restrict the value of  $\xi_\gamma$  and so find bounds for  $\xi_\gamma$  and  $\eta_\gamma$ .

First of all we restrict the values of  $\xi_\gamma$  and  $\eta_\gamma$ , so that they are close to the values of  $\xi_\alpha$  and  $\eta_\alpha$  for  $\tilde{\alpha}$  in  $\mathcal{C}$ , by the inequalities

$$\xi_\alpha - 0.1 \leq \xi \leq \xi_\alpha + 0.1 \quad 4.7$$

$$\text{and} \quad \eta_\alpha - 0.1 \leq \eta \leq \eta_\alpha + 0.1$$

The constant 0.1 is chosen bearing in mind that the critical value for  $C$ , so far as the Euclidean Algorithm is concerned, is 1.

With  $\xi_\gamma, \eta_\gamma$  satisfying 4.7, a sufficient condition for  $\xi_\gamma$  satisfying 4.4 is

$$|(\xi_\alpha - \xi_\gamma)| < \frac{C}{\rho_n^2} \quad |N(\tilde{\beta} - \tilde{\gamma})| < C$$

where  $\rho_n^2 = \max |(\xi - \xi_\gamma)^2 + (\eta - \eta_\gamma)^2|$ ,

and the maximum is over all points  $\tilde{\beta} = (\xi, \eta, \zeta)$  where  $\tilde{\beta}$  is in  $\mathcal{C}^{(i)}$  and  $|N(\tilde{\beta} - \tilde{\gamma})| < C$ . A necessary bound for  $\rho_n$  is then  $\frac{1}{2}(\xi_n^2 + \eta_n^2)^{\frac{1}{2}}$ .

$\frac{1}{2} \max\{\xi_n, \eta_n\}$  was in fact used as this lower bound since it was thought to be a more efficient expression to use when programming, although

$$\frac{1}{\sqrt{2}} \left( \frac{1}{2} (\xi_n^2 + \eta_n^2) \right)^{\frac{1}{2}} \leq \frac{1}{2} \max\{\xi_n, \eta_n\} \leq \frac{1}{2} (\xi_n^2 + \eta_n^2)^{\frac{1}{2}}$$

so that some extra points would be included in  $\mathcal{J}$ . However, it



was found necessary to impose bounds on the coefficients  $x_\nu, y_\nu, z_\nu$ , where  $\xi_\nu = x_\nu + y_\nu M'(\theta) + z_\nu Q'(\theta)$ ; with the result that  $\xi_\nu$  does not come near the bound calculated, using either of the above bounds for  $\rho_h$ , when  $C$  is close to 1. Thus a necessary condition for the integer point  $\tilde{\delta}$  to be included in the set  $\mathcal{J}$  is that when  $\tilde{\delta}$  satisfies 4.7 it further satisfies

$$|\xi_\alpha - \xi_\nu| < \frac{4C}{(\max\{\xi_h, \eta_h\})^2} \quad \text{for all } \tilde{\alpha} \text{ in } \mathcal{C}^{(i)}$$

for which a necessary condition is

$$\xi_\nu - \frac{4C}{(\max\{\xi_h, \eta_h\})^2} \leq \xi_\alpha \leq \xi_u + \frac{4C}{(\max\{\xi_h, \eta_h\})^2} \quad 4.8$$

We now find necessary conditions for  $\tilde{\delta}$  in  $\mathcal{J}$  by first restricting the value of  $\xi_\nu$  so that it is close to the values taken by  $\xi_\alpha$  for  $\tilde{\alpha}$  in  $\mathcal{C}$ , we use the inequality

$$\xi_\nu - 0.1 \leq \xi_\alpha \leq \xi_u + 0.1 \quad 4.9$$

With  $\xi_\nu$  satisfying 4.9 a sufficient condition for  $\xi_\nu, \eta_\nu$  satisfying 4.6 is

$$|(\xi_u - \xi_\nu)^2 + (\eta_u - \eta_\nu)^2| < \frac{C}{Z_h} \quad \text{where } Z_h = \max |\xi - \xi_\nu|$$

and the maximum is taken over all points  $\tilde{\beta} = (\xi, \eta, \zeta)$  where  $\xi_\nu^{(i)} \leq \xi \leq \xi_u^{(i)}, \eta_\nu^{(i)} \leq \eta \leq \eta_u^{(i)}$  and  $|N(\tilde{\beta} - \tilde{\delta})| < C$ . A necessary lower bound for  $Z_h$  is then  $\frac{1}{2} \xi_u$ . Thus a necessary condition for the integer point  $\tilde{\delta}$  to be included in  $\mathcal{J}$  is that when it satisfies 4.9 it further satisfies

$$|(\xi_u - \xi_\nu)^2 + (\eta_u - \eta_\nu)^2| < \frac{2C}{\xi_u}$$

for which necessary conditions are

$$\xi_\nu - \left(\frac{2C}{\xi_u}\right)^{\frac{1}{2}} \leq \xi_\alpha \leq \xi_u + \left(\frac{2C}{\xi_u}\right)^{\frac{1}{2}} \quad 4.10$$

and 
$$\eta_\nu - \left(\frac{2C}{\xi_u}\right)^{\frac{1}{2}} \leq \eta_\alpha \leq \eta_u + \left(\frac{2C}{\xi_u}\right)^{\frac{1}{2}}$$

4.12 Having found the regions  $R_1, \dots, R_m$  for which there is no integer point  $\tilde{\gamma}$  satisfying 4.2, it is necessary to determine with which of these regions the translates of the regions,  $E(R_1), \dots, E(R_m)$ , intersect.

We suppose that, for some  $n$  where  $1 \leq n \leq m$ ,  $R_n$  is the set of points  $(\xi, \eta, \zeta)$  which satisfy

$$\xi_1^{(n)} \leq \xi \leq \xi_2^{(n)}$$

$$\eta_1^{(n)} \leq \eta \leq \eta_2^{(n)}$$

$$\zeta_1^{(n)} \leq \zeta \leq \zeta_2^{(n)}$$

For each of these vertices,  $(\xi_i^{(n)}, \eta_j^{(n)}, \zeta_k^{(n)})$   $i = 1, 2, j = 1, 2, k = 1, 2$ , we find the cylindrical polar co-ordinates  $(\rho_{(i,j)}^{(n)}, \alpha_{(i,j)}^{(n)}, \zeta_k^{(n)})$ .

If the fundamental unit  $\epsilon$  has cylindrical polar co-ordinates

$(\rho_\epsilon, \alpha_\epsilon, \zeta_\epsilon)$ , the transform of a point  $(\rho, \alpha, \zeta)$  by  $\epsilon$  is

$(\rho\rho_\epsilon, \alpha + \alpha_\epsilon, \zeta\zeta_\epsilon)$ ; thus transforms of the vertices of  $R_n$  will

be  $(\rho_{(i,j)}^{(n)}\rho_\epsilon, \alpha_{(i,j)}^{(n)} + \alpha_\epsilon, \zeta_k^{(n)}\zeta_\epsilon)$  for  $i = 1, 2, j = 1, 2, k = 1, 2$ .

$E(R_n)$  is then the cuboid with these points as vertices. We

suppose now that  $E(R_n)$  is the set of points  $(\xi, \eta, \zeta)$  which satisfy

$$\xi_{1\epsilon}^{(n)} \leq \xi \leq \xi_{2\epsilon}^{(n)}$$

$$\eta_{1\epsilon}^{(n)} \leq \eta \leq \eta_{2\epsilon}^{(n)}$$

$$\zeta_{1\epsilon}^{(n)} \leq \zeta \leq \zeta_{2\epsilon}^{(n)}$$

We wish to find the set  $J$  of those points  $\tilde{\gamma}$  which cause

$E(R_n) - \tilde{\gamma}$  to intersect with the original fundamental region  $J$ ,

in order that we may find which of  $R_1, \dots, R_m$ , if any, intersect

with the translates of  $E(R_n)$ .

For each triplet  $(i, j, k)$  we find a triad  $(x_{ijk}^{(n)}, y_{ijk}^{(n)}, z_{ijk}^{(n)})$

such that

$$\xi_{i\epsilon}^{(s)} = x_{ijk}^{(s)} + y_{ijk}^{(s)} \operatorname{Re} M'(\phi) + z_{ijk}^{(s)} \operatorname{Re} Q'(\phi)$$

$$\eta_{j\epsilon}^{(s)} = y_{ijk}^{(s)} \operatorname{Im} M'(\phi) + z_{ijk}^{(s)} \operatorname{Im} Q'(\phi)$$

$$\zeta_{k\epsilon}^{(s)} = x_{ijk}^{(s)} + y_{ijk}^{(s)} M'(\theta) + z_{ijk}^{(s)} Q'(\theta).$$

The rational integers  $X_\ell, X_u, Y_\ell, Y_u, Z_\ell, Z_u$  are then defined

such that  $X_\ell, Y_\ell, Z_\ell$  are the greatest and  $X_u, Y_u, Z_u$  are

the least rational integers to satisfy

$$X_\ell \leq x_{ijk}^{(s)} \leq X_u \quad \text{and} \quad \text{etc.}$$

$$Y_\ell \leq y_{ijk}^{(s)} \leq Y_u$$

$$Z_\ell \leq z_{ijk}^{(s)} \leq Z_u$$

over all  $i, j, k$  where  $i = 1, 2, j = 1, 2, k = 1, 2$ . The possible

integers  $\tilde{\delta} = (\xi_x, \eta_x, \zeta_x)$  are then among those which satisfy

$$\xi_x = x + y \operatorname{Re} M'(\phi) + z \operatorname{Re} Q'(\phi) \quad i = 1, 2, j = 1, 2,$$

$$\eta_x = y \operatorname{Im} M'(\phi) + z \operatorname{Im} Q'(\phi) \quad i = 1, 2, j = 1, 2,$$

$$\zeta_x = x + y M'(\theta) + z Q'(\theta) \quad i = 1, 2, j = 1, 2,$$

where

$$X_\ell \leq x \leq X_u$$

$$Y_\ell \leq y \leq Y_u$$

$$Z_\ell \leq z \leq Z_u;$$

we define  $\mathcal{J}$  to be the set of integer points which satisfy these

conditions. For each integer  $\tilde{\delta}$  in  $\mathcal{J}$  we then test whether,

for  $p = 1, \dots, m$ ,  $R_p + \tilde{\delta}$  intersects with  $E(R_n)$ .

4.13 The program which performs this testing calculates the basis of the lattice and then reads the values of the coefficients of the vertices of  $R_1, \dots, R_m$  as data, in a format with ten

decimal places. To ensure that no points are lost due to rounding error the values of  $\xi_1^{(n)}$ ,  $\eta_1^{(n)}$  and  $\zeta_1^{(n)}$  for  $n = 1, \dots, m$  are decreased by  $10^{-10}$  and the values of  $\xi_2^{(n)}$ ,  $\eta_2^{(n)}$  and  $\zeta_2^{(n)}$  are increased by  $10^{-10}$  before any transformations are performed. The values of  $\rho_\epsilon$ ,  $\alpha_\epsilon$ ,  $\zeta_\epsilon$  will be accurate within the limits of the machine, that is fifteen significant figures, as the coefficients of  $\epsilon$  relative to the basis  $(1, M'(\theta), Q'(\theta))$  are read in as rational integers and  $\rho_\epsilon$ ,  $\alpha_\epsilon$ ,  $\zeta_\epsilon$  are determined using the basis which is calculated.

Since  $0 < \epsilon < 1$  we also have

$$\begin{aligned} \text{thus, if } \zeta &= \frac{\zeta}{\rho}, & \zeta_{1\epsilon}^{(n)} &= \zeta_1^{(n)} \zeta_\epsilon \\ & & \zeta_{2\epsilon}^{(n)} &= \zeta_2^{(n)} \zeta_\epsilon \\ & & \xi_{1\epsilon}^{(n)} &= \min_{i,j} (\rho_{(i,j)}^{(n)} \rho_\epsilon \cos(\alpha_{(i,j)}^{(n)} + \alpha_\epsilon)) \quad i = 1, 2, j = 1, 2. \\ & & \eta_{1\epsilon}^{(n)} &= \min_{i,j} (\rho_{(i,j)}^{(n)} \rho_\epsilon \sin(\alpha_{(i,j)}^{(n)} + \alpha_\epsilon)) \quad i = 1, 2, j = 1, 2. \\ & & \xi_{2\epsilon}^{(n)} &= \max_{i,j} (\rho_{(i,j)}^{(n)} \rho_\epsilon \cos(\alpha_{(i,j)}^{(n)} + \alpha_\epsilon)) \quad i = 1, 2, j = 1, 2. \\ & & \eta_{2\epsilon}^{(n)} &= \max_{i,j} (\rho_{(i,j)}^{(n)} \rho_\epsilon \sin(\alpha_{(i,j)}^{(n)} + \alpha_\epsilon)) \quad i = 1, 2, j = 1, 2. \end{aligned}$$

bounds are obtained by showing that values of  $\Re(\frac{\omega}{1-\epsilon^n} - \delta)$  are equal to those taken by a restricted set of  $\delta$ 's, since the norm remains unchanged under certain transformations of the lattice of integers.

D.2 In the trivial case when  $\frac{\omega}{1-\epsilon^n}$  is an integer then  $\delta = 0$ . In the following it will be assumed that  $\frac{\omega}{1-\epsilon^n}$  is not an integer of  $\mathbb{K}$ .

THE MINIMUM OF  $\frac{\alpha}{1 - \epsilon^n}$  transforms into itself if the elements

5.1 For a given number  $\frac{\alpha}{1 - \epsilon^n}$ , where  $\alpha$  is an algebraic integer, we wish to find the minimum  $C$  of  $|N(\frac{\alpha}{1 - \epsilon^n} - \delta)|$  for integers  $\delta$  in  $K$ ; from 1.3, if  $C$  is greater than or equal to 1 then the field is non-Euclidean. The special number  $\frac{\alpha}{1 - \epsilon^n}$  is chosen since for every number  $\zeta$  in  $K$  there is a positive rational integer  $n$  such that

$$\epsilon^n \zeta \equiv \zeta \pmod{1 - \epsilon^n} \quad (1)$$

thus, if  $\zeta = \frac{\alpha}{\beta}$ ,

$$(1 - \epsilon^n) \frac{\alpha}{\beta} \equiv 0 \pmod{1 - \epsilon^n} \quad (1)$$

hence  $\beta$  divides  $1 - \epsilon^n$ ; therefore, without loss of generality, we may restrict  $\beta$  to values of the form  $1 - \epsilon^n$ .

We obtain bounds for the coefficients  $p, q, r$  in the canonical representation of  $\delta$  as  $p\lambda + q\theta + r$ ; using these bounds we can conduct a computer search for possible  $\delta$ . The bounds are obtained by showing that values of  $N(\frac{\alpha}{1 - \epsilon^n} - \delta)$  are equal to those taken by a restricted set of  $\delta$ 's, since the norm remains unchanged under certain transformations of the lattice of integers.

5.2 In the trivial case when  $\frac{\alpha}{1 - \epsilon^n}$  is an integer then  $C = 0$ . In the following it will be assumed that  $\frac{\alpha}{1 - \epsilon^n}$  is not an integer of  $K$ .

$$A = \{ \delta \mid \delta \text{ integer in } K \text{ and } \tau_c \in \delta < \frac{\tau_c - \alpha}{\epsilon^n} \}$$

$$B = \{ \delta \mid \delta \text{ integer in } K \text{ and } \frac{\tau_c - \alpha}{\epsilon^n} < \delta < \tau_c \}$$



LEMMA 5.1 Let  $\alpha$  be an integer in  $K$  which satisfies

The number  $\frac{\alpha}{1-\epsilon^n}$  transforms into itself if the elements of  $K$  are transformed by  $T$  belonging to  $A$  for any rational integer  $i$ , in particular for  $T: \gamma \rightarrow \frac{\gamma-\alpha}{\epsilon^n}$  for every rational integer  $i$ . We have 5.1

PROOF OF LEMMA 5.1 trivial.  $T(\gamma) = \frac{\gamma-\alpha}{\epsilon^n}$

For any given integer  $\gamma$  in  $K$

$$\begin{aligned} N\left(\frac{\alpha}{1-\epsilon^n} - T(\gamma)\right) &= N\left(\frac{\alpha}{1-\epsilon^n} - \frac{\gamma-\alpha}{\epsilon^n}\right) \\ &= N\left(\frac{\alpha\epsilon^n + \alpha - \alpha\epsilon^n - \gamma(1-\epsilon^n)}{\epsilon^n(1-\epsilon^n)}\right) \\ &= N\left(\frac{1}{\epsilon^n} \left(\frac{\alpha}{1-\epsilon^n} - \gamma\right)\right) \\ &= N\left(\frac{\alpha}{1-\epsilon^n} - \gamma\right) \end{aligned}$$

since  $N(\epsilon^n) = 1$  and  $N$  is multiplicative.

We now have

LEMMA 5.2

$$N\left(\frac{\alpha}{1-\epsilon^n} - T^i(\gamma)\right) = N\left(\frac{\alpha}{1-\epsilon^n} - \gamma\right)$$

for every rational integer  $i$ .

5.3 If  $\tau_l$  is an arbitrary real number such that  $\tau_l > \frac{\alpha}{1-\epsilon^n}$ , under the transformation 5.1 the set of integers  $\gamma$  in  $K$  for which  $\gamma \geq \tau_l$  is transformed into that set for which

$$\gamma \geq \frac{\tau_l - \alpha}{\epsilon^n} > \tau_l \quad \text{since } 0 < \epsilon^n < 1 \text{ when } n \geq 1.$$

If  $\tau_u$  is an arbitrary real number such that  $\tau_u < \frac{\alpha}{1-\epsilon^n}$ , the set of integers  $\gamma$  for which  $\gamma \leq \tau_u$  is transformed into that set for which

$$\gamma \leq \frac{\tau_u - \alpha}{\epsilon^n} < \tau_u \quad \text{since } 0 < \epsilon^n < 1 \text{ when } n \geq 1.$$

Now define the sets  $A$  and  $B$ :

$$A = \left\{ \gamma : \gamma \text{ integer in } K \text{ and } \tau_l \leq \gamma < \frac{\tau_l - \alpha}{\epsilon^n} \right\}$$

$$B = \left\{ \gamma : \gamma \text{ integer in } K \text{ and } \frac{\tau_u - \alpha}{\epsilon^n} < \gamma \leq \tau_u \right\}.$$

Suppose  $\gamma$  is an integer in  $K$  which satisfies

$$\frac{\tau_i \alpha}{1 - \epsilon^n} < \gamma < \tau_i$$

and that  $T^i(\gamma)$  does not belong to  $A$  for any rational integer  $i$ , in particular for any positive rational integer  $i$ . We have

$$\gamma > \frac{\tau_i \alpha}{1 - \epsilon^n} \quad \gamma > \frac{\alpha}{1 - \epsilon^n}, \quad T(\gamma) = \frac{\gamma - \alpha}{\epsilon^n}$$

so that  $T(\gamma)$  does not belong to  $A$  hence

$$T(\gamma) - \frac{\alpha}{1 - \epsilon^n} = \frac{\gamma - \frac{\alpha}{1 - \epsilon^n}}{\epsilon^n} > 0$$

since  $0 < \epsilon^n < 1$ . Thus

$$T^i(\gamma) > \frac{\alpha}{1 - \epsilon^n}$$

for every positive rational integer  $i$ . Therefore, we have

Since  $\gamma < \tau_i$ ,  $T(\gamma) < \frac{\tau_i - \alpha}{\epsilon^n}$ , but  $T(\gamma)$  does not belong to  $A$  hence

$$T(\gamma) < \tau_i.$$

Thus  $T^i(\gamma) < \tau_i$  for every positive rational integer  $i$ .

We now have arbitrarily large  $i$  since  $\tau_i$  and  $\alpha$  are fixed,

$$0 < \epsilon^n < 1 \quad \frac{\alpha}{1 - \epsilon^n} < T^i(\gamma) < \tau_i \epsilon^n = \frac{\tau_i (1 - \epsilon^n)}{\epsilon^n (1 - \epsilon^n)} > \tau_i.$$

for every positive rational integer  $i$ , so that

$$\frac{\alpha}{1 - \epsilon^n} < \frac{\gamma - \alpha \sum_{j=0}^{i-1} \epsilon^{nj}}{\epsilon^{ni}} < \tau_i$$

for every positive rational integer  $i$ ,

$$\text{hence for some } \tau_i \quad \gamma < \epsilon^{ni} \tau_i + \alpha \sum_{j=0}^{i-1} \epsilon^{nj} \quad T^i(\gamma) \text{ belongs to } A.$$

for every positive rational integer  $i$ , hence

$$\gamma < e + \frac{\alpha}{1 - \epsilon^n}$$

where  $e$  is arbitrarily small since  $\tau_i$  is fixed and  $0 < \epsilon^n < 1$ ;

but this contradicts 5.3 and 5.4, for any integer  $i$ , in  $A$ .

there is an integer  $\frac{\alpha}{1 - \epsilon^n} < \gamma$  such that

Thus, for some positive rational integer  $i$ ,  $T^i(\gamma)$  belongs to  $A$ .

Thus, Now suppose that  $\gamma$  is an integer in  $K$  which satisfies

$$\gamma > \frac{\tau_i - \alpha}{\epsilon^n}$$

and that  $T^i(\gamma)$  does not belong to  $A$  for any rational integer  $i$ , in particular for any negative rational integer  $i$ .

$$\gamma > \frac{\tau_i - \alpha}{\epsilon^n} \quad \text{hence} \quad T^{-1}(\gamma) > \tau_i.$$

$T^{-1}(\gamma)$  does not belong to  $A$  hence

$$T^{-1}(\gamma) > \frac{\tau_i - \alpha}{\epsilon^n},$$

consequently

$$T^{-i}(\gamma) > \frac{\tau_i - \alpha}{\epsilon^n}$$

for every positive rational integer  $i$ . Therefore, we have

$$\gamma \epsilon^{in} + \sum_{j=0}^{i-1} \alpha \epsilon^{jn} > \frac{\tau_i - \alpha}{\epsilon^n}$$

for every positive rational integer  $i$ , thus

$$\gamma > \epsilon^{-in} \left\{ \frac{\tau_i - \alpha}{\epsilon^n} - \sum_{j=0}^{i-1} \alpha \epsilon^{jn} \right\}$$

for every positive rational integer  $i$  since  $0 < \epsilon^n < 1$ .

Hence  $\gamma > M$  for arbitrarily large  $M$  since  $\tau_i$  and  $\alpha$  are fixed,

$$0 < \epsilon^n < 1 \quad \text{and} \quad \frac{\tau_i - \alpha}{\epsilon^n} - \frac{\alpha}{1 - \epsilon^n} = \frac{\tau_i(1 - \epsilon^n) - \alpha}{\epsilon^n(1 - \epsilon^n)} > 0.$$

This contradicts any particular choice of  $\gamma$ ; hence, if  $\gamma$  is an integer such that

$$\gamma > \frac{\alpha}{1 - \epsilon^n} \tag{5.3}$$

then for some rational integer  $i$ ,  $T^i(\gamma)$  belongs to  $A$ .

Similarly, if  $\gamma$  is an integer such that

$$\gamma < \frac{\alpha}{1 - \epsilon^n} \tag{5.4}$$

then for some rational integer  $i$ ,  $T^i(\gamma)$  belongs to  $B$ .

From lemma 5.2, 5.3 and 5.4, for any integer  $\gamma_1$  in  $K$  there is an integer  $\gamma_2$  in  $A \cup B$  such that

$$N\left(\frac{\alpha}{1 - \epsilon^n} - \gamma_1\right) = N\left(\frac{\alpha}{1 - \epsilon^n} - \gamma_2\right).$$

Thus, if  $C_1$  is any positive real number and there is an integer  $\delta_1$  in  $K$  such that

$$|N(\frac{\alpha}{1 - \epsilon^n} - \delta_1)| < C_1, \quad 5.5$$

there is an integer  $\delta_2$  in  $A \cup B$  such that 5.5 holds with  $\delta_1$  replaced by  $\delta_2$ . In particular if

$$C = \min |N(\frac{\alpha}{1 - \epsilon^n} - \delta)|$$

where the minimum is over integers  $\delta$  in  $K$ , we have

$$C = \min |N(\frac{\alpha}{1 - \epsilon^n} - \delta)|$$

where the minimum is over integers  $\delta$  in  $A \cup B$ .

5.4 Following the notation described in the introduction, the number  $\delta = p_r \lambda + q_r \theta + r_r$  may be represented in the space  $Q$  by the point  $\tilde{\delta} = (\xi_r, \eta_r, \zeta_r)$  where

$$\zeta_r = p_r \lambda + q_r \theta + r_r$$

$$\xi_r = p_r \operatorname{Re} \psi + q_r \operatorname{Re} \phi + r_r$$

$$\eta_r = p_r \operatorname{Im} \psi + q_r \operatorname{Im} \phi.$$

If

5.5 Consider,  $\Delta = \begin{vmatrix} \lambda & \theta & 1 \\ \operatorname{Re} \psi & \operatorname{Re} \phi & 1 \\ \operatorname{Im} \psi & \operatorname{Im} \phi & 0 \end{vmatrix}$  which satisfy 5.7. If

$$|(\zeta_r - \zeta_{r'}) \{ (\xi_r - \xi_{r'})^2 + (\eta_r - \eta_{r'})^2 \}| \leq C_1,$$

since

$$\zeta_r = \begin{vmatrix} \lambda & \theta & 1 \\ \operatorname{Re} \psi & \operatorname{Re} \phi & 1 \\ \operatorname{Im} \psi & \operatorname{Im} \phi & 0 \end{vmatrix}$$

we have

$$p_r = \frac{(\zeta_r - \zeta_{r'}) \begin{vmatrix} \lambda & \theta & 1 \\ \operatorname{Re} \psi & \operatorname{Re} \phi & 1 \\ \operatorname{Im} \psi & \operatorname{Im} \phi & 0 \end{vmatrix}}{\Delta} = \dots$$

Thus, in searching for  $\tilde{\delta}$  satisfying 5.5 and 5.7, it is sufficient to consider those  $\tilde{\delta}$  which also satisfy

$$\zeta_{r'} - \sigma_1 \leq \zeta_r \leq \zeta_{r'} + \sigma_1,$$

$$\eta_{r'} - \sigma_1 \leq \eta_r \leq \eta_{r'} + \sigma_1.$$

From the bound  $q_r = \begin{vmatrix} \lambda_r & \xi_r & 1 \\ \operatorname{Re} \psi & \xi_r & 1 \\ \operatorname{Im} \psi & \eta_r & 0 \end{vmatrix} / \Delta$  we calculate  $p_r^{(1)}, p_r^{(2)}$ ,  $q_r^{(1)}, q_r^{(2)}$  and  $r_r^{(1)}, r_r^{(2)}$  to find those  $\tilde{\gamma}$  satisfying 5.6 and 5.7, it is sufficient to consider those  $\tilde{\gamma}$  for which

$$r_r^{(1)} = \begin{vmatrix} \lambda & \theta & \xi_r \\ \operatorname{Re} \psi & \operatorname{Re} \phi & \xi_r \\ \operatorname{Im} \psi & \operatorname{Im} \phi & \eta_r \end{vmatrix} / \Delta \quad (5.8)$$

$$\quad \quad \quad (5.9)$$

$$\quad \quad \quad (5.11)$$

$\frac{\alpha}{1 - \epsilon^n} = \xi_{\alpha^*}$  may be represented by the point  $\alpha$ , since

$$\tilde{\xi}_{\alpha^*} = (\xi_{\alpha^*}, \eta_{\alpha^*}, \xi_{\alpha^*})$$

for. In searching for the points  $\tilde{\gamma}$  for which satisfied 5.6

$$\text{and 5.7 we have } |N(\frac{\alpha}{1 - \epsilon^n} - \tilde{\gamma})| \leq C_1 \quad (5.6)$$

for some chosen positive real number  $C_1$ , if  $\tau_l, \tau_u$  are arbitrarily chosen real numbers such that  $\operatorname{Im} \phi$  is positive,

$$\tau_u < \frac{\alpha}{1 - \epsilon^n} < \tau_l,$$

from the above it is sufficient to consider those  $\tilde{\gamma}$  for which

$$\tau_l \leq \xi_r < \frac{\tau_l - \alpha}{\epsilon^n} \quad (5.7)$$

$$\text{and } \tau_u \geq \xi_r > \frac{\tau_u - \alpha}{\epsilon^n}. \quad (5.8)$$

In this way we obtain bounds on the values of  $p_r, q_r$  and  $r_r$ .

5.5 Consider, first of all, those  $\tilde{\gamma}$  which satisfy 5.7. If

and so have a finite number of possibilities for  $\tilde{\gamma}$  satisfying

$$|(q_r - \xi_{\alpha^*}) \{ (\xi_r - \xi_{\alpha^*})^2 + (\eta_r - \eta_{\alpha^*})^2 \}| \leq C_1,$$

since  $\xi_r \geq \tau_l > \xi_{\alpha^*}$  (1) then, for any  $\tilde{\gamma}$  with  $q_r$ ,

we have

$$(\xi_r - \xi_{\alpha^*})^2 + (\eta_r - \eta_{\alpha^*})^2 \leq \frac{C_1}{\tau_l - \xi_{\alpha^*}} = \sigma_1^2 \quad \text{say.} \quad (5.10)$$

Thus, in searching for those  $\tilde{\gamma}$  satisfying 5.6 and 5.7, it is

sufficient to consider those  $\tilde{\gamma}$  which also satisfy

$$\xi_{\alpha^*} - \sigma_1 \leq \xi_r \leq \xi_{\alpha^*} + \sigma_1,$$

$$\eta_{\alpha^*} - \sigma_1 \leq \eta_r \leq \eta_{\alpha^*} + \sigma_1.$$



From the bounds on  $\xi_r, \xi_r, \eta_r$ , we may calculate  $p_r^{(1)}, p_r^{(2)}, q_r^{(1)}, q_r^{(2)}, r_r^{(1)}, r_r^{(2)}$  such that, to find those  $\tilde{\gamma}$  satisfying 5.6 and 5.7, it is sufficient to consider those  $\tilde{\gamma}$  for which

$$p_r^{(1)} \leq p_r \leq p_r^{(2)} \tag{5.9}$$

$$q_r^{(1)} \leq q_r \leq q_r^{(2)} \tag{5.10}$$

$$r_r^{(1)} \leq r_r \leq r_r^{(2)}. \tag{5.11}$$

We now suppose  $p_r^{(2)} - p_r^{(1)} \leq q_r^{(2)} - q_r^{(1)}$ , then, since

$$\eta_r = p_r \operatorname{Im} \psi + q_r \operatorname{Im} \phi,$$

for any  $\tilde{\gamma}$  with  $p_r$  satisfying 5.9, which also satisfies 5.6 and 5.7 we have

$$\frac{\eta_{\alpha^*} - \sigma_1 - p_r \operatorname{Im} \psi}{\operatorname{Im} \phi} \leq q_r \leq \frac{\eta_{\alpha^*} + \sigma_1 - p_r \operatorname{Im} \psi}{\operatorname{Im} \phi} \tag{5.12}$$

since  $\phi$  may be chosen such that  $\operatorname{Im} \phi$  is positive.

$$\xi_r = p_r \operatorname{Re} \psi + q_r \operatorname{Re} \phi + r_r$$

thus, for any  $\tilde{\gamma}$  with  $p_r$  satisfying 5.9 and  $q_r$  satisfying 5.12, for which 5.6 and 5.7 hold

$$\xi_{\alpha^*} - \sigma_1 - p_r \operatorname{Re} \psi - q_r \operatorname{Re} \phi \leq r_r \leq \xi_{\alpha^*} + \sigma_1 - p_r \operatorname{Re} \psi - q_r \operatorname{Re} \phi.$$

In this way we obtain bounds on the values of  $p_r, q_r$  and  $r_r$ ; and so have a finite number of possibilities for  $\tilde{\gamma}$  satisfying 5.6 and 5.7.

If  $p_r^{(2)} - p_r^{(1)} > q_r^{(2)} - q_r^{(1)}$  then, for any  $\tilde{\gamma}$  with  $q_r$  satisfying 5.10 which also satisfies 5.6 and 5.7 we have

$$\frac{\eta_{\alpha^*} - \sigma_1 - q_r \operatorname{Im} \phi}{\operatorname{Im} \psi} \leq p_r \leq \frac{\eta_{\alpha^*} + \sigma_1 - q_r \operatorname{Im} \phi}{\operatorname{Im} \psi} \tag{5.13}$$

since  $\operatorname{Im} \psi > 0$  as  $\operatorname{Im} \psi = \frac{2\operatorname{Re} \phi \operatorname{Im} \phi + \theta \operatorname{Im} \phi + a}{2}$ ,  $\phi$  is chosen such

that  $\operatorname{Im} \phi$  is positive,  $\theta + 2\operatorname{Re} \phi = a$  hence  $\operatorname{Re} \phi = \frac{a - \theta}{2} > 0$

since  $a \geq 1$  and  $0 < \theta < 1$ .

We now have, for any  $\tilde{\gamma}$  with  $q_r$  satisfying 5.10 and  $p_r$

satisfying 5.13, turn, those  $\tilde{\gamma}$  satisfying

$$\xi_{\alpha^*} - \epsilon_1 - p_1 \operatorname{Re} \psi - q_1 \operatorname{Re} \phi \leq r_1 \leq \xi_{\alpha^*} + \epsilon_1 - p_1 \operatorname{Re} \psi - q_1 \operatorname{Re} \phi.$$

Again we have bounds on the possible values of  $p_1, q_1$  and  $r_1$ , and so a finite number of possibilities for  $\tilde{\gamma}$  satisfying 5.6 and 5.7.

5.6 Now we consider those  $\tilde{\gamma}$  satisfying 5.8

$$\tau_u \geq \xi_1 > \frac{\tau_u - \alpha}{\epsilon^n} \quad \text{hence} \quad \xi_{\alpha^*} > \tau_u \geq \xi_1$$

and if

$$|(\xi_1 - \xi_{\alpha^*})\{(\xi_1 - \xi_{\alpha^*})^2 + (\eta_1 - \eta_{\alpha^*})^2\}| \leq C_1$$

we have a loss of accuracy in the double precision arithmetic.

$$(\xi_1 - \xi_{\alpha^*})^2 + (\eta_1 - \eta_{\alpha^*})^2 \leq \frac{C_1}{\xi_{\alpha^*} - \tau_u} = C_2^2 \quad \text{say.}$$

Thus in the same way as that used for  $\tilde{\gamma}$  satisfying 5.6 and 5.7 with  $\epsilon_1$  replaced by  $\epsilon_2$ , we may find a finite number of possibilities for  $\tilde{\gamma}$  satisfying 5.6 and 5.8.

5.7 In some cases the range of values of  $\gamma$  in A and B

may be very large, we, therefore, choose integers  $\tau_1^{(1)}, \dots, \tau_1^{(k)}$  such that

$$\frac{\alpha}{1 - \epsilon^n} < \tau_1 = \tau_1^{(1)} < \dots < \tau_1^{(m)} = \frac{\tau_1 - \alpha}{\epsilon^n}$$

and integers  $\tau_u^{(1)}, \dots, \tau_u^{(m)}$  such that

$$\frac{\tau_u - \alpha}{\epsilon^n} = \tau_u^{(m)} < \dots < \tau_u^{(1)} = \tau_u < \frac{\alpha}{1 - \epsilon^n}.$$

Then, using the method outlined above for  $\tilde{\gamma}$  satisfying 5.7, we consider, in turn, those  $\tilde{\gamma}$  satisfying

$$\tau_1^{(i)} \leq \xi_1 < \tau_1^{(i+1)} \quad \text{for } i = 1, \dots, k-1$$

replacing  $\tau_1$  by  $\tau_1^{(i)}$  throughout.

We then use the method outlined above for  $\tilde{\gamma}$  satisfying 5.8

to consider, in turn, those  $\tilde{\gamma}$  satisfying

$$\tau_u^{(i+1)} < \tilde{\gamma} \leq \tau_u^{(i)} \quad \text{for } i = 1, \dots, m-1$$

replacing  $\tau_u$  by  $\tau_u^{(i)}$  throughout.

5.1 This chapter gives an example of the application of each

5.7 When determining  $|N(\frac{\alpha}{1-\epsilon^n} - \gamma)|$  for  $\gamma \in A \cup B$ , we first determine the integer  $\alpha - \gamma(1 - \epsilon^n)$  then calculate its norm and divide this by  $N(1 - \epsilon^n)$ . The integers involved in these calculations often become very large; consequently double

precision arithmetic is used, and a message is monitored when each of the methods described involves the determination the quantities involved are likely to cause integer overflow, of the real zero  $\theta$  of the defining polynomial of the field, that is a loss of accuracy in the double precision entities. a basis for the field, and a basis for the lattice of integers.

If a field has defining polynomial  $x^3 - ax^2 + bx - c$ , the zeros are the roots of the equation  $x^3 - ax^2 + bx - c = 0$ ; we make the substitution  $y = 3x - a$ , that is  $x = \frac{y+a}{3}$ , then

$$y^3 - qy - u = 0 \tag{5.1}$$

where  $q = 3a^2 - 9b$ ,  $u = 27c - 9ab + 2a^3$ ; if we now let  $q = 3uv$ ,  $u = u^3 + v^3$  then  $\theta = u + v$  is the real zero of the polynomial 5.1. This theory is used to determine  $\theta$  and  $\theta = \frac{\theta + a}{3}$ , the former is required for the application of Voronoi's algorithm which is initially referred to in chapter 1.

The cubic field  $K$  with discriminant  $-108$  has defining polynomial  $x^3 - 2x^2 + 5x - 3$  thus

$$u = 7, \quad v = -33$$

and we obtain

$$\theta = .73727772\text{e}+00$$

$$\theta = .21193316\text{e}+00$$

A NUMERICAL CONSIDERATION OF THE METHODS EMPLOYED.

For the lattice of integers

6.1 This chapter gives an example of the application of each of the methods of this thesis, using the cubic field with discriminant -199 as example. Also included in this chapter are descriptions of some of the basic routines employed and observations on the limitations of the methods.

6.2 Each of the methods described involves the determination of the real zero  $\theta$  of the defining polynomial of the field, a basis for the field, and a basis for the lattice of integers.

If a field has defining polynomial  $x^3 - ax^2 + bx - c$ , the zeroes are the roots of the equation  $x^3 - ax^2 + bx - c = 0$ ; we make the substitution  $y = 3x - a$ , that is  $x = \frac{y+a}{3}$ , then

$$y^3 - qy - n = 0 \tag{6.1}$$

where  $q = 3a^2 - 9b$ ,  $n = 27c - 9ab + 2a^3$ ; if we now let  $q = 3uv$ ,  $n = u^3 + v^3$  then  $\omega = u + v$  is the real zero of the polynomial 6.1. This theory is used to determine  $\omega$  and  $\theta = \frac{\omega + a}{3}$ , the former is required for the application of Voronoi's algorithm which is initially referred to in chapter 1.

The cubic field  $K$  with discriminant -199 has defining polynomial  $x^3 - 2x^2 + 5x - 3$  thus  $n = 7$ ,  $q = -33$

and we obtain

$$\theta = .73727772.....$$

$$\omega = .21183316.....$$

The index of the polynomial relative to the field is 1, thus the basis is  $(1, \theta, \lambda)$  where  $\lambda = \theta^2$ . We now have the basis for the lattice of integers  $(0,0,0), (1,0,1), (.63136113\dots, 1.9158303\dots, .73727772\dots), (-3.2717892\dots, 2.4191617\dots, .54357843\dots)$ .

6.3 Most of the methods employed require the determination of the coefficients of  $1, \theta, \lambda$  in the expressions for  $\theta^2, \lambda^2$  and  $\theta\lambda$ ; these values are subsequently used in routines for the product and the quotient of two algebraic numbers. For the field  $K$  we find, since  $\lambda = \theta^2$ ,

$$\theta^2 = -1\lambda + 0\theta + 0$$

$$\lambda^2 = -1\lambda - 7\theta + 6$$

$$\theta\lambda = 2\lambda - 5\theta + 3$$

Some of the methods require the determination of the norm of an algebraic number, the product of two algebraic numbers or the quotient of two algebraic numbers. It is in the routines for these operations that there is the greatest danger of integer overflow; to prevent undetected overflow, the coefficients with respect to the basis of the field of the algebraic integer or integers involved are tested to ensure that the coefficients of the resultant integer will not require more significant figures than possible with the computer in order to give an exact result.

When the product of  $n$  coefficients is involved, for single precision arithmetic we ensure that no product may exceed  $10^{14-t}$ , and for double precision arithmetic we ensure that no product may



exceed  $10^{29-t}$ , where, in each case,  $t$  is the smallest power of 10 which exceeds the greatest constant multiplier of the products of coefficients involved in the particular subroutine in question.

6.4 I now turn to the adaptation of Cassels' result described in chapter 2. We start with the ideal

$$(1, \theta, \theta^2) = (1, \frac{6 + 3\omega}{9}, \frac{4 + 4\omega + \omega^2}{9})$$

which has reduced form itself, and multiples of it by the fundamental

$$\text{unit. However, } (1, \frac{3 - 3\omega}{9}, \frac{7 + \omega + \omega^2}{9}),$$

and the next ideal in the loop is also this ideal. The reduced

ideal, expressed in terms of  $\theta$ , is given by

$$(1, 1 - \theta, 1 - \theta + \theta^2),$$

and we see that one complete sequence of relative minima is

$$1, 1 - \theta$$

and  $\epsilon = 1 - \theta$ . We then have

The method  $M = 3.8063007$  thus  $\frac{|D|}{720M} = .0726$  ;

in fact  $\frac{|D|}{715.244M} = .0731$   $\frac{|\Delta|}{412.944} = .0342$ .

Using the method for fields in the table of (1) no results concerning the existence of the Euclidean algorithm could be found.

The field with discriminant  $-160087$  was shown to have 44

distinct relative minima, thus the chain calculated consisted of

$$\left[ \frac{44 + 1}{3} \right] = 15 \text{ relative minima "before" } \theta^2 \text{ and } 29 \text{ "after" } \theta^2$$

The resultant value of  $M$  was 188.308... so that

$$\frac{|D|}{720M} = 1.1807 > 1,$$

hence the field does not possess a Euclidean Algorithm.

Similarly, for the field with discriminant  $-169571$   
 $M = 167.136\dots$  thus

$$\frac{|D|}{720M} = 1.4091 > 1$$

and this field does not possess a Euclidean Algorithm.

6.5 For the method described in chapter 3 we have  $\epsilon = 1 - \theta$

hence  $\epsilon - 1 = -\theta$  and  $N(\epsilon - 1) = -3$ . Thus the only divisor of  $\epsilon - 1$  is  $\epsilon - 1$  itself, and multiples of it by the fundamental unit. However, any integer of norm 2 or 3 will have the fundamental unit of the field, or its negative, as a member of its non-zero residue classes; consequently, no result can be obtained by considering residue classes modulo  $\epsilon - 1$ .

$\epsilon + 1 = -\theta + 2$  has norm 7, thus any factor, not a power of the fundamental unit, must be a multiple of  $\epsilon + 1$  by a power of the fundamental unit.

The method used for calculating integers of a given norm causes congruences to be considered modulo  $\beta = \theta^2 - \theta + 5 = (\epsilon + 1)/\epsilon$ .

We then have  $\epsilon \equiv -1 (\beta)$  thus  $-\epsilon \equiv 1 (\beta)$ . There are no integers of norm 2 and so none of norms 4 or 6.

$3 = \theta(\theta^2 - 2\theta + 5)$  where  $\theta^2 - 2\theta + 5$  is an algebraic prime of norm 9. Thus integers of norm 3 are either congruent to  $\theta \equiv 2 (\beta)$  or to  $-\theta \equiv 5 (\beta)$ . There are no integers of norm 5.

Thus we see that there are no integers of norm less than  $|N(\beta)|$  congruent to 3 or 4 modulo  $\beta$ , and so the field does not possess a Euclidean Algorithm.

When using this method it was found that, as the size of the coefficients of the fundamental unit with respect to the field increased, the magnitude of the norms of  $\epsilon - 1$  and of  $\epsilon + 1$  also increased and had prime factors of norm greater than 500, with the result that no conclusion was reached; also, for such fields, overflow warning messages were given, both in calculating the integers of a given norm and in establishing to which congruence class any particular integer belonged, thus invalidating any results which may have been obtained. With these points in mind only those fields with discriminant of absolute value less than 5,000 were investigated by means of the program, but some additional ones were investigated "by hand".

Generally, when  $\alpha$  is an integer of norm less than  $|N(\beta)|$ , we determine to which congruence class it belongs by establishing for which  $\zeta_j$ , where  $1 \leq j \leq r_\beta$ ,  $\beta$  divides  $\zeta_j - \alpha$ . However, this method becomes time consuming when  $|N(\beta)|$  is large; we note that when  $|N(\beta)|$  is a rational prime the residue classes modulo  $\beta$  are represented by the rational integers  $0, 1, \dots, |N(\beta)| - 1$  and we have  $\theta \equiv i_\theta (\beta)$  and  $\lambda \equiv i_\lambda$  for some rational integers  $i_\theta$  and  $i_\lambda$  satisfying  $0 \leq i_\theta \leq |N(\beta)| - 1$  and  $0 \leq i_\lambda \leq |N(\beta)| - 1$ ; hence, for any  $\alpha = p\lambda + q\theta + r$  we have  $\alpha \equiv pi_\lambda + qi_\theta + r (\beta)$ . By reducing the expression modulo  $|N(\beta)|$  we establish to which residue class  $\alpha$  belongs. In some cases this modification reduced the time taken to establish a result from over 120 seconds of central processor time to less than 5 seconds.

6.6 We now consider the adaptation of the method of Barnes and Swinnerton-Dyer as described in chapter 4. Starting from the basis for  $K$  and the corresponding basis for  $\mathcal{A}$  described above, and following the arguments of 4.8, we obtain the basis

$(1, \theta - 1, \theta^2 - \theta + 2)$  for  $K$  then the corresponding basis for  $\mathcal{A}$  is

$(0,0,0), (1,0,1), (-.3686388604, 1.9158303963, -.2627222793),$   
 $(-1.9031503584, .5033313285, 1.8063007167).$

This results in  $\mathcal{C}$  being the set of points  $(\xi, \eta, \zeta)$  given by

$$|\xi| \leq 1.6358946094$$

$$|\eta| \leq 1.2095808624$$

$$|\zeta| \leq 1.5345114980.$$

In the following a region which satisfies 4.2 will be said to be covered.

On subdivision of  $\mathcal{C}$  into sub-regions, with  $C = .9999$ , there are 14 uncovered regions with  $\zeta \geq 0$ ; but only one of these regions contains uncovered sub-regions itself, namely  $R_1$ , that given by

$$-1.3087156875 \leq \xi \leq -.9815367656$$

$$.9676646899 \leq \eta \leq 1.2095808624$$

$$.3069022996 \leq \zeta \leq .6133045992.$$

$R_1$  contains five uncovered sub-regions but only two of these contain uncovered sub-regions themselves, namely

$R_{11}$  :

$$-1.2759977953 \leq \xi \leq -1.2432799031$$

$$1.0160479244 \leq \eta \leq 1.0402395416$$

$$.3682827595 \leq \zeta \leq .3989729895$$



$$R_{12} : \dots \in (\mathbb{Z} - (R \cup R_1)) \pmod{1},$$

$$\text{Thus any point } -1.2759977953 \leq \xi \leq -1.2432799031$$

$$1.0402395416 \leq \eta \leq 1.0644311589$$

$$\text{or } .3682827595 \leq \zeta \leq .3989729895.$$

On further subdivision,  $R_{11}$  contains three uncovered sub-regions, two of which themselves contain uncovered sub-regions, namely

$$R_{111} : \dots = a(2\theta - 1) + 2\theta - 1 = \beta,$$

$$\text{where } -1.2629106384 \leq \xi \leq -1.2596388492$$

$$\text{and we have } 1.0354012182 \leq \eta \leq 1.0378203799$$

$$(2^2 - 2\theta + 1)\beta \quad .3744208055 \leq \zeta \leq .3774898285$$

$$R_{112} : \dots = \beta = \frac{-2\theta^2 + \theta}{\theta^2 - 1} = \frac{-2\theta^2 + \theta}{\theta^2 - 1}.$$

$$\text{It now remains } -1.2596388492 \leq \xi \leq -1.2563670600$$

$$1.0354012182 \leq \eta \leq 1.0378203799$$

$$\text{where the minimum } .3744208055 \leq \zeta \leq .3774898285 \text{ will be done}$$

$R_{12}$  contains four uncovered sub-regions, but none of these contain uncovered sub-regions themselves.  $\tilde{\alpha}$  is not equal to  $\beta$

$R_{111}$  contains two and  $R_{112}$  five uncovered sub-regions

which are all contained in the region  $R$  given by

6.7 When using the method of chapter 4, it was found that as

$$-1.2599660281 \leq \xi \leq -1.2589844914$$

the absolute value of the discriminant of the field increased,

$$1.0366107991 \leq \eta \leq 1.0378203800$$

the number of uncovered sub-regions of a particular region became

$$.3756484147 \leq \zeta \leq .3759553170,$$

large with  $C = 15000$ , the "critical" value in the investigation thus any point  $\tilde{\alpha}$  in  $\mathcal{J}$  with  $|N(\tilde{\alpha} - \tilde{\gamma})| > .9999$ , where  $\tilde{\gamma}$  is of the Euclidean Algorithm. To avoid very large amounts of an integer in  $\mathcal{J}$ , must be in  $R$  or in  $R_-$ , where  $R_-$  is the output, when the number of uncovered sub-regions of a particular set of those points  $\tilde{\alpha}$  for which  $-\tilde{\alpha}$  is in  $R$ .

region avoided a bound given as data at run-time, only the

We find that any point  $\tilde{\alpha}$  in  $R$  satisfies either

number of sub-regions, and not their nature, was printed; when

$$E(\tilde{\alpha}) - \tilde{\gamma} \text{ is in } R_- \text{ where } \gamma = (1 + 2(\theta - 1))$$



or nature of the  $E(\tilde{\alpha}) \in (\mathfrak{F} - (R \cup R_-)) \pmod{1}$ . Thus any point  $\tilde{\alpha}$  in  $R$  satisfies either  $E^2(\tilde{\alpha}) - E(\tilde{\gamma}) + \tilde{\gamma}$  is in  $R$  suboid, with sides or  $E(\tilde{\alpha}) - \tilde{\gamma} \in (\mathfrak{F} - (R \cup R_-)) \pmod{1}$ . In the latter case  $\tilde{\alpha}$  must have minimum less than .9999.

The required fixed point  $\tilde{\beta}$  of  $E^2$  is given by the number of roots  $\epsilon^2\beta - \epsilon(2\theta - 1) + 2\theta - 1 = \beta$ , where  $\epsilon = -\theta + 1$ ,  $\epsilon^2 = \theta^2 - 2\theta + 1$  to make the method and we have

$$(\theta^2 - 2\theta + 1)\beta - (-2\theta^2 + 3\theta - 1) + 2\theta - 1 = \beta$$

$$\beta = \frac{-2\theta^2 + \theta}{\theta^2 - 2\theta} = \frac{-2\theta^2 + \theta}{\epsilon^2 - 1}.$$

It now remains to calculate  $\min_{\delta} |N(\tilde{\beta} - \tilde{\delta})|$ , where the minimum is over integers  $\delta$  in  $K$ ; this will be done using the method described in chapter 5. From above, we know that for every  $\tilde{\alpha}$  in  $\mathfrak{F}$  such that  $\tilde{\alpha}$  is not equal to  $\tilde{\beta}$

$$\min_{\delta} |N(\tilde{\alpha} - \tilde{\delta})| < .9999.$$

6.7 When using the method of chapter 4, it was found that as the absolute value of the discriminant of the field increased, the number of uncovered sub-regions of a particular region became large with  $C = .9999$ , the "crucial" value in the investigation of the Euclidean Algorithm. To avoid very large amounts of output, when the number of uncovered sub-regions of a particular region exceeded a bound given as data at run-time, only the number of sub-regions, and not their nature, was printed; when

the nature of the sub-regions was required, a modified version of the program was used, this determined and printed the bounds on the co-ordinates of points in the smallest cuboid, with sides parallel to the co-ordinate axes, which contained all of the uncovered sub-regions of a given region.

However, even with the modified program, it was found that the number of regions to be considered, when attempting to satisfy II(a) or II(b) of chapter 4, became so large as to make the method require a large amount of computing for those fields of discriminants -680 and -687, and completely impractical for that with discriminant -1004. Consequently, with the exception of the field with discriminant -1004, no investigation was made of fields after that with discriminant -687 in the table of (1).

6.8 I now turn to the determination of the minimum of

$$\frac{-2\theta^2 + \theta}{e^2 - 1} = \frac{2\theta^2 - \theta}{1 - e^2} \text{ using the method of chapter 5.}$$

Note that

$$\frac{2\theta^2 - \theta}{1 - e^2} \equiv \frac{3\theta}{1 - e^2} \pmod{1}$$

and

$$\frac{3\theta}{1 - e^2} = \frac{3}{-2 + 2} = \frac{3\theta^2 + 15}{7} = \frac{3.5435784374 + 15}{7} = 2.375819331 = \xi_{\alpha^*}$$

$\alpha^* = \frac{3\theta}{1 - e^2} = \frac{3\theta^2 + 15}{7}$  is fixed under  $e^2$ , and  $\alpha^* > 0$ ;  $\alpha^*$  is used here instead of  $\frac{2\theta^2 - \theta}{1 - e^2}$  to simplify calculations.

$$\tilde{\alpha}^* \text{ is the point } \left\{ \operatorname{Re} \left\{ \frac{3\phi^2 + 15}{7} \right\}, \operatorname{Im} \left\{ \frac{3\phi^2 + 15}{7} \right\}, \frac{3\theta^2 + 15}{7} \right\}$$

$$\operatorname{Re} \left\{ \frac{3\phi^2 + 15}{7} \right\} = \frac{3\operatorname{Re}\phi^2 + 15}{7} = .740661764 = \xi_{\alpha^*}$$

$$\operatorname{Im} \left\{ \frac{3\phi^2 + 15}{7} \right\} = \frac{3\operatorname{Im}\phi^2}{7} = 1.036783597 = \eta_{\alpha^*}$$

that is,  $\tilde{\alpha}^*$  is the point  $\{ .740661764, 1.036783597, 2.375819331 \}$   
 $= \{ \xi_{\alpha^*}, \eta_{\alpha^*}, \zeta_{\alpha^*} \}.$

If  $\tau_l = 5$  and  $\tau_u = -5$  we have  $\tau_u < \zeta_{\alpha^*} < \tau_l$ . We wish to consider those points  $\tilde{\gamma} = (\xi_r, \eta_r, \zeta_r)$  for which

$$|N(\tilde{\alpha}^* - \tilde{\gamma})| \leq 2$$

where 2 is chosen as a value greater than 1. We note that if

$$\begin{aligned} \xi_r &= p_r \operatorname{Re} \phi^2 + q_r \operatorname{Re} \phi + r_r \\ \eta_r &= p_r \operatorname{Im} \phi^2 + q_r \operatorname{Im} \phi \\ \zeta_r &= p_r \theta^2 + q_r \theta + r_r \end{aligned} \tag{6.2}$$

we have

$$\begin{aligned} p_r &= -.271619232\dots \xi_r - .015016455\dots \eta_r + .271619232\dots \zeta_r \\ q_r &= .342979656\dots \xi_r + .540928484\dots \eta_r - .342979656\dots \zeta_r \\ r_r &= -.105224901\dots \xi_r - .390651899\dots \eta_r + 1.105224901\dots \zeta_r \end{aligned}$$

check: the determinant of the coefficients is  $-.14177624 = -\sqrt{\frac{4}{199}}$  and  $-\sqrt{\frac{199}{4}}$  is the determinant of the coefficients of the equations 6.2.

Following the reasoning of chapter 5, we first consider those

points  $\tilde{\gamma}$  for which

$$\tau_l \leq \zeta_r \leq \frac{\tau_l - 3\theta}{\epsilon^2};$$

using the notation of chapter 5

$$\sigma_1 = \left(\frac{2}{5 - \zeta_{\alpha^*}}\right)^{\frac{1}{2}} = .873007817$$

and we have

$$\begin{aligned} \zeta_{\alpha^*} - \sigma_1 &= -.132346053 \leq \zeta_r \leq \zeta_{\alpha^*} + \sigma_1 = 1.613669581 \\ \eta_{\alpha^*} - \sigma_1 &= .163775780 \leq \eta_r \leq \eta_{\alpha^*} + \sigma_1 = 1.909791414 \\ \tau_l &= 5 \leq \zeta_r \leq \frac{\tau_l - 3\theta}{\epsilon^2} = 40.394753039. \end{aligned}$$

We now obtain the bounds

$$\begin{aligned} .891114170 &\leq p_r \leq 11.005480200 \\ -13.811379521 &\leq q_r \leq -.128381867 \\ 4.610262640 &\leq r_r \leq 44.595233710 \end{aligned}$$

hence

$$1 \leq p_y \leq 11$$

$$-13 \leq q_y \leq -1$$

$$5 \leq r_y \leq 44.$$

For each value of  $p_y$  we may obtain smaller ranges for  $q_y$ , then for  $r_y$ , as follows

$$p_y = 10 \text{ requires } q_y = \frac{r_y - p_y \text{Im}\phi^2}{\text{Im}\phi} = 41 \Rightarrow |N(\alpha^* - \tilde{\gamma})| = 1$$

$$\text{hence } q_y = -10 \quad r_y = 45 \Rightarrow |N(\alpha^* - \tilde{\gamma})| = \frac{309}{21}$$

$$.085485532 - 1.262722280p_y \leq q_y \leq .996847852 - 1.262722280p_y$$

$$r_y = \xi_y - p_y \text{Re}\phi^2 - q_y \text{Re}\phi$$

thus

$$-.132346053 + 3.271789219p_y - .6313611396q_y$$

$$\leq r_y \leq 1.613669581 + 3.271789219p_y - .6313611396q_y.$$

$$p_y = 1 \text{ requires } \left\{ \begin{array}{l} q_y = -1 \quad r_y = 4 \Rightarrow |N(\alpha^* - \tilde{\gamma})| = 1 \\ \text{or } r_y = 5 \Rightarrow |N(\alpha^* - \tilde{\gamma})| = 1 \end{array} \right.$$

$$p_y = 2 \text{ requires } \left\{ \begin{array}{l} q_y = -2 \quad r_y = 8 \Rightarrow |N(\alpha^* - \tilde{\gamma})| = \frac{33}{21} \\ \text{or } r_y = 9 \Rightarrow |N(\alpha^* - \tilde{\gamma})| = \frac{27}{21} \end{array} \right.$$

$$p_y = 3 \text{ requires } \left\{ \begin{array}{l} q_y = -3 \quad r_y = 12 \Rightarrow |N(\alpha^* - \tilde{\gamma})| = \frac{81}{21} \\ \text{or } r_y = 13 \Rightarrow |N(\alpha^* - \tilde{\gamma})| = \frac{111}{21} \end{array} \right.$$

$p_y = 4$  gives no integer values of  $q_y$ .

$$p_y = 5 \text{ requires } q_y = -6 \quad r_y = 21 \Rightarrow |N(\alpha^* - \tilde{\gamma})| = \frac{72}{21}$$

$$p_y = 6 \text{ requires } \left\{ \begin{array}{l} q_y = -7 \quad r_y = 24 \Rightarrow |N(\alpha^* - \tilde{\gamma})| = \frac{261}{21} \\ \text{or } r_y = 25 \Rightarrow |N(\alpha^* - \tilde{\gamma})| = 1 \end{array} \right.$$

$$p_y = 7 \text{ requires } \left\{ \begin{array}{l} q_y = -8 \quad r_y = 28 \Rightarrow |N(\alpha^* - \tilde{\gamma})| = \frac{339}{21} \\ \text{or } r_y = 29 \Rightarrow |N(\alpha^* - \tilde{\gamma})| = \frac{216}{21} \end{array} \right.$$

$$p_y = 8 \text{ requires } q_y = -10 \begin{cases} r_y = 33 \Rightarrow |N(\tilde{\alpha}^* - \tilde{\gamma})| = \frac{441}{21} \\ \text{or } r_y = 34 \Rightarrow |N(\tilde{\alpha}^* - \tilde{\gamma})| = \frac{783}{21} \end{cases}$$

$$p_y = 9 \text{ requires } q_y = -11 \begin{cases} r_y = 37 \Rightarrow |N(\tilde{\alpha}^* - \tilde{\gamma})| = \frac{87}{21} \\ \text{or } r_y = 38 \Rightarrow |N(\tilde{\alpha}^* - \tilde{\gamma})| = \frac{591}{21} \end{cases}$$

$$p_y = 10 \text{ requires } q_y = -12 \quad r_y = 41 \Rightarrow |N(\tilde{\alpha}^* - \tilde{\gamma})| = 1$$

$$p_y = 11 \text{ requires } q_y = -13 \quad r_y = 45 \Rightarrow |N(\tilde{\alpha}^* - \tilde{\gamma})| = \frac{369}{21}$$

We must now consider those points  $\tilde{\gamma}$  for which

$$\tau_u \geq \xi_y \geq \frac{\tau_u - \alpha}{\epsilon^n}$$

and we have

$$\epsilon_2 = \left( \frac{2}{\xi_{u^*} + 5} \right)^{\frac{1}{2}} = .52072672\dots$$

hence

$$\xi_{u^*} - \epsilon_2 \leq \xi_y \leq \xi_{u^*} + \epsilon_2$$

$$\eta_{u^*} - \epsilon_2 \leq \eta_y \leq \eta_{u^*} + \epsilon_2$$

$$\frac{\tau_u - 3\theta}{\epsilon^2} \leq \xi_y \leq \tau_u$$

and we proceed as above to reach eventually the conclusion

$$\min_y |N\left(\frac{3\theta}{1 - \epsilon^2} - \gamma\right)| = 1.$$

We now have the result that the cubic field with discriminant -199 does not possess a Euclidean Algorithm; it has inhomogeneous minimum 1 and this minimum is attained at the numbers congruent to  $\pm \frac{(1 + 3\theta^2)}{7} \pmod{1}$ .

6.9 For the method of chapter 5, as the value of  $n$  increases  $\epsilon^n$  becomes so small that the range of values of  $p_y, q_y, r_y$  to be considered becomes too large to be practical. For the field of discriminant -680 it was found that no result could be obtained with  $n = 3$  within 60 seconds central processor time.



CHAPTER 7

THE PROGRAM BELMIN

This program is that used for the method of chapter 2. In this and other program descriptions the subroutines will be described before the main program; widely used values are passed between the subroutines by means of COMMON variables.

PART II

FUNCTION IOP(1)

The answer returned is the greatest power of 10 which is required when we are determining bounds which are used for the prevention of integer overflow.

SUBROUTINE BASHN

The field has defining polynomial  $x^3 - I_A x^2 + I_B x - I_C$ , discriminant  $DISC$ , and the index of the polynomial over the field is  $INDEX$ . The routine determines the values  $PR = \textcircled{p}$ ,  $IQ = q = q$ ,  $RNE = RN = n$ ,  $RHH = \textcircled{h}$  and  $RK = RHH - q$ , where  $\textcircled{p}$ ,  $q$  and  $n$  are as defined in 5.2. These values are all required in the application of Vercaut's algorithm in the subroutine MINIMA.

The values  $IL1, IL2, IL3, IP1, IP2, IP3$  where

$$\lambda^3 = IL1 \lambda + IL2 \theta + IL3$$

$$\theta^3 = IP1 \lambda + IP2 \theta + IP3$$

$$\theta \lambda = IP1 \lambda + IP2 \theta + IP3$$

are also determined. From these values we now have a basis for the lattice of integers

$$(0, \theta, \theta^2), (R(1), U(1), H(1)), (R(2), U(2), H(2)), (R(3), U(3), H(3))$$

Finally the subroutines in CHAPTER 7 determine the upper bounds on

### THE PROGRAM RELMIN

subroutines MULTCD and DIVID, which are used for double precision

This program is that used for the method of chapter 2. In multiplication and division respectively, this and other program descriptions the subroutines will be described

before the main program; widely used values are passed between the subroutines by means of COMMON variables.

### FUNCTION IOF(I)

The answer returned is the greatest power of 10 which exceeds the value of I. This value is required when we are determining bounds which are used for the prevention of integer overflow.

### SUBROUTINE BASEH

The field has defining polynomial  $x^3 - IA.x^2 + IB.x - IC$ , discriminant IDET, and the index of the polynomial over the field is INDEX. The routine determines the values  $PH = \textcircled{M}$ ,  $IQ = Q = q$ ,  $INN = RN = n$ ,  $HHH = \textcircled{N}^2$  and  $ZZ = HHH - q$ , where  $\textcircled{M}$ ,  $q$  and  $n$  are as defined in 6.2. These values are all required in the

application of Voronoi's algorithm in the subroutine MINIMA. The values IL1, IL2, IL3, IT1, IT2, IT3, IP1, IP2, IP3 where

$$\lambda^2 = IL1.\lambda + IL2.\theta + IL3 \left( \frac{L2.\lambda + L2.\theta + L2}{L1} \right)$$

again the coefficients and denominator are first

$$\theta\lambda = IP1.\lambda + IP2.\theta + IP3$$

are also determined. From these values we now have a basis for the lattice of integers

$$(0,0,0), (R(1),U(1),H(1)), (R(2),U(2),H(2)), (R(3),U(3),H(3))$$

Finally the subroutine is used to determine the upper bounds on the product of the coefficients of the algebraic integers for the subroutines MULTCD and DIVCD, which are used for double precision multiplication and division respectively.

The subroutines SUB(I,J,K,L), INVER(I,J,K,KDET), ICF(IJ,IK,IL,IH), ICF2(IJ,IK,IL,IH) and MULT(J1,J2,J3,K1,K2,K3,L1,L2,L3) are those used by Angell in (1) chapter 6.

SUBROUTINE PHITH(I1,I2,I3,ID,J1,J2,J3,JD)

Given I1, I2, I3, ID, this routine finds J1, J2, J3, JD such that

$$\frac{I1 + I2.\theta + I3.\theta^2}{ID} = \frac{J1.\lambda + J2.\theta + J3.\theta^2}{JD}$$

SUBROUTINE DIVCD(L1,M1,N1,LN,L2,M2,N2,LD,L,M,N,LR)

L, M, N and LR are determined such that

$$\frac{L.\lambda + M.\theta + N}{LR} = \left( \frac{L1.\lambda + M1.\theta + N1}{LN} \right) / \left( \frac{L2.\lambda + M2.\theta + N2}{LD} \right);$$

the coefficients in the numerator and denominator are first tested for the possibility of overflow.

SUBROUTINE MULTCD(L1,M1,N1,LN,L2,M2,N2,LD,L,M,N,LR)

L, M, N and LR are determined such that

$$\frac{L.\lambda + M.\theta + N}{LR} = \left( \frac{L1.\lambda + M1.\theta + N1}{LN} \right) . \left( \frac{L2.\lambda + M2.\theta + N2}{LD} \right);$$

again the coefficients in the numerator and denominator are first tested for the possibility of overflow.

DOUBLE FUNCTION DCI(X,Y,Z,RD)

The answer returned is the co-ordinate of  $X + Y.\theta + Z.\lambda$

in the  $\xi$ ,  $\eta$  or  $\zeta$  direction when RD is the array R, U or H respectively.

**SUBROUTINE MINIMA**

SECTION 1 of this routine closely follows the subroutine

UNIT of (1) chapter 6. (On exit from SECTION 1 we have a loop

of N lattices and  $IOF(I)$  are as for

$$\left( 1, \frac{IAN(1,J) + IAN(2,J) \cdot \Theta + IAN(3,J) \cdot \Theta^2}{NP(J)}, \frac{JAN(1,J) + JAN(2,J) \cdot \Theta + JAN(3,J) \cdot \Theta^2}{NP(J)} \right)$$

for  $J = 1, \dots, N$ , and the first and Nth lattices are the same.

From these lattices we obtain a sequence of relative minima

$$L(1,ILAT) \cdot \lambda + L(2,ILAT) \cdot \theta + L(3,ILAT)$$

for  $ILAT = 1, \dots, N$ ; and if  $NQP = \lfloor N/3 \rfloor + 1$  we have

$$L(1,NQP) \cdot \lambda + L(2,NQP) \cdot \theta + L(3,NQP) = 1.$$

From these relative minima we obtain the value REMN which is

the value M of chapter 2; thus we obtain the value of

$DET/720 \cdot REMN$  where DET is a real variable with the same value

as IDET.

**THE MAIN PROGRAM.**

The main program simply co-ordinates the execution of the subroutines, and reads and prints relevant information about rational prime factors of  $NN$ . Finally, I is the number of the fields.

Following the reasoning of chapter 2, this program only gives a meaningful result when  $INDEX = 1$ .

has at least one factor less than  $n^2$ , unless it is prime.

IDEAL is a simplified version of the subroutine FACTOR  
 THE PROGRAM CONG

of (1) chapter 6; IDEAL finds the ideals of norm N, where N

The method of chapter 3 is put into practice by using this

program. The subroutines ICF2(IJ,IK,IL,IH), SUB(I,J,K,L),

ICF(IJ,IK,IL,IH), INVER(I,J,K,KDET), MULT(J1,J2,J3,K1,K2,K3,L1,L2,L3),

DIVCD(L1,M1,N1,LN,L2,M2,N2,LD,L,M,N,R) and IOF(I) are as for

the program RELMIN. The subroutine MULT2(A,B,C,D,E,F,G,H,I)

is the double precision version of MULT as described in chapter 6

of (1).

SUBROUTINE BASE

This routine is basically similar to BASEH in the program

RELMIN. BASE does not calculate a basis for the lattice of

integers. Additionally, BASE finds upper bounds for the product

of the coefficients of the algebraic integers for the subroutines

MULTC and DIVC.

SUBROUTINE FACTOR(NN,NF,NP,NL,I)

This routine is used to determine the rational prime factors

of NN which are less than 500. NL is chosen by the calling

routine to be an integer greater than the number of distinct

rational prime factors of NN. Finally I is the number of

distinct rational prime factors of NN and these factors are in

the array NF such that NP(I) is the power to which NF(I)

is a factor of NN. In the routine we note that an integer n

has at least one factor less than  $n^{\frac{1}{2}}$ , unless it is prime.



SUBROUTINE IDEAL(N,ISC)

IDEAL is a simplified version of the subroutine FACTOR of (1) chapter 6; IDEAL finds the ideals of norm N, where N is restricted to be a rational prime.

$$\text{NORL}(\text{KB}) = N((\theta^2 + \text{KT} \cdot \theta + \text{KB})/L)$$

where L is the index of the polynomial  $x^3 - \text{IA}x^2 + \text{IB}x - \text{IC}$ ,

and KT, KS are such that  $\lambda = \frac{\theta^2 + \text{KT}\theta + \text{KS}}{L}$ ; and

$$\text{NORN}(\text{J}) = N(\theta + \text{J}).$$

The ideals of norm N are

$$(\text{A}(\text{I}), \text{B}(\text{I}) \cdot \theta + \text{C}(\text{I}), \text{D}(\text{I}) \cdot \lambda + \text{E}(\text{I}) \cdot \theta + \text{G}(\text{I}))$$

for  $\text{I} = 1, \dots, \text{ISC}$ .

SUBROUTINE CHANGE(I,J,K,L,I1,J1,K1)

The integers I, J, K, L are supplied to the routine; I1, J1, K1 are found such that  $\text{I1} \cdot \lambda + \text{J1} \cdot \theta + \text{K1} = (\text{I} \cdot \theta^2 + \text{J} \cdot \theta + \text{K})/L$ .

SUBROUTINE PRIN(INK)

This subroutine, as the subroutine MINIMA of chapter 7, is based on the subroutine UNIT of (1).

When INK is equal to 1, starting from the basis  $(1, \theta, \lambda)$  the routine finds the reduced basis for the unit ideal of the field, which is given by

$$\left( 1, \frac{\text{MU1} + \text{MU2}\theta + \text{MU3}\theta^2}{\text{IGU1}}, \frac{\text{NU1} + \text{NU2}\theta + \text{NU3}\theta^2}{\text{IGU1}} \right).$$

When INK is greater than 1, given the ideal

$(\text{JA}, \text{JB} \cdot \theta + \text{JC}, \text{JD} \cdot \lambda + \text{JE} \cdot \theta + \text{JF})$  of norm INK the routine finds an integer  $\text{JP} \cdot \lambda + \text{JR} \cdot \theta + \text{JS}$  which produces the ideal.

These calculations are based on the theory of chapter 1.

PRIN also includes tests on the arguments of the routines MULT2 and CHANGE to check for the possibility of integer overflow.

DOUBLE FUNCTION FDET(I1,I2,I3,J1,J2,J3,K1,K2,K3)

The answer returned is the value of the determinant

$$\begin{vmatrix} I1 & I2 & I3 \\ J1 & J2 & J3 \\ K1 & K2 & K3 \end{vmatrix}$$

In the program RELMIN this was written as a function statement in the subroutine DIVCD; but the compiler used for this program appeared to loop when attempting to optimise the expression for the determinant, this was overcome by compiling this subroutine alone without full optimisation. The different compiler was used since the one for this program considerably increased the compile-time but greatly reduced the run-time; the latter was important for this program due to the extent to which it was used.

FUNCTION IALMOD(IP,IR,IS,JP,JR,JS)

If  $JP.\lambda + JR.\theta + JS$  divides  $IP.\lambda + IR.\theta + IS$ , the answer returned is 0; the answer 1 is returned otherwise. The result is obtained by testing whether the denominator of the quotient of these two algebraic integers is 1.

SUBROUTINE INCON(JP,JR,JS,M,ICP,ICR,ICS,LIC)

$$NORN(JJ,J) = N(JJ.\theta + J)$$

The routine sets  $ICP(I).\lambda + ICR(I).\theta + ICS(I)$  for  $I = 1, \dots, M$

to representatives of the non-zero residue classes modulo  $JP.\lambda + JR.\theta + JS$  where  $M = |N(JP.\lambda + JR.\theta + JS)| - 1$ . Also  $LIC(I)$  is set to 1 for  $I = 1, \dots, M$ .

SUBROUTINE MODAL(KP, KR, KS, ICP, ICR, ICS, LIC, M, JP, JR, JS)

$ICP(IM)$  is again  $|N(JP.\lambda + JR.\theta + JS)| - 1$ . The routine finds the values of  $I$ , where  $1 \leq I \leq M$  such that  $KP.\lambda + KR.\theta + KS$  or  $-(KP.\lambda + KR.\theta + KS)$  is congruent to  $ICP(I).\lambda + ICR(I).\theta + ICS(I)$  modulo  $JP.\lambda + JR.\theta + JS$ . For these values of  $I$ ,  $LIC(I)$  is set to 0.

When  $IPRIME$  equals 0,  $JP.\lambda + JR.\theta + JS$  does not have norm a rational prime and the representatives of the residue classes must be tested in turn until the relevant ones are found.

When  $IPRIME$  equals 1,  $JP.\lambda + JR.\theta + JS$  is an algebraic prime and we have  $\theta \equiv ITHETA$  and  $\lambda \equiv ILAM$  hence

$KP.\lambda + KR.\theta + KS \equiv KP.ILAM + KR.ITHETA + KS \equiv KPRS$

and  $-(KP.\lambda + KR.\theta + KS) \equiv -KPRS$ .

When  $IPRIME$  equals 2,  $KP.\lambda + KR.\theta + KS$  is equal to  $\theta$  or to  $\lambda$  and we only wish to find the residue class to which  $+\theta$  or  $+\lambda$  belong.

SUBROUTINE TEST

Each integer of norm  $MF$  is considered, where  $MF$  is a factor of  $N(\epsilon + 1)$  or of  $N(\epsilon - 1)$ . We have  $IM.\lambda + JM.\theta + KM$  equal to  $\epsilon + 1$  or to  $\epsilon - 1$ , as appropriate, and  $JP.\lambda + JR.\theta + JS$  is the integer of norm  $MF$  currently being considered.  $ICL = MF - 1$ .

Originally  $MF$  is the power to which  $MF$  is a factor of

$IM.\lambda + JM.\theta + KM$ , it is reduced by 1 each time that an integer of norm  $MF$  has been considered in the manner described if so, we find the representatives of the set of non-zero residue classes modulo  $JP.\lambda + JR.\theta + JS$ , otherwise we turn to the next integer of norm  $MF$ . The set of representatives is  $ICP(I).\lambda + ICR(I).\theta + ICS(I)$  for  $I = 1, \dots, ICL$ . We note that if and only if  $JP.\lambda + JR.\theta + JS$  has norm a rational prime, the set of representatives of its non-zero residue classes will be  $1, \dots, ICL$  hence  $ICS(ICL) = ICL$ . When  $JP.\lambda + JR.\theta + JS$  has norm a rational prime,  $IPRIME$  is set to 2 and we find the values  $ILAM$  and  $ITHETA$  such that  $ILAM \equiv \lambda (JP.\lambda + JR.\theta + JS)$  and  $ITHETA \equiv \theta (JP.\lambda + JR.\theta + JS)$ ,  $IPRIME$  is then set to 1; the initial setting  $ILAM = ITHETA = 0$  allows for the case when  $JP.\lambda + JR.\theta + JS$  divides  $\theta$  or  $\lambda$ . Those values  $LIC(I)$  which have been set to 0 in finding  $ILAM$  and  $ITHETA$  are reset to 1. In the case when the norm of  $JP.\lambda + JR.\theta + JS$  does not have norm a rational prime  $IPRIME$  is set to 0.

We now test for the residue classes to which  $+\epsilon$  and  $-\epsilon$  belong, and then find which other residue classes contain integers of norm less than  $MF$ ; finally  $JCOUNT$  is the number of residue classes which do not contain such an integer. If  $JCOUNT$  equals 0 we go on to the next integer of norm  $MF$ . When  $JCOUNT$  does not equal 0 the representatives of the residue classes which do not contain an integer of norm less than  $MF$  are printed together with the nature of  $JP.\lambda + JR.\theta + JS$ , and we set  $EUCLID$  to 1 to indicate that the field does not possess a Euclidean Algorithm.

Originally  $MP$  is the power to which  $MF$  is a factor of



$N( IM.\lambda + JM.\theta + KM )$ , it is reduced by 1 each time that an integer of norm  $MF$  has been considered in the manner described above, which indicates that the integer is a factor of  $IM.\lambda + JM.\theta + KM$ ;  $MP = 0$  causes exit from the routine.

SUBROUTINE DIVC(L1,M1,N1,L2,M2,N2,L,M,N)

KDET(I1,I2,I3,J1,J2,J3,K1,K2,K3) is the single precision

version of the double precision function FDET.

The routine yields the quotient

$L.\lambda + M.\theta + N = (L1.\lambda + M1.\theta + N1)/(L2.\lambda + M2.\theta + N2)$  of  $\epsilon - 1$ ,

and is used only when we know that the quotient is an algebraic

integer. The routine also includes tests for the possibility

of integer overflow.

SUBROUTINE MULTC(L1,M1,N1,L2,M2,N2,L,M,N)

The routine gives the product

$L.\lambda + M.\theta + N = (L1.\lambda + M1.\theta + N1).(L2.\lambda + M2.\theta + N2)$

and also checks for the possibility of integer overflow.

DOUBLE FUNCTION DNORM(IP,IR,IS)

The answer returned is

$N( IP.\lambda + IR.\theta + IS )$

and checks are made for the possibility of integer overflow.

THE MAIN PROGRAM

The program reads the details of the field to give defining

polynomial  $x^3 - IA.x^2 + IB.x - IC$ , discriminant  $-IDET$  and the

index of the polynomial relative to the field is INDEX; the

fundamental unit of the field is  $\frac{I.\theta^2 + J.\theta + K}{L}$ . The subroutine



BASE is called; then we use the routine CHANGE to give the rational integers  $I_1, J_1, K_1$  such that  $\epsilon = I_1.\lambda + J_1.\theta + K_1$ . We now set the rational integers AA, BB, CC, DD, EE, FF such that the ideal  $(AA, BB.\theta + CC, DD.\lambda + EE.\theta + FF)$  is the unit ideal, then call PRIN(1) to find the reduced basis of this ideal.

EUCLID is set to 0 to indicate that the Euclidean property of the field is unknown, and ICHECK is set to -1 as we begin to consider congruences modulo integers which are factors of  $\epsilon - 1$ .

Eventually  $IP(I).\lambda + IR(I).\theta + IS(I)$  for  $I = 1, \dots, MAPC - 1$  will be the integers so far calculated which produce distinct ideals; thus MAPC points to the first "empty" element of the arrays

IP, IR, IS and so is set to 1 here. We now set

$IM.\lambda + JM.\theta + KM = \epsilon - 1$  and MAX to be the greatest member of the set

$\{ f: f = p \text{ or } f = p^3, p \text{ a rational prime, } f \text{ divides } |N(\epsilon - 1)| \text{ and } f < 500 \}$ .

On this cycle MLX will be set to 2 and we now turn to the determination of all distinct algebraic integers of norm of absolute value between 2 and MAX inclusive. These integers are determined according to the theory of chapter 3, and we have the integers of norm INORM as  $IP(I).\lambda + IR(I).\theta + IS(I)$  for I between MAP(INORM,1) and MAP(INORM,2).

The method allows for a maximum of 2000 integers of norm of absolute value at most MAX. Reaching this limit, or an attempt to exceed it, would result in the printing of a diagnostic

and the value of MAX being reduced so that the above condition is satisfied.

We now go on to consider the rational integers MF, where MF is either a rational prime or the cube of a rational prime, MF divides  $N( IM.\lambda + JM.\theta + KM )$ , and MF is at most MAX. If, after a call on the subroutine TEST, EUCLID is equal to 1, we know that the field does not possess a Euclidean Algorithm and so proceed no further. When all rational integers MF have been considered we set  $IM.\lambda + JM.\theta + KM$  to  $\epsilon + 1$ , and ICHECK to 1 to indicate that factors of  $\epsilon + 1$  are being considered. We now repeat the above process noting that:

if MLX is first set to MAX, MAX is set to the

greatest member of the set

$\{ f: f = p \text{ or } f = p^3, p \text{ a rational prime, } f \text{ divides } |N(\epsilon + 1)| \text{ and } f < 500 \}$ ,

we only need to calculate the integers of norm of absolute value between MLX and MAX on this second cycle.

We consider congruences modulo integers with norm the cube of a rational prime to ensure that all the rational prime factors of  $\epsilon + 1$  and  $\epsilon - 1$  are taken into account.

THE PROGRAMS CUBOID, FCUB AND CUBX

These three programs, which are basically similar, are those used to put the method of chapter 4 into practice. CUBOID is that used for the first subdivision of the smallest cuboid which contains the fundamental region, then FCUB is used for subsequent subdivisions. CUBX is used to find the smallest cuboid which contains the uncovered sub-regions of a particular region, when the number of such sub-regions is large.

The function subroutine CI is the single precision version of the function subroutine DCI of chapter 7.

SUBROUTINE BASIS

This subroutine is basically the same as the subroutine BASEH of chapter 7, with the following exceptions. Only the values T and S where  $\lambda = (\theta^2 + T\theta + S)/INDEX$  are found and not the other coefficients in the expressions for  $\theta^2, \lambda^2$  and  $\theta\lambda$ . No bounds to be used for the prevention of integer overflow are found since they are not necessary for this program. Additionally, this routine adjusts the basis, given in the arrays R, U and H, to be that one which corresponds to the basis  $(1, M'(\theta), Q'(\theta))$  for the field as described in 4.8. Details of the basis of the field and that of the lattice are printed only for the program CUBOID.

SUBROUTINE BOUND(R,AL,AU)

We note that a basis for a fundamental region of the lattice

is part of  $\mathcal{J}$  is found such that we show that the sub-region with  
 vertices  $(0,0,0)$ ,  $(1,0,1) = (R(1),U(1),H(1))$ ,  $(R(2),U(2),H(2))$ ,

$$(X(I + \epsilon_1), Y(J + \epsilon_2), Z(K + \epsilon_3)) = (R(3),U(3),H(3)) \quad 9.1$$

where  $R, U, H$  are the COMMON arrays of the main program.

Thus the vertices of this region are given by  $(\xi, \eta, \zeta)$  where

$$\begin{aligned} \xi &= x.R(1) + y.R(2) + z.R(3) \\ \eta &= x.U(1) + y.U(2) + z.U(3) \\ \zeta &= x.H(1) + y.H(2) + z.H(3) \end{aligned}$$

and  $x, y, z$  each take the values  $(0, 1)$  which are in regions

BOUND is used to find the maximum and minimum values of  $\xi, \eta$  or  $\zeta$  when  $R$  is the main program array  $R, U$  or  $H$  respectively.

SUBROUTINE USE-DEF, defining polynomial  $x^3 - 14x^2 + 14x - 1$

This routine is used to determine the set  $\mathcal{J}$  of integer points to be used in testing.  $\mathcal{J}$  consists of the points corresponding to the integers of  $K$  given by  $PM(I) + PN(I)\theta + PO(I)\lambda$  for  $I = 1, \dots, KZ$ . If  $KZ$  reaches the value 1000 no more integers are added to the set.

SUBROUTINE COVER(X,Y,Z,L,ICOUNT)

For each element of  $\mathcal{J}$ , this routine determines for which sub-regions of the region under consideration 4.2 is satisfied, that is we have these sub-regions covered. If at any stage we find that all sub-regions are covered, a return to the calling program is made immediately. ICOUNT is the number of sub-regions left uncovered after consideration of any element of  $\mathcal{J}$ . When an



element of  $\mathcal{f}$  is found such that we show that the sub-region with vertices

$$(X(I + t_i), Y(J + t_j), Z(K + t_k)) \quad \text{for } t_i = 0,1; \quad t_j = 0,1; \quad t_k = 0,1, \quad 9.3$$

is covered, we set  $L(I,J,K)$  to 0.

We note that for the program CUBOID we are considering the whole of the cuboid containing the fundamental region, and this cuboid is symmetrical about the origin; so we need only consider those regions containing points  $(\xi, \eta, \zeta)$  which are in regions 9.3 where  $K$  takes values between 6 and 10.

#### THE MAIN PROGRAM CUBOID

The details of the field are first read; the field has discriminant  $-IDET$ , defining polynomial  $x^3 - IA \cdot x^2 + IB \cdot x - IC$  and this polynomial has index relative to the field  $INDEX$ .

The subroutine BASIS is called to determine the basis of the field. We now use the subroutine BOUND to find the values  $XL, XU, YL, YU, ZL, ZU$  such that the smallest cuboid containing the fundamental region with vertices 9.2 is the set of points  $(\xi, \eta, \zeta)$  such that

$$XL \leq \xi \leq XU, \quad YL \leq \eta \leq YU, \quad ZL \leq \zeta \leq ZU.$$

We then adjust these values so that the similar fundamental region which is symmetric about the origin is contained in the cuboid which is the set of points  $(\xi, \eta, \zeta)$  such that



For now with  $|\xi| \leq XU = -XL$  for those sub-regions  
 with the not integer  $|\eta| \leq YU = -YL$  fundamental regions. We will  
 show that  $|\zeta| \leq ZU = -ZL$ , the others,

BK is assigned the value C of chapter 4. CIN = .1 is the ratio of the side of a sub-region which we are attempting to cover to that of the original cuboid containing the fundamental region, that is the value  $10^{-h}$  of chapter 4. AM was a value used when the program was employed to find the inhomogeneous minimum of a field; it was a measure of the amount by which the value of C was to be altered if the number of uncovered sub-regions was greater than the value LIMIT. Since we are only concerned here with the Euclidean property of the field, LIMIT was made so large that the value of C was not adjusted.

IBC is the greatest absolute value of the coefficients of the integer points, this bound was used to prevent the number of elements of  $\mathcal{J}$  becoming too large. For the fields with the smaller values of IDET this value was fixed at 10 in the program, but it was found necessary to be able to make it smaller as the value of IDET increased. The subroutine USE is then called.

The sub-regions are to be those with vertices 9.3 for  $I = 1, \dots, 10; J = 1, \dots, 10; K = 1, \dots, 10;$  and by symmetry we need only consider the values  $K = 6, \dots, 10$ . For these values of I, J, K we set  $L(I,J,K)$  to 1 to indicate that the corresponding sub-region has not yet been "covered", in the sense of 6.6. Then we determine the values of the elements of the arrays X, Y and Z.

We now wish to set  $L(I, J, K)$  to 2 for those sub-regions which do not intersect with the fundamental region. We note that if  $(\xi, \eta, \zeta)$  is a point of the cuboid,

$$\begin{aligned} \xi &= x.R(1) + y.R(2) + z.R(3) \\ \eta &= x.U(1) + y.U(2) + z.U(3) \\ \zeta &= x.H(1) + y.H(2) + z.H(3) \end{aligned} \tag{9.4}$$

for some real numbers  $x, y, z$ . This point is also a point of the fundamental region only if

$$|x| \leq \frac{1}{2}, \quad |y| \leq \frac{1}{2}, \quad |z| \leq \frac{1}{2}$$

Thus a sub-region intersects with the fundamental region only if at least one of its vertices satisfies these conditions.

It should be pointed out that the determinant of the coefficients in 9.4 is equal to 1

$$\begin{vmatrix} 1 & \text{Re}\phi - k & \text{Re}\psi - p & \text{Re}\phi - q \\ 0 & \text{Im}\phi & \text{Im}\psi - p & \text{Im}\phi \\ 1 & \theta - k & \lambda - p & \theta - q \end{vmatrix} = \begin{vmatrix} 1 & \text{Re}\phi & \text{Re}\psi \\ 0 & \text{Im}\phi & \text{Im}\psi \\ 1 & \theta & \lambda \end{vmatrix} = \frac{1}{2i} \begin{vmatrix} 1 & \phi & \psi \\ 1 & \bar{\phi} & \bar{\psi} \\ 1 & \theta & \lambda \end{vmatrix}$$

which has absolute value  $\frac{i\Delta}{2i} = \frac{\Delta}{2}$ , where  $\Delta = \text{IDET}^{\frac{1}{2}}$ .

But, from the choice of  $\text{Im}\phi$  as positive, we have

$$\begin{vmatrix} 1 & \text{Re}\phi & \text{Re}\psi \\ 0 & \text{Im}\phi & \text{Im}\psi \\ 1 & \theta & \lambda \end{vmatrix} = \frac{1}{\text{INDEX}} \begin{vmatrix} 1 & \text{Re}\phi & \text{Re}\phi^2 \\ 0 & \text{Im}\phi & \text{Im}\phi^2 \\ 1 & \theta & \theta^2 \end{vmatrix}$$

$$\begin{aligned}
 \text{THE MAIN PROGRAM} &= \frac{(\text{Im}\phi \cdot \theta^2 + \text{Re}\phi \cdot \text{Im}\phi^2 - \text{Re}\phi^2 \cdot \text{Im}\phi - \text{Im}\phi^2 \cdot \theta)}{\text{INDEX}} \\
 &= (\text{Im}\phi \cdot \theta^2 + \text{Re}\phi \cdot 2 \cdot \text{Re}\phi \cdot \text{Im}\phi - ((\text{Re}\phi)^2 - (\text{Im}\phi)^2) \cdot \text{Im}\phi \\
 &\quad - 2 \cdot \text{Re}\phi \cdot \text{Im}\phi \cdot \theta) / \text{INDEX} \\
 &= \frac{\text{Im}\phi}{\text{INDEX}} (\theta^2 - 2 \cdot \text{Re}\phi \cdot \text{Im}\phi + (\text{Re}\phi)^2 + (\text{Im}\phi)^2) \\
 &> 0
 \end{aligned}$$

thus the value of the determinant is  $\frac{1}{2} \text{IDET}^{\frac{1}{2}}$ . Consequently, the efficiency of this program could have been improved by using this expression.

We now call the subroutine COVER to find the number ICOUNT of sub-regions which are left uncovered, and print the value of ICOUNT. If ICOUNT is not equal to 0, and is less than LIMIT, the values of X(I), Y(J), Z(K) for those regions for which L(I,J,K) is equal to 1 are printed; these values are also punched onto cards for later use by FCUB.

#### THE MAIN PROGRAM FCUB

This program finds the smallest cuboid containing the fundamental region, then reads the values BK and IBC as in CUBOID. The value CIN =  $10^{-h}$  must be read by this program, then the subroutine USE is called.

LIMIT is again the maximum number of uncovered regions for which the vertices are to be printed. M is the number of regions which are to be subdivided. Each of these regions is now treated as the cuboid containing the fundamental region was treated by CUBOID; with the exception that we must now consider the values  $I = 1, \dots, 10$ ;  $J = 1, \dots, 10$ ;  $K = 1, \dots, 10$ .

This program is basically the same as FCUB with the following exceptions. We do not read a value LIMIT. On return from the

subroutine COVER we find the values XLW, XUP, YLW, YUP, ZLW, ZUP

such that the uncovered sub-regions of the particular region in question are contained in the cuboid which consists of the set of

points  $(\xi, \eta, \zeta)$  where

$$XLW \leq \xi \leq XUP$$

$$YLW \leq \eta \leq YUP$$

FUNCTION ADJ(R,U,I,J,ZLW  $\leq \zeta \leq$  ZUP.

This determines the value of

$$R(I),U(J) - R(J),U(I) = \begin{vmatrix} R(I) & U(I) \\ R(J) & U(J) \end{vmatrix}$$

SUBROUTINE UPLN(RL,RU,RL)

On return from this subroutine

$$RL = \min(RL, RU)$$

$$RU = \max(RU, RL)$$

SUBROUTINE NCHG(I,IS)

IS takes the value of IS rounded to the nearest integer.

FUNCTION INDEX(I,\*)

When this function subroutine is called, I is a bound on the co-ordinates of one of the measured points. J = 1 indicates that I is a lower bound and the answer returned is I + 0.5. J = 2 indicates that I is an upper bound and the answer returned is I - 0.5. In this way we avoid the possibility of rounding

## CHAPTER 10

### THE PROGRAM TRANS

From the programs of chapter 9 we find regions  $R_1, \dots, R_n$ , we wish to determine the intersections of the transformations of these regions by the fundamental unit with the regions themselves; this program is used to achieve this. The subroutine BASIS and the function subroutine CI are the same as for the programs of chapter 9.

### FUNCTION ADJ(R,U,I,J)

This determines the value of

$$R(I).U(J) - R(J).U(I) = \begin{vmatrix} R(I) & U(I) \\ R(J) & U(J) \end{vmatrix}$$

### SUBROUTINE UPLW(RL,RU,TR)

On return from this subroutine

$$RL = \min(RL, TR)$$

$$RU = \max(RU, TR)$$

### SUBROUTINE ROUNDI(RC,IR)

IR takes the value of RC rounded to the nearest integer.

### FUNCTION DCORR(X,J)

When this function subroutine is called, X is a bound on the co-ordinates of one of the uncovered regions. J = 1 indicates that X is a lower bound and the answer returned is X - 0.1. J = 2 indicates that X is an upper bound and the answer returned is X + 0.1. In this way we avoid the possibility of rounding



error concealing the intersection of an uncovered region with the transform of some uncovered region.

#### THE MAIN PROGRAM TRANS

We note that if the point  $(\xi, \eta, \zeta)$  has cylindrical polar co-ordinates  $(\rho, \alpha, \zeta)$ , and the point of the lattice corresponding to the fundamental unit has cylindrical polar co-ordinates  $(\rho_\epsilon, \alpha_\epsilon, \zeta_\epsilon = \epsilon)$ , the transform of the point by the fundamental unit has cylindrical polar co-ordinates

$$(\rho\rho_\epsilon, \alpha + \alpha_\epsilon, \zeta\zeta_\epsilon).$$

We will define  $R_{-J}$  to be the set of points  $\tilde{\xi}$  where  $-\tilde{\xi}$  is in  $R_J$ .  $R_J + \tilde{\gamma}$  will be defined as the set of points  $\tilde{\xi} + \tilde{\gamma}$  where  $\tilde{\xi}$  is in  $R_J$ .

TRANS reads the values EXC, EYC, EZC where

$$\epsilon = \text{EXC} \cdot \text{H}(1) + \text{EYC} \cdot \text{H}(2) + \text{EZC} \cdot \text{H}(3).$$

It should be noted that EXC, EYC, EZC are the coefficients of  $\epsilon$  relative to the basis  $(1, \mathcal{M}'(\theta), \mathcal{Q}'(\theta))$  and not  $(1, \theta, \lambda)$ .  $(\text{EX}, \text{EY}, \text{EZ})$  are the cartesian co-ordinates of the point  $\tilde{\epsilon}$  of the lattice of integers, and  $(\text{ER}, \text{EALP}, \text{EZ})$  are its cylindrical polar co-ordinates.

N is the number of regions to be considered; these regions are  $R_I$  for  $I = 1, \dots, N$  where  $R_I$  is the set of points  $(\xi, \eta, \zeta)$  such that

$$X(1, I) \leq \xi \leq X(2, I)$$

$$Y(1, I) \leq \eta \leq Y(2, I)$$

$$Z(1, I) \leq \zeta \leq Z(2, I)$$

for  $I = 1, \dots, N$ . When N is large, which usually means

greater than 15, the amount of time taken to consider the intersection of the transform of each region with the appropriate translates of the original region becomes large; consequently, M1 and M2 are read such that the transforms of the regions  $R_I$  for  $I = M1, \dots, M2$  are considered in this run.

We now consider each region in turn. For the vertex  $(X(II,I), Y(JJ,I), Z(KK,I))$  we determine its cylindrical polar co-ordinates  $(RA,ALP,Z(KK,I))$ , and from these the cylindrical polar co-ordinates of its transform by the fundamental unit  $(TRA,TALP,TZ)$ ; we note that RA and ALP, and consequently TRA and TALP, are independent of  $Z(KK,I)$  and so may be determined in an outer loop. This vertex is then the point  $\tilde{\xi}$  where  $\xi = TXC.H(1) + TYC.H(2) + TZC.H(3)$ . After inspecting the eight vertices, the transformed region is the set of points  $(\xi, \eta, \zeta)$  where

$$XL \leq \xi \leq XU$$

$$YL \leq \eta \leq YU$$

$$ZL \leq \zeta \leq ZU;$$

and if  $\xi = x.H(1) + y.H(2) + z.H(3)$

$$XLC \leq x \leq XUC$$

$$YLC \leq y \leq YUC$$

$$ZLC \leq z \leq ZUC.$$

Since the fundamental region consists of the points  $\tilde{\xi}$  where  $\xi = x.H(1) + y.H(2) + z.H(3)$  and  $|x| \leq \frac{1}{2}$ ,  $|y| \leq \frac{1}{2}$ ,  $|z| \leq \frac{1}{2}$ , if XLC, XUC, YLC, YUC, ZLC, ZUC, when rounded to the nearest integer, become IXL, IXU, IYL, IYU, IZL, IZU respectively, we

must consider for each integer point  $\tilde{\delta}$ , where

$\tilde{\delta} = x.H(1) + y.H(2) + z.H(3)$ ,  $x$  is between  $IXL$  and  $IXU$ ,  $y$  is between  $IYL$  and  $IYU$ ,  $z$  is between  $IZL$  and  $IZU$ , whether  $R_J + \tilde{\delta}$  or  $R_{-J} + \tilde{\delta}$  intersects with  $E(R_I)$  for  $J$  between 1 and  $N$ .  $IFLAG$  is set to 1 when at least one such region  $R_J$  is found.

## CHAPTER 11

### THE PROGRAM EXCEP

The method of chapter 5 is implemented by this program. The subroutine BASES is basically the same as the subroutine BASEH of RELMIN, but the bounds which it finds for the prevention of integer overflow are for the arguments of DNORM and MULTD. The subroutine MULTD is the double precision version of the subroutine MULTC of the program CONG. The function subroutine DCI is the same routine as that for the program RELMIN. The function subroutine DAJ is the double precision version of the function ADJ of the program TRANS. The subroutine CHANGE and the function subroutine DNORM are the same as those for the program CONG.

### LOWUP(XT,XV,XW,XL,XU)

The routine finds XL and XU such that

$$XL \leq XT.Y(1) + XV.Y(2) + XW.Y(3)$$

where  $YL(I) \leq Y(I) \leq YU(I)$  for  $I = 1,2,3$

and the arrays YL and YU are passed across to the subroutine as COMMON arrays.

### LIMITL(XXL,IXL)

IXL is the least integer greater than XXL. A diagnostic is given if the resultant integer would require more than 14 digits; that is if the integer would come close to the limits of accuracy of the machine.

## LIMITU (XXU, IXU)

IXU is the greatest integer less than XXU. A diagnostic is again given if the resultant integer would require more than 14 digits.

## THE MAIN PROGRAM EXCEP

The field being considered has discriminant  $-IDET$ , defining polynomial  $x^3 - IA \cdot x^2 + IB \cdot x - IC$  and the index of the polynomial relative to the field is INDEX. BASES is called to find the details of the basis of the field and of the lattice of integers.

If we are finding the minimum of  $\frac{\alpha}{1 - \epsilon^n}$ , then  $\alpha = \frac{IZ \cdot \theta^2 + IY \cdot \theta + IX}{ID}$

and  $1 - \epsilon^n = \frac{JZ \cdot \theta^2 + JY \cdot \theta + JX}{JD}$  and  $H(2) = \theta$ . We then have

$$AX + AY \cdot \theta + AZ \cdot \lambda = I1X + I1Y \cdot \theta + I1Z \cdot \lambda = \frac{IZ \cdot \theta^2 + IY \cdot \theta + IX}{ID}$$

$$BX + BY \cdot \theta + BZ \cdot \lambda = J1X + J1Y \cdot \theta + J1Z \cdot \lambda = \frac{JZ \cdot \theta^2 + JY \cdot \theta + JX}{JD}$$

and we obtain  $AN = \alpha$ ,  $BE = 1 - \epsilon^n$ .

The transformation on the lattice is  $\Omega \rightarrow \frac{\Omega - \alpha}{\epsilon^n} = \frac{\Omega - AN}{1 - BE}$ , so we determine the values  $BP = 1 - BE$  and  $ABP = -AN / (1 - BE)$ .  $CK = SK$  is the value of  $C_1$  of chapter 5;  $CK$  is a double precision quantity,  $SK$  is single precision.

The conjugates of  $\alpha$  are  $ANV \pm i \cdot ANW$  and those of  $1 - \epsilon^n$  are  $BEV \pm i \cdot BEW$ , then the conjugates of  $\frac{\alpha}{1 - \epsilon^n}$  are  $ABV \pm i \cdot ABW$ .  $BNORM$  is the norm of  $1 - \epsilon^n$ .

$TL1 = SL1$  is the value  $\tau_1$  of chapter 5; then  $TU1 = \frac{TL1 - \alpha}{\epsilon^n}$ . We are to investigate the range of values of  $\zeta_4$  of chapter 5 in sections given by

$$TL \leq \zeta_4 < TU \quad \text{such that} \quad TU - TL \leq 100.$$



With this in mind we set  $TL = TL1$  and  $TU = \min(TU1, TL + 100)$ .

RMINIM is to eventually take the value of the minimum of  $\frac{\alpha}{1 - \epsilon^n}$ ; it is initially set to the value SK, which should be chosen by the user so that it will be greater than the minimum.

NEXT is set to 1 when we are considering values of  $\zeta_x$  greater than  $\frac{\alpha}{1 - \epsilon^n}$ , and to 2 when we are considering values of  $\zeta_x$  less than  $\frac{\alpha}{1 - \epsilon^n}$ .

SB is the value  $G_1$  of chapter 5. We must consider the integers  $\tilde{\zeta} = (\zeta_x, \eta_x, \zeta_y)$  such that

$$VL \leq \zeta_x \leq VU$$

$$WL \leq \eta_x \leq WU$$

$$TL \leq \zeta_y \leq TU.$$

If 
$$\zeta_y = p \cdot \lambda + q \cdot \theta + r$$

for some rational integers p, q and r, we investigate the values of p and q where

$$IPL \leq p \leq IPU$$

$$IQL \leq q \leq IQU,$$

the range of values for r will be found for any particular pair of values of p and q. If the range of integer values for p or that for q is null, a diagnostic is printed and the program terminates. We determine which of p and q has the smaller range of values. If p has the smaller range, for each value of p we determine the integer values IQLT, IQUT such that we must consider the reduced range of values of q given by

$$IQLT \leq q \leq IQUT.$$

Then for each value of q in this range, we determine the integer

values ISLT, ISUT such that we must consider the values of  $r$  in the range

$$ISLT \leq r \leq ISUT.$$

For each value of  $r$  we then determine the value of  $N\left(\frac{\alpha}{1 - \epsilon^n} - \delta\right)$  in the form  $N(\alpha - \delta(1 - \epsilon^n))/N(1 - \epsilon^n)$ . When  $q$  has the smaller range, for each value of  $q$  we determine the range for  $p$ , then for each value of  $p$  we determine the range for  $r$ .

Having considered the range of values of  $\xi_r$  given by  $TL \leq \xi_r \leq TU$ , if  $TU$  is less than  $TU1$ ,  $TL$  is set to the value  $TU$  and  $TU$  to  $\min(TU1, TU + 100)$  then the above process is repeated; when  $TU = TU1$  we go on to consider values of  $\xi_r$  less than  $\frac{\alpha}{1 - \epsilon^n}$ . First  $NEXT$  is set to  $NEXT + 1 = 2$ ,  $TU1 = SU1$  is the value of  $\tau_n$  of chapter 5, then  $TL1 = \frac{TU1 - \alpha}{\epsilon^n}$ . Again we are to consider values of  $\xi_r$  in "sections", therefore, we set  $TU = TU1$  and  $TL = \max(TU - 100, TL1)$ .  $SB$  is now the value  $\sigma_2$  of chapter 5. The relevant integers are investigated in the manner described above; then, if  $TL > TL1$ , we set  $TU = TL$  and then  $TL = \max(TU - 100, TL1)$ , otherwise we print the value of  $RMINIM$  and stop.

## CHAPTER 12

### THE RESULTS OBTAINED

The following table includes all the results obtained by the methods of this thesis. For each field we give the discriminant,  $D$ , the coefficients  $a, b, c$  of the defining polynomial  $x^3 - ax^2 + bx - c$ , and the index,  $L$ , of the polynomial relative to the field, as given in the table in (1).

The "property" of the field is indicated by  $E$  if the field possesses a Euclidean Algorithm and by  $N$  if it does not. The last column has an entry when further information about the field is included at the end of this chapter, the entry is a reference to that information.

The results for the first fifteen fields, those with discriminants between  $-23$  and  $-152$  inclusive, were obtained by Godwin (18); these fields were not investigated by the methods of this thesis, but the results are included here for the sake of completeness.

D	L	a, b, c			Property	
-23	1	4	5	1	E	(1)
-31	1	3	4	1	E	(1)
-44	1	4	6	2	E	(1)
-59	1	3	5	2	E	(1)
-76	1	2	4	2	E	(1)
-83	1	5	9	4	E	(1)
-87	1	4	7	3	E	(1)
-104	1	6	11	4	E	(1)
-107	1	2	4	1	E	(1)
-108	1	6	12	6	E	(1)
-116	1	5	8	2	E	(1)
-135	1	3	6	3	E	(1)
-139	1	4	6	1	E	(1)
-140	1	3	5	1	E	(1)
-152	1	8	19	10	E	(1)
-172	1	7	15	7	E	(2)
-175	1	5	10	5	E	(3)
-199	1	2	5	3	N	(4)
-200	1	4	7	2	E	
-204	1	5	9	3	E	
-211	1	6	10	1	E	
-212	1	2	5	2	E	
-216	1	3	6	2	E	
-231	1	7	16	9	E	

D	L	a, b, c			Property
-239	1	6	11	3	E
-243	1	6	12	5	E
-244	1	10	29	18	E
-247	1	6	13	7	E
-255	1	5	8	1	E
-268	1	7	13	1	E
-300	1	8	18	6	E
-307	1	4	8	3	N
-324	1	9	24	14	E
-327	1	8	19	9	N (5)
-335	1	4	9	5	N
-339	1	7	15	6	N
-351	1	3	6	1	N
-356	2	5	12	4	E
-364	1	3	7	3	N
-367	1	4	7	1	N
-379	1	5	9	2	E
-411	1	2	6	3	E
-419	1	9	23	10	E
-424	2	7	16	4	E
-431	2	9	26	16	E
-436	1	6	13	6	N (6)
-439	1	7	14	3	N
-440	2	6	14	4	E



D	L		a, b, c		Property	
-451	1	10	28	13	E	
-459	1	9	21	2	N	
-460	1	2	6	2	E	
-472	1	9	22	6	E	
-484	1	4	9	4	E	
-492	1	4	8	2	E	
-499	1	3	7	2	E	
-503	2	7	18	8	E	
-515	1	8	20	11	E	
-516	2	8	22	12	E	
-519	1	10	29	17	E	(7)
-524	1	5	11	5	N	
-527	1	3	8	5	N	
-543	1	5	10	3	E	(8)
-547	1	8	18	5	N	
-567	1	9	24	13	N	
-620	1	11	35	23	N	
-628	2	7	20	12	E	
-652	2	4	12	4	E	
-655	1	7	16	7	N	
-671	1	6	11	1	N	
-675	1	6	12	3	N	
-679	1	6	13	5	N	
-687	1	4	9	3	E	(9)

D	L		a, b, c		Property
-695	1	8	21	13	N
-755	1	4	10	5	N
-759	1	2	7	3	N
-808	1	5	10	2	N
-812	1	11	33	13	N
-823	1	9	22	5	N
-839	1	8	19	7	N
-863	1	7	18	11	N
-883	2	8	23	12	N
-908	2	9	23	3	N
-940	1	9	25	15	N
-959	1	4	11	7	N
-972	1	6	12	2	N
-983	1	2	7	1	N
-1004	1	6	14	6	N
-1007	1	7	14	1	N
-1011	1	11	35	22	N
-1036	2	6	16	4	N
-1048	2	5	16	8	N
-1059	1	7	17	8	N
-1075	2	7	18	4	N
-1087	1	5	12	5	N
-1135	1	10	27	7	N
-1147	1	10	30	19	N

D	L		a, b, c		Property
-1164	1	7	15	3	N
-1172	2	7	21	11	N
-1175	1	3	8	1	N
-1187	1	2	8	5	N
-1191	1	14	55	39	N
-1196	1	5	13	7	N
-1207	1	11	34	17	N
-1208	2	4	13	2	N
-1219	1	12	40	21	N
-1231	1	10	29	15	N
-1235	1	10	28	11	N
-1259	1	13	47	28	N
-1267	1	2	8	3	N
-1291	1	9	25	14	N
-1292	1	12	38	10	N
-1295	1	7	16	5	N
-1315	1	13	45	16	N
-1316	2	4	14	4	N
-1319	1	9	26	17	N
-1327	1	6	13	3	N
-1351	1	8	21	11	N
-1355	1	6	14	5	N
-1363	1	11	31	2	N
-1383	1	4	11	5	N

D	L	a, b, c			Property
-1388	1	2	8	2	N
-1407	1	11	32	7	N
-1431	1	6	15	7	N
-1448	2	6	17	4	N
-1452	1	10	26	2	N
-1547	1	7	15	2	N
-1567	1	8	19	5	N
-1579	1	6	16	9	N
-1580	1	4	10	2	N
-1583	1	11	36	25	N
-1599	3	7	24	9	N
-1603	2	12	43	28	N
-1615	1	3	10	5	N
-1619	1	8	22	13	N
-1647	1	12	39	15	N
-1675	1	4	12	7	N
-1687	1	9	22	3	N
-1700	2	2	13	4	N
-1708	1	10	28	10	N
-1736	2	9	29	17	N
-1743	1	13	48	33	N
-1751	3	11	36	9	N
-1755	2	9	30	20	N
-1763	1	14	52	19	N

D	L	a, b, c			Property
-1772	1	7	19	11	N
-1807	1	2	9	5	N
-1815	1	5	12	3	N
-1868	2	7	23	13	N
-1871	1	2	9	7	N
-1895	1	12	41	25	N
-1955	1	11	35	20	N
-1959	1	2	9	3	N
-1967	3	8	27	9	N
-2023	1	4	11	3	N
-2036	2	2	14	8	N
-2039	1	7	20	13	N
-2047	1	3	10	3	N
-2063	3	10	31	3	N
-2159	1	10	29	13	N
-2167	1	13	44	9	N
-2199	3	9	30	9	N
-2207	1	10	27	5	N
-2228	1	6	17	10	N
-2235	1	10	28	9	N
-2283	3	9	33	18	N
-2315	1	5	13	4	N
-2316	1	7	17	5	N
-2372	1	4	13	8	N



D	L		a, b, c		Property
-2403	1	12	42	29	N
-2420	2	8	25	10	N
-2423	3	11	42	27	N
-2444	1	9	23	5	N
-2479	1	12	37	3	N
-2491	2	10	31	10	N
-2503	1	7	18	7	N
-2515	1	12	40	19	N
-2540	2	11	39	25	N
-2579	1	11	33	10	N
-2591	3	10	39	27	N
-2599	1	11	36	23	N
-2604	1	2	10	6	N
-2627	1	2	10	7	N
-2636	2	12	44	23	N
-2647	1	4	13	7	N
-2668	1	9	25	11	N
-2699	2	9	26	4	N
-2708	2	9	29	13	N
-2723	1	3	11	4	N
-2732	1	6	14	2	N
-2759	1	5	12	1	N
-2791	1	10	31	19	N
-2795	2	5	18	4	N

D	L		a, b, c		Property
-2796	1	5	13	3	N
-2803	1	5	15	8	N
-2855	1	14	55	37	N
-2860	2	10	32	12	N
-2879	1	4	11	1	N
-2895	3	9	30	5	N
-2915	1	8	20	5	N
-2951	1	5	14	5	N
-2956	1	6	16	6	N
-3011	1	9	23	4	N
-3039	1	7	20	11	N
-3043	2	6	19	2	N
-3055	1	15	58	11	N
-3059	2	6	23	14	N
-3107	1	13	43	2	N
-3127	1	16	69	37	N
-3148	2	16	72	52	N
-3159	1	3	12	7	N
-3159	1	12	39	13	N
-3176	2	3	17	7	N
-3179	1	13	45	14	N
-3191	1	12	41	23	N
-3212	2	6	20	4	N
-3235	1	9	25	10	N

D	L		a, b, c		Property
-3244	1	12	40	18	N
-3259	1	4	14	9	N
-3263	1	9	26	13	N
-3308	2	8	24	4	N
-3311	1	8	21	7	N
-3331	1	5	13	2	N
-3359	1	4	13	5	N
-3371	1	3	11	2	N
-3404	2	2	16	4	N
-3439	1	14	51	11	N
-3451	1	10	28	7	N
-3483	2	3	18	12	N
-3495	1	10	27	3	N
-3543	3	10	37	15	N
-3547	3	8	28	3	N
-3560	2	13	45	5	N
-3575	3	7	28	9	N
-3591	1	3	12	5	N
-3599	1	11	34	13	N
-3615	1	2	11	7	N
-3619	1	10	32	21	N
-3647	1	7	18	5	N
-3671	1	8	19	1	N
-3687	3	6	27	11	N

D	L		a, b, c		Property
-3695	1	2	11	5	N
-3711	1	5	16	9	N
-3756	2	5	19	3	N
-3767	1	6	17	7	N
-3783	3	5	26	13	N
-3820	2	4	20	12	N
-3831	1	11	36	21	N
-3915	2	9	30	12	N
-3916	1	7	19	7	N
-3935	3	4	25	15	N
-3943	1	13	46	19	N
-3980	2	11	35	5	N
-3991	1	2	11	3	N
-3999	3	9	36	19	N
-4007	1	10	31	17	N
-4012	1	8	24	14	N
-4023	3	6	27	9	N
-4039	1	5	14	3	N
-4075	1	14	52	17	N
-4076	2	7	27	17	N
-4111	1	16	71	51	N
-4147	1	8	20	3	N
-4175	1	8	23	11	N
-4191	3	3	24	17	N

D	L		a, b, c		Property
-4239	1	3	12	3	N
-4255	1	7	20	9	N
-4291	2	14	51	2	N
-4295	3	11	44	25	N
-4307	1	9	23	2	N
-4359	1	17	80	61	N
-4364	1	5	17	11	N
-4411	2	4	19	4	N
-4455	1	9	24	5	N
-4487	1	14	53	23	N
-4511	1	10	33	23	N
-4555	2	7	26	12	N
-4556	2	10	36	20	N
-4559	1	9	26	11	N
-4567	1	9	28	17	N
-4575	3	2	23	19	N
-4639	1	6	19	11	N
-4691	1	8	24	13	N
-4699	1	14	54	29	N
-4715	2	15	62	28	N
-4727	1	5	16	7	N
-4735	1	14	55	35	N
-4771	2	8	27	8	N
-4823	1	13	48	29	N



D	L	a, b, c			Property	
-4831	1	17	76	29	N	
-4844	1	8	20	2	N	
-4859	1	2	12	7	N	
-4903	3	10	41	23	N	
-4908	1	2	12	6	N	
-4935	3	7	30	9	N	
-4963	1	18	88	59	N	
-10011	3	10	42	9	N	
-10019	2	16	75	56	N	
-10028	2	17	75	7	N	
-19899	3	18	96	61	N	
-19927	2	6	31	2	N	
-19939	1	9	31	12	N	
-19976	2	22	129	68	N	
-160087	1	0	1	77	N	(10)
-169571	1	1	35	22	N	(10)

We now have the additional information concerning the indicated fields.

- (1) These results are those obtained by Godwin (18).
- (2) The inhomogeneous minimum is  $\frac{3}{4}$  and is attained at numbers congruent to  $\pm\left(\frac{9\theta^2 - 25\theta + 14}{1 - \epsilon}\right)$ .
- (3) The inhomogeneous minimum is  $\frac{3}{5}$  and is attained at numbers congruent to  $\pm\left(\frac{3\theta^2 - 6\theta + 3}{1 - \epsilon}\right)$ .
- (4) The inhomogeneous minimum is 1 and is attained at numbers congruent to  $\pm\left(\frac{3\theta}{1 - \epsilon^2}\right)$ .
- (5) The inhomogeneous minimum is  $\frac{101}{99}$  and is attained at numbers congruent to  $\pm\left(\frac{4\theta^2 - 4\theta + 1}{1 - \epsilon^2}\right)$  and transforms of them by the fundamental unit.
- (6) The inhomogeneous minimum is  $\frac{79}{78}$  and is attained at numbers congruent to  $\pm\left(\frac{29\theta^2 - 33\theta + 10}{1 - \epsilon^2}\right)$ .
- (7) Numbers congruent to  $\pm\left(\frac{19\theta^2 - 213\theta - 155}{1 - \epsilon^4}\right)$ , and transforms of them by the fundamental unit, have minimum  $\frac{44712}{45747}$ , thus the inhomogeneous minimum of the field is at least this value but is less than 1.
- (8) Numbers congruent to  $\pm\left(\frac{-5\theta^2 + 68\theta - 23}{1 - \epsilon^3}\right)$ , and transforms of them by the fundamental unit, have minimum  $\frac{7593}{8343}$ , thus the inhomogeneous minimum of the field is at least this value but is less than 1.
- (9) Numbers congruent to  $\pm\left(\frac{9\theta^2 - 27\theta + 8}{1 - \epsilon^2}\right)$ , and transforms of them by the fundamental unit, have minimum  $\frac{937}{945}$ , thus the

inhomogeneous minimum of the field is at least this value  
but is less than 1.

(10) These results were obtained by the method of chapter 2.

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and their subroutines. Each program listing is preceded by a title page which gives:

1. the program name;
2. the names of those subroutines which it uses and which are listed in the same section;
3. the names of those subroutines which it uses and which have been listed in a previous section.

#### APPENDIX

Listings of those routines copied from Angel(1), or which are modifications of subroutines of Angel are included for completeness.

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ROYAL HOLLOWAY COLLEGE, LONDON

Supervisor : PROFESSOR H J GODWIN

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and their subroutines. Each program listing is preceded by a title page which gives

1. the program name
2. the names of those subroutines which it uses and which are listed in the same section
3. the names of those subroutines which it uses and which have been listed in a previous section.

Listings of those routines copied from Angell(1), or which are modifications of subroutines of Angell are included for completeness.

```

PROGRAM RELMIN(OUTPUT)
C... THIS PROGRAM FINDS A LOWER BOUND FOR THE INHOMOGENEOUS MINIMUM OF
C... THE FIELD.
C... 01/10/73 JTC, ENCF, IDCI
C... 02/11/73 IL2, IL3, IT1, IT2, IT3, IP1, IP2, IP3
C... IN DETAILS OF THE FIELD, A BLANK CARD WHICH SETS IDCI TO ZERO.
C... INDICATES THE END OF THE DATA DECK.
C... 03/10/73 IDCI, INDE, IS, IS*E
C... 04/10/73
C... 05/10/73 STOP
C... THE FIELD NOW BEING CONSIDERED HAS DISCRIMINANT -IDET, IT IS XIXI
C... WHERE XIXI IS A SQUARE.
15 10/10/73
PRINT 102, TOE, INDE, IS, IS*E
100 FORMAT(1H1,1F,5H THE FIELD HAS DISCRIMINANT,1F,6H INDEX,1J,2BH AND
100 POLYNOMIAL COEFFICIENTS,1F,10H
C... NOW FIND X WHERE
C... THE PROGRAM RELMIN IS A BASIS FOR THE FIELD.
BASE, BASEH
PRINT 103
103 09/10/73
ASSIG=0
PRINT 104, ASSIG, IT1, IT2, IT3
104 10/10/73
PRINT 105, ASSIG, IT1, IT2, IT3
105 10/10/73
C... NOW FIND THE
C... MINIMUM FOR THIS FIELD
C... CALC MINIMA
C... REPEAT THE CYCLE
GO TO 102
END

```

THE PROGRAM RELMIN IS A BASIS FOR THE FIELD.

SUBROUTINES - ALL LISTED IN THIS SECTION

BASEH	5	MULT	19
MINIMA	8	PHITH	20
SUB	15	DIVCD	21
INVER	16	MULTCD	22
ICF	17	DCI	23
ICF2	18	IOF	24



```

PROGRAM RELMIN(INPUT,OUTPUT)
C
C...THIS PROGRAM FINDS A LOWER BOUND FOR THE INHOMOGENEOUS MINIMUM OF
C...THE FIELD.
5 C
COMMON/D1/IA,IB,IC,INDEX,IDET
COMMON/D2/IL1,IL2,IL3,IT1,IT2,IT3,IP1,IP2,IP3
C...READ IN DETAILS OF THE FIELD. A BLANK CARD WHICH SETS IDET TO ZERO.
C...INDICATES THE END OF THE DATA DECK.
10 2 READ 101,IDET,INDEX,IA,IB,IC
101 FORMAT(I6,I2,3I5)
IF (IDET.EQ.0) STOP
C...THE FIELD NOW BEING CONSIDERED HAS DISCRIMINANT -IDET, IT IS K(X)
C...WHERE  $X^{*3} - IA * X^{*2} + IB * X - IC = 0$ .
15 IDETM = -IDET
PRINT 102, IDETM, INDEX, IA, IB, IC
102 FORMAT(1H1,1X,26H THE FIELD HAS DISCRIMINANT,1B,6H INDEX,1I,28H AND
1 POLYNOMIAL COEFFICIENTS,3I6)
C...NOW FIND X WHERE THE FIELD IS K(X) AND ALSO A BASIS FOR THE FIELD.
20 CALL BASEH
PRINT 103
103 FORMAT(/,1X,45H THE FIELD IS K(X) AND HAS BASIS (Y,X,1) WHERE)
ASSIG1=4HX**2 & ASSIG2=4HY**2 & ASSIG3=4HX*Y
25 PRINT 104,ASSIG1,IT1,IT2,IT3
PRINT 104,ASSIG2,IL1,IL2,IL3
PRINT 104,ASSIG3,IP1,IP2,IP3
104 FORMAT(12X,A4,2H =,1I0,3HY +,1I0,3HX +,1I10)
C...NOW FIND THE LOOP OF IDEALS STARTING WITH 1 AND A CHAIN OF RELATIVE
C...MINIMA FOR THIS FIELD.
30 CALL MINIMA
C...REPEAT THE CYCLE
GO TO 2
END

```

```

SUBROUTINE BASEH
DOUBLE R,U,H,A,B,C,V1,V2,V3,V4,V5,V6,V7,D,T,S,XIN
DOUBLE IORD,IORDD
COMMON/D1/IA,IB,IC,INDEX,IDET
5 COMMON/D2/IL1,IL2,IL3,IT1,IT2,IT3,IP1,IP2,IP3
COMMON/D3/R(3),U(3),H(3)
COMMON/D10/PH,ZZ,HHH/D9/IQ,INN,Q,RN
COMMON/D16/IORD,IORDD
C
10 C...WHEN CONSIDERING THE FIELD K(X) SUPPOSE PH=3*X-IA THEN PH
C...SATISFIES AN EQUATION PH**3-IQ*PH-INN=0. FIND PH AND PH*PH-IQ
C...FOR USE IN THE SUBROUTINE MINIMA.
C
15 A=FLOAT(IA) & B=FLOAT(IB) & C=FLOAT(IC)
IQ=3*IA*IA-9*IB
INN=27*IC-9*IA*IB+2*IA*IA*IA
Q=FLOAT(IQ)
RN=FLOAT(INN)
20 V1=FLOAT(IQ/3)
V2=FLOAT(INN)
V3=DSQRT(V2**2-4.000*V1**3)
V4=(V2+V3)/2.000
V5=(V2-V3)/2.000
IF(V4.NE.0.000) GO TO 51
25 V6=0.000 & GO TO 52
51 D=1.000
IF(V4.GE.0.000) GO TO 53
V4=-V4 & D=-1.000
53 V6=DEXP((DLOG(V4))/3.000)*D
30 52 IF(V5.NE.0.000) GO TO 54
V7=0.000 & GO TO 55
54 D=1.000
IF(V5.GE.0.000) GO TO 56
35 V5=-V5 & D=-1.000
56 V7=DEXP((DLOG(V5))/3.000)*D
55 PH=SNGL(V6+V7) & HHH=PH*PH
ZZ=HHH-Q
H(2)=(V6+V7+A)/3.000
R(2)=(A-H(2))/2.000
40 U(2)=DSQRT(C/H(2)-R(2)*R(2))
H(1)=1.000 & R(1)=1.000 & U(1)=0.000
C...R(N)+I.U(N) AND R(N)-I.U(N) ARE THE CONJUGATES OF H(N) FOR N=1,2,3.
C
C...(H(1),H(2),H(3))=(I,X,Y) IS A BASIS FOR THE FIELD.
45 C...Y=(X**2+IT*X+IS)/INDEX, IT AND IS ARE NOW CALCULATED.
C
IN=INDEX
IN2=IN**2 & IN3=IN2*IN
50 IF(INDEX.NE.1) GO TO 1
IT=0 & IS=0 & GO TO 2

```

```

1 K1=IA**2-2*IB & K4=IB**2-2*IA*IC & K5=IA*IB-3*IC
DO 3 ITC=1,IN
IT=ITC-1
DO 4 ISC=1,IN
55 IS=ISC-1
JA=K1+IT*IA+3*IS
IF (MOD(JA,IN),NE.0) GO TO 4
JB=K4+IT**2*IB+3*IS**2+IT*K5+2*IS*K1+2*IS*IT*IA
IF (MOD(JB,IN2),NE.0) GO TO 4
60 JC=IS**3+IT**3*IC+IC**2+IS**2*IT*IA+IS*K4+IS*IT**2*IB+IT*IB*IC
1+IT**2*IA*IC+IS**2*K1+IS*IT*K5
IF (MOD(JC,IN3),NE.0) GO TO 4
GO TO 2
4 CONTINUE
65 3 CONTINUE
2 IT1=INDEX
IT2=-IT
IT3=-IS
ML=2*IA*IT+IA*IA+2*IS+IT*IT-IB
70 IL1=ML/IN
IL2=(-IT*ML+2*IS*IT+IC-2*IB*IT-IA*IB)/IN2
IL3=(-IS*ML+2*IC*IT+IA*IC+IS*IS)/IN2
MP=IA+IT
75 IP1=MP
IP2=(IS-IB-IT*MP)/IN
IP3=(IC-IS*MP)/IN
T=FLOAT(IT) & S=FLOAT(IS) & XIN=FLOAT(INDEX)
H(3)=(H(2)*H(2)+T*H(2)+S)/XIN
R(3)=(R(2)*R(2)-U(2)*U(2)+T*R(2)+S)/XIN
80 U(3)=(2.0D0*R(2)*U(2)+T*U(2))/XIN
C
C...NOW FIND A BOUND ON THE RELEVANT PRODUCTS OF ARGUMENTS OF MULTCD AND
C...DIVCD FOR TESTING FOR THE POSSIBILITY OF OVERFLOW.
C
85 IL1A=IABS(IL1) & IL2A=IABS(IL2) & IL3A=IABS(IL3)
IT1A=IABS(IT1) & IT2A=IABS(IT2) & IT3A=IABS(IT3)
IP1A=IABS(IP1) & IP2A=IABS(IP2) & IP3A=IABS(IP3)
IOBL=MAX0(IL1A,IL2A,IL3A)
IOBT=MAX0(IT1A,IT2A,IT3A)
90 IOBP=MAX0(IP1A,IP2A,IP3A)
ILTP=MAX0(IOBL,IOBT,IOBP)
ILTP1=IOF(ILTP)
IOBS=29-ILTP1
IORD=10.0D0**IOBS
IF (IOBL.GT.IOBP) GO TO 20
IF (IOBP.GT.IOBT) GO TO 21
IOBAS=IOBP*IOBT
GO TO 22
100 21 IOBAS=IOBP*IOBP
GO TO 22

```



```

SUBROUTINE MINIMA
DOUBLE L(3,99),DA,DB,DC,ZPL,DD,DE,DF,ZPJ,Z0,X1,Y1,XY1
DOUBLE REMN,RNMN,R,U,H,DCI
COMMON/D1/IA,IB,IC,INDEX,IDET
5 COMMON/D2/KX,KY,KZ,IX,IY,IZ,JX,JY,JZ
COMMON/D3/R(3),U(3),H(3)
COMMON/D7/M1,M2,M3,N1,N2,N3,A,B,C/DB/MB(3)
COMMON/D10/PH,ZZ,HHH/D9/I0,IN,Q,RN
10 DIMENSION LI(8),JAN(3,99),K(4,3),NPJ(99),NP(99),IAN(3,99)
C...RHO IS THE SQUARE OF THE DISTANCE OF THE POINT P+V*PH+W*PH**2
C...FROM THE REAL AXIS.
RHO(P,V,W)=((P+W*Q)**2-V*(V*Q+W*RN))+(W*W*RN-V*P)*PH
1+ (V*V-W*(P+W*Q))*HHH
N=1
15
C**** SECTION 1 ****
C...BEGIN WITH THE IDEAL (1,X,Y).
LI(1)=1 & LI(4)=0 & LI(3)=INDEX & LI(2)=0
20 LI(7)=1 & LI(6)=-IY & LI(5)=-IZ & LI(8)=INDEX
C...BASIS FOR LATTICE IS
C... ( I )
1 (1,(M1+M2*PH+M3*PH**2)/IG,(N1+N2*PH+N3*PH**2)/IG).
IG=9*LI(1)*LI(8) & M1=9*LI(2)+3*IA*LI(3)+LI(4)*IA*IA
N1=9*LI(5)+3*IA*LI(6)+LI(7)*IA*IA
25 M2=3*LI(3)+2*LI(4)*IA
N2=3*LI(6)+2*LI(7)*IA
M3=LI(4) & N3=LI(7)
C..... A,B,C ARE COEFFICIENTS OF THE ASSOCIATED BINARY QUADRATIC FORM
A=R=C=0.
30 C...PARTS (I),.....,(VIII) ARE THE STEPS OF VORONOI'S ALGORITHM AS
DESCRIBED IN DELONE AND FADDEEV.
C.... ( I ) .....
C..... TEST THAT M2*N3-M3*N2 IS GREATER THAN 0 .....
35 10 IF(M2*N3-M3*N2.GT.0.) GO TO 20
I1=M1 & I2=M2 & I3= M3
M1=M1 & M2=N2 & M3= N3
N1=I1 & N2=I2 & N3=I3
C.... ( II ) .....
40 C..... PRODUCTION OF INITIAL VALUES OF A,B,C .....
20 R1=M1 & R2=M2 & R3=M3
S1=M1 & S2=N2 & S3=N3
G=IG
A=R2*R2+R2*R3*PH+R3*R3*ZZ
45 B=R2*S2+(R2*S3+R3*S2)*PH*0.5+R3*S3*ZZ
C=S2*S2+S2*S3*PH+S3*S3*ZZ
C.... ( III ) .....
C... (III) AND (IV) FIND THE TWO BASIS ELEMENTS OF THE REDUCED HEXAGON
50 C...OF ZELLING.
IF(R.GT.0.) GO TO 22
CALL SUB(0,1,-1,0)

```



```

C.... ( TV ) .....
22 IF (A-R.LT.0.) GO TO 32
   IF (C-R.GT.0.) GO TO 50
55 32 IF (A-C.GT.0.) GO TO 40
   ID=INT(B/A)
   CALL SUB(1,-ID,0,1)
   GO TO 22
   40 ID=INT(B/C)
   CALL SUB(1,0,-ID,1)
60 GO TO 22
C.... ( V ) .....
C...TO FIND THE TWO ZELLING HEXAGON BASIS ELEMENTS WHICH COVER THE
C...NEGATIVE 'XSI' AXIS.
65 50 R2=M2 & R3=M3 & S2=N2 & S3=N3
   RB=R2-R3*PH & RD=S2-S3*PH
   RUP=RD-RB
   IF (RB.LT.0.) GO TO 60
   IF (RD.LT.0.) GO TO 90
   IF (RUP.GT.0.) GO TO 55
70 CALL SUB(0,-1,1,1)
   GO TO 90
55 CALL SUB(-1,-1,1,0)
   GO TO 90
60 IF (RD.LT.0.) GO TO 70
75 CALL SUB(-1,0,0,-1)
   GO TO 90
70 IF (RD-RB.GT.0.) GO TO 80
   CALL SUB(1,1,-1,0)
   GO TO 90
80 80 CALL SUB(0,1,-1,-1)
C.... ( VI ) .....
C...FINDING THE PINHEADS CORRESPONDING TO (M1+M2*PH+M3*PH**2)/IG AND
C...(N1+N2*PH+N3*PH**2)/IG.
85 90 R1=M1 & R2=M2 & R3=M3
   S1=N1 & S2=N2 & S3=N3
   T=(R1+R2*PH+R3*PH**2)/G
   I=IFIX(T)
   IF (T.GT.0.) GO TO 4001
   I=I-1
90 4001 M1=M1-I*IG
   S=(S1+S2*PH+S3*PH**2)/G
   I=IFIX(S)
   IF (S.GT.0.) GO TO 4002
   I=I-1
95 4002 N1=N1-I*IG
C.... ( VII ) .....
C...CHOOSING THE TWO ELEMENTS OF THE REDUCED BASIS OF THE LATTICE FROM
C...THE SEVEN POSSIBILITIES BY FINDING THE TWO WITH MINIMUM RHO VALUE.
100 200 R1=M1 & S1=N1 & G2=2.*G
   XA=(2.*(R1+R3*Q)-R2*PH-R3*PH**2)/G2

```

```

XC=(2.*(S1+S3*Q)-S2*PH-S3*HHH)/G2
IF(XA-0.5.GT.0.) GO TO 201
K(1,1)=M1 & K(1,2)=M2 & K(1,3)=M3
GO TO 202
105 201 K(1,1)=IG-M1 & K(1,2)= -M2 & K(1,3)= -M3
202 IF(XC-0.5.GT.0.) GO TO 203
K(2,1)=N1 & K(2,2)=N2 & K(2,3)=N3
GO TO 204
110 203 K(2,1)=IG-N1 & K(2,2)= -N2 & K(2,3)= -N3
204 T=(R1+R2*PH+R3*HHH)/G
S=(S1+S2*PH+S3*HHH)/G
IF(T-S.LT.0.) GO TO 220
IF(XA-XC.GT.0.5) GO TO 210
K(3,1)=M1-N1 & K(3,2)=M2-N2 & K(3,3)=M3-N3
115 GO TO 240
210 K(3,1)=IG-M1+N1 & K(3,2)= -M2+N2 & K(3,3)= -M3+N3
GO TO 240
220 IF(XC-XA.GT.0.5) GO TO 230
K(3,1)=N1-M1 & K(3,2)=N2-M2 & K(3,3)=N3-M3
120 GO TO 240
230 K(3,1)=IG-N1+M1 & K(3,2)= -N2+M2 & K(3,3)= -N3+M3
240 RZ0=RHO(FLOAT(K(1,1)),FLOAT(K(1,2)),FLOAT(K(1,3)))
RZ1=RHO(FLOAT(K(2,1)),FLOAT(K(2,2)),FLOAT(K(2,3)))
RZ2=RHO(FLOAT(K(3,1)),FLOAT(K(3,2)),FLOAT(K(3,3)))
125 IF(RZ2.LT.RZ0) GO TO 280
IF(RZ2.LT.RZ1) GO TO 260
IF(T+S.GT.1.)GO TO 250
IF(XA+XC.GT.0.5) GO TO 250
130 K(4,1)=M1+N1 & K(4,2)=M2+N2 & K(4,3)=M3+N3
RZ3=RHO(FLOAT(K(4,1)),FLOAT(K(4,2)),FLOAT(K(4,3)))
IF(RZ3.GT.RZ0) GO TO 245
IF(RZ3.GT.RZ1) GO TO 244
M1=K(4,1) & M2=K(4,2) & M3=K(4,3)
IF(RZ0.GT.RZ1) GO TO 243
135 N1=K(1,1) & N2=K(1,2) & N3=K(1,3)
GO TO 300
243 N1=K(2,1) & N2=K(2,2) & N3=K(2,3)
GO TO 300
244 M1=K(2,1) & M2=K(2,2) & M3=K(2,3)
140 N1=K(4,1) & N2=K(4,2) & N3=K(4,3)
GO TO 300
245 IF(RZ3.GT.RZ1) GO TO 250
M1=K(1,1) & M2=K(1,2) & M3=K(1,3)
N1=K(4,1) & N2=K(4,2) & N3=K(4,3)
145 GO TO 300
250 IF(RZ0.GT.RZ1) GO TO 255
M1=K(1,1) & M2=K(1,2) & M3=K(1,3)
N1=K(2,1) & N2=K(2,2) & N3=K(2,3)
150 GO TO 300
255 M1=K(2,1) & M2=K(2,2) & M3=K(2,3)

```

```

      N1=K(1,1) & N2=K(1,2) & N3=K(1,3)
      GO TO 300
260 M1=K(1,1) & M2=K(1,2) & M3=K(1,3)
      N1=K(3,1) & N2=K(3,2) & N3=K(3,3)
155 GO TO 300
      280 IF(RZ1.GT,RZ2) GO TO 290
      M1=K(2,1) & M2=K(2,2) & M3=K(2,3)
      N1=K(3,1) & N2=K(3,2) & N3=K(3,3)
      GO TO 300
160 290 IF(RZ0.GT,RZ1) GO TO 295
      M1=K(3,1) & M2=K(3,2) & M3=K(3,3)
      N1=K(1,1) & N2=K(1,2) & N3=K(1,3)
      GO TO 300
165 295 M1=K(3,1) & M2=K(3,2) & M3=K(3,3)
      N1=K(2,1) & N2=K(2,2) & N3=K(2,3)
      C..... ( VIII ) .....
300 JAN(1,N)=N1 & JAN(2,N)=N2 & JAN(3,N)=N3 & NP(N)=IG
      IAN(1,N)=M1 & IAN(2,N)=M2 & IAN(3,N)=M3
      N00=N-1
170 C..... FIND THE INVERSE OF THE SECOND BASIS ELEMENT,THE FIRST RELATIVE
      C..... OF THE LATTICE .....
8988 CALL INVER(M1,M2,M3,JDET)
      I1=MB(1)*IG & I2=MB(2)*IG & I3=MB(3)*IG
      CALL ICF(I1,I2,I3,IF1)
175 C..... DIVIDE LATTICE BY SECOND BASIS ELEMENT AND PRODUCE A NEW LATTICE
      CALL MULT(MB(1),MB(2),MB(3),N1,N2,N3,J1,J2,J3)
      CALL ICF(J1,J2,J3,IF2)
      CALL ICF(IF1,IF2,JDET,IF)
180 M1=J1/IF & M2=J2/IF & M3=J3/IF
      N1=I1/IF & N2=I2/IF & N3=I3/IF & IG=JDET/IF
      IF(N.EQ.1) GO TO 501
      C...CHECK IF OLD LATTICE COMPLETED THE LOOP.
      IF(IAN(1,N).NE.IAN11) GO TO 502
185 IF(IAN(2,N).NE.IAN21) GO TO 502
      IF(IAN(3,N).NE.IAN31) GO TO 502
      IF(JAN(1,N).NE.JAN11) GO TO 502
      IF(JAN(2,N).NE.JAN21) GO TO 502
      IF(JAN(3,N).NE.JAN31) GO TO 502
      IF(NP(N).EQ.NP1) GO TO 600
190 GO TO 502
      501 IAN11=IAN(1,1) & IAN21=IAN(2,1) & IAN31=IAN(3,1)
      JAN11=JAN(1,1) & JAN21=JAN(2,1) & JAN31=JAN(3,1) & NP1=NP(1)
      502 N=N+1
195 C.....IF N.GT.99 PRINT DIAGNOSTIC.....
      IF(N.LT.99) GO TO 10
      PRINT 1111
1111 FORMAT(39H LATTICE LOOP HAS MORE THAN 99 MEMBERS )
      GO TO 7
200 C**** SECTION 2 *****

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```

C...NOW PRINT OUT THE LATTICES IN THIS CHAIN.
600 PRINT 112
205 112 FORMAT(/,1X,81HTHE ELEMENTS OF THE LOOP ARE, WHEN EXPRESSED IN TER
MS OF THE BASIS (1,PHI,PHI**2))
DO 700 JIM=1,N
PRINT 113,IAN(1,JIM),IAN(2,JIM),IAN(3,JIM),NP(JIM),JAN(1,JIM),
1JAN(2,JIM),JAN(3,JIM),NP(JIM)
113 FORMAT(/,1X,3I15,* /*,15,/,3X,3I15,* /*,15)
210 700 CONTINUE
PRINT 114
114 FORMAT(/,1X,57HWHEN EXPRESSED IN TERMS OF THE BASIS (Y,X,1) THESE
1 BECOME)
C...WHEN THE THIRD ELEMENT OF EACH LATTICE BASIS HAS BEEN PRINTED THEN
215 C...IT IS NO LONGER NEEDED AND SO THE ARRAY JAN MAY BE USED TO STORE THE
C...VALUES OF THE COEFFICIENTS OF THE SECOND ELEMENT IN TERMS OF THE
C...BASIS (Y,X,1), SO THAT SECOND ELEMENT IS
C... JAN(1,JIM)*Y+JAN(2,JIM)*X+JAN(3,JIM).
220 DO 800 JIM=1,N
CALL PHITH(IAN(1,JIM),IAN(2,JIM),IAN(3,JIM),NP(JIM),NIAN1,NIAN2,
NIAN3,NNPI)
CALL PHITH(JAN(1,JIM),JAN(2,JIM),JAN(3,JIM),NP(JIM),NJAN1,NJAN2,
225 NJAN3,NNPJ)
PRINT 113,NIAN1,NIAN2,NIAN3,NNPI,NJAN1,NJAN2,NJAN3,NNPJ
230 JAN(1,JIM)=NIAN1
JAN(2,JIM)=NIAN2
JAN(3,JIM)=NIAN3
NPJ(JIM)=NNPI
800 CONTINUE
C
C...THE RELATIVE MINIMA ARE TO BE L(1,ILAT)*Y+L(2,ILAT)*X+L(3,ILAT) FOR
C...ILAT=1,N. THE RELATIVE MINIMUM 1 IS TO BE PUT AT 'ABOUT ONE THIRD'
C...OF THE WAY ALONG THE CHAIN.
235 NQH=N/3
NQP=NQH+1
NQM=N
L(1,NQP)=0.000
L(2,NQP)=0.000
L(3,NQP)=1.000
240 C...NOW CALCULATE THOSE RELATIVE MINIMA WHICH COME 'BEFORE' 1.
DO 900 ILAT=1,NQH
NOPP=NQP
NQP=NQP-1
NQM=NQM-1
245 DD=JAN(1,NQM) & DE=JAN(2,NQM) & DF=JAN(3,NQM)
ZPJ=NPJ(NQM)
CALL DIVCD(L(1,NQPP),L(2,NQPP),L(3,NQPP),1.000,DD,DE,DF,ZPJ,DA,DB,
1DC,7PL)
L(1,NQP)=DA
250 L(2,NQP)=DB

```

```

      L(3,NQP)=DC
900 CONTINUE
C...NOW CALCULATE THOSE RELATIVE MINIMA WHICH COME 'AFTER' 1.
255     NQU=NQO-NQH & NQPP=NQH+1
        DO 901 ILAT=1,NQU
          NQP=NQPP
          NQPP=NQPP+1
          DD=JAN(1,ILAT) & DE=JAN(2,ILAT) & DF=JAN(3,ILAT)
          ZPJ=NPJ(ILAT)
260     CALL MULTCD(L(1,NQP),L(2,NQP),L(3,NQP),1.D0,DD,DE,DF,ZPJ,DA,DB,
          1DC,ZPL)
          L(1,NQPP)=DA
          L(2,NQPP)=DB
          L(3,NQPP)=DC
265     901 CONTINUE
C
C...NOW PRINT OUT THIS CHAIN OF RELATIVE MINIMA.
C...FOR N=1 TO NQO CALCULATE Z(N)*(X(N+1)*X(N+1)+Y(N+1)*Y(N+1)) WHERE
270 C...Z(N) IS THE NTH RELATIVE MINIMUM OF THE CHAIN AND X(N)+I.Y(N),
      C...X(N)-I.Y(N) ARE ITS CONJUGATES.
C
      PRINT 115
115  FORMAT(/,1X,34HTHE RELATIVE MINIMA ARE AS FOLLOWS)
      Z0=DCI(L(3,1),L(2,1),L(1,1),H)
275     REMN=1.000
          PRINT 116,L(1,1),L(2,1),L(3,1)
116  FORMAT(/,1X,3D24.16)
      DO 902 ILAT=2,N
          DD=L(1,ILAT) & DE=L(2,ILAT) & DF=L(3,ILAT)
280     PRINT 116,DD,DE,DF
          X1=DCI(DF,DE,DD,R)
          Y1=DCI(DF,DE,DD,U)
          XY1=X1*X1+Y1*Y1
285     RNMN=DABS(Z0*XY1)
          PRINT 990,Z0,X1,Y1,XY1,RNMN
990  FORMAT(/,1X,3HZ0=,D24.16,4H X1=,D24.16,4H Y1=,D24.16,/,1X,11HGIVIN
          1G XY1=,D24.16,10H AND RNMN=,D24.16)
          Z0=DCI(DF,DE,DD,H)
          IF (RNMN.GT.REMN) REMN=RNMN
290     902 CONTINUE
C
C...NOW FIND AND PRINT THE CALCULATED LOWER BOUND ON THE MINIMUM OF THE
C...FIELD.
C
295     PRINT 970,REMN
970  FORMAT(6H REMN=,D24.16)
      REMNR=REMN
      DET=FLOAT(IDET)
      VALUE=DET/(720.0*REMNR)
300     PRINT 980,VALUE

```





```

SUBROUTINE SUB(I ,J,K,L)DET1
C
C . . . ALTERS THE BASIS OF THE LATTICE AND CALCULATES THE NEW VALUES OF A, B AND C.
C . . . THE INVARIANT IS  $(I^2+J^2+K^2+L^2)^2$ 
5 C . . . COMMON/D7/M1,M2,M3,N1,N2,N3,A,B,C
C . . . RI=I & RJ=J & RK=K & RL=L
C . . . A1=A*RI*RI+2.*B*RI*RK+C*RK*RK
C . . . B1=RI*RJ*A+B*(RI*RL+RJ*RK)+C*RK*RL
10 C . . . C1=A*RJ*RJ+2.*B*RJ*RL+C*RL*RL
C . . . A=A1 & B=B1 & C=C1
C . . . I1=M1 & I2=M2 & I3=M3 & J1=N1 & J2=N2 & J3=N3
C . . . M1=I*I1+K*J1 & M2=I*I2+K*J2 & M3=I*I3+K*J3
C . . . N1=J*I1+L*J1 & N2=J*I2+L*J2 & N3=J*I3+L*J3
15 C . . . RETURN
C . . . END DET1
C . . . DO 100 I=1,100
C . . . 100 CONTINUE
C . . . CALL TOP(N1,N2,N3,N1,N2,N3)
C . . . CALL TOP(I2,I2,KDET,I2)
C . . . N1=N1/21 & N2=N2/21 & N3=N3/21
C . . . KDET=KDET/21
C . . . RETURN
C . . . END

```

```

      SURROUTINE INVER(I,J,K,KDET)
C
C...FINDS THE INVERSE OF THE ALGEBRAIC INTEGER I+J*PH+K*PH**2
C...THE INVERSE IS (N(1,1)+N(1,2)*PH+N(1,3)*PH**2)/KDET.
5  C
      COMMON/D9/IQ,IN,Q,RN/D8/N(3)
      LA=I  &  LB=J  &  LC=K
      LD=IN*K  &  LE=I+IQ*K  &  LF=J
      LG=IN*J  &  LH=IQ*J+IN*K  &  LI=I+IQ*K
10  N(1)=LE*LI-LH*LF
      N(2)=LC*LH-LB*LI  &  N(3)=LB*LF-LE*LC
      MA=N(1)  &  MB=LG*LF-LD*LI  &  MC=LD*LH-LG*LE
      JDET=LA*MA+LB*MB+LC*MC
15  KDET=IARS(JDET)
      IF(JDET.GT.0) GO TO 10
      DO 5 JZ=1,3
      N(JZ)=-N(JZ)
      5  CONTINUE
      10  CONTINUE
20  CALL TCF(N(1),N(2),N(3),IZ)
      CALL TCF(IZ,IZ,KDET,IY)
      N(1)=N(1)/IY  &  N(2)=N(2)/IY  &  N(3)=N(3)/IY
      KDET=KDET/IY
25  RETURN
      END

```

```

SUBROUTINE ICF(IJ,IK,IL,IH)
C
C..... FINDS IH ,THE H.C.F. OF THE ABSOLUTE VALUES OF 3 INTEGERS IJ,IK,IL.
C
5      N=1 & J=IABS(IJ) & K=IABS(IK) & L=IABS(IL)
1      IF(K.EQ.0) GO TO 7
      IF(J.NE.0) GO TO 9
      J=K & GO TO 1
7      IF(J.NE.0) GO TO 8
10     IH=IL & GO TO 50
8      K=J & GO TO 1
9      IF(J.LT.K) GO TO 10
      I=J & J=K & K=I
10     M=MOD(K,J)
15     IF(M.EQ.0) GO TO 30
      IF(M.EQ.1) GO TO 31
      K=J & J=M & GO TO 10
30     IH=J
      IF(N.EQ.2) GO TO 50
20     N=N+1 & J=IH & K=L & GO TO 1
31     IH=1
50     RETURN
      END

```

```

SUBROUTINE ICF2(IJ,IK,IL,IH,J,K,L,H,I)
C
C .. DOUBLE PRECISION VERSION OF ICF
C
5   C   DOUBLE IJ,IK,IL,IH,J,K,L,H,I
      N=1
      J=IJ & K=IK & L=IL
      IF (IJ.GT.0.) GO TO 100
      J=-J
10  100 IF (IK.GT.0.) GO TO 101
      K=-K
      101 IF (IL.GT.0.) GO TO 102
      L=-L
      102 CONTINUE
15  1   IF (K.EQ.0) GO TO 7
      IF (J.NE.0) GO TO 9
      J=K & GO TO 1
      7   IF (J.NE.0) GO TO 8
      IH=L & GO TO 50
20  8   K=J & GO TO 1
      9   IF (J.LT.K) GO TO 10
      I=J & J=K & K=I
      10  H=DMOD(K,J)
      IF (H.EQ.0.) GO TO 30
25  IF (H.EQ.1) GO TO 31
      K=J & J=H & GO TO 10
      30  IH=J
      IF (N.EQ.2) GO TO 50
      N=N+1 & J=IH & K=L & GO TO 1
30  31  IH=1
      50  RETURN
      END

```

```

SUBROUTINE MULT(J1,J2,J3,K1,K2,K3,L1,L2,L3)
C
C...MULTIPLIES J1*J2*PH+J3*PH**2 BY K1+K2*PH+K3*PH**2 TO GIVE
5 C...L1+L2*PH+L3*PH**2.
C
COMMON/D9/I1,I2,0,RN
L1=J1*K1+I2*(J3*K2+J2*K3)
L2=J1*K2+K1*J2+I1*(J3*K2+J2*K3)+I2*J3*K3
10 L3=J1*K3+J2*K2+J3*K1+I1*J3*K3
RETURN
END

```

```

SUBROUTINE PHITH(I1,I2,I3,IO,J1,J2,J3,JD)
C
C... (I1+I2*PH+I3*PH**2)/IO=(J1*Y+J2*X+J3)/JD
C... (I1+I2*PH+I3*PH**2)/IO=(J1*Y+J2*X+J3)/JD
5 COMMON/D1/IA,IB,IC,INDEX,IDET
COMMON/D2/IL1,IL2,IL3,IT1,IT2,IT3,IP1,IP2,IP3
J3=I1-IA*I2+I3*IA*IA & J2=3*I2-6*IA*I3
J1=9*I3 & JD=IO
10 J3=J3+J1*IT3 & J2=J2+J1*IT2 & J1=J1*IT1
CALL ICF(JD,J1,J2,JF)
JF2=JF
CALL ICF(J3,JF,JF2,JFF)
J1=J1/JFF & J2=J2/JFF & J3=J3/JFF & JD=JD/JFF
15 RETURN
END

```



```

SUBROUTINE DIVCD(L1,M1,N1,LN,L2,M2,N2,LD,L,M,N,LR)
C
C...DIVCD IS A DOUBLE PRECISION ROUTINE WHICH FINDS THE QUOTIENT OF
C... (L1*Y+M1*X+N1)/LN AND (L2*Y+M2*X+N2)/LD; IT IS (L*Y+M*X+N)/LR.
5 C
      DOUBLE L1,M1,N1,LN,L2,M2,N2,LD,L,M,N,LR,I1,I2,I3,J1,J2,J3,K1,K2,K3
      DOUBLE JDET,KDET,IHF,LG,LRA
      DOUBLE IORD,IORDD,LA,MA,NA,DDIV1,DDIV2,DDIV
      COMMON/D2/IL1,IL2,IL3,IT1,IT2,IT3,IP1,IP2,IP3
10 COMMON/D16/IORD,IORDD
      KDET(I1,I2,I3,J1,J2,J3,K1,K2,K3)=I1*J2*K3+I2*J3*K1+I3*J1*K2
      1-I1*J3*K2-I2*J1*K3-I3*J2*K1
C...FIRST TEST FOR THE POSSIBILITY OF OVERFLOW.
      LA=DABS(L1) & MA=DABS(M1) & NA=DABS(N1)
15 DDIV1=DMAX1(LA,MA,NA)
      LA=DABS(L2) & MA=DABS(M2) & NA=DABS(N2)
      DDIV2=DMAX1(LA,MA,NA)
      IF(DDIV1.GT.DDIV2) GO TO 1
      DDIV=DDIV2*DDIV2*DDIV2
20 GO TO 2
      1 DDIV=DDIV1*DDIV2*DDIV2
      2 IF(DDIV.LT.IORDD) GO TO 3
      PRINT 50
50 FORMAT(1X,10(1H*),27HDANGER OF OVERFLOW IN DIVCD,10(1H*))
25 3 I1=L2*IL1+IP1*M2+N2
      I2=L2*IL2+IP2*M2
      I3=L2*IL3+IP3*M2
      J1=IP1*L2+M2*IT1
30 J2=IP2*L2+M2*IT2+N2
      J3=IP3*L2+M2*IT3
      K1=L2 & K2=M2 & K3=N2
      JDET=KDET(I1,I2,I3,J1,J2,J3,K1,K2,K3)
      L=KDET(L1,M1,N1,J1,J2,J3,K1,K2,K3)*LD
      M=KDET(I1,I2,I3,L1,M1,N1,K1,K2,K3)*LD
35 N=KDET(I1,I2,I3,J1,J2,J3,L1,M1,N1)*LD
      LR=JDET*LN
      CALL ICF2(L,M,N,IHF)
      LRA=LR
      CALL ICF2(IHF,LRA,LR,LG)
40 L=L/LG & M=M/LG & N=N/LG & LR=LR/LG
      RETURN
      END

```

```

SUBROUTINE MULTCD(L1,M1,N1,LN,L2,M2,N2,LD,L,M,N,LR)
C
C...MULTCD IS A DOUBLE PRECISION ROUTINE WHICH FINDS THE PRODUCT OF
C...(L1*Y+M1*X+N1)/LN AND (L2*Y+M2*X+N2)/LD, IT IS (L*Y+M*X+N)/LR.
5 C
DOUBLE L1,M1,N1,LN,L2,M2,N2,LD,L,M,N,LR,IP,Y1,Y2,Z
DOUBLE IORD,IORDD,LA,MA,NA,MULT1,MULT2,MULT
COMMON/D2/IL1,IL2,IL3,IT1,IT2,IT3,IP1,IP2,IP3
COMMON/D16/IORD,IORDD
10 C...FIRST TEST FOR THE POSSIBILITY OF OVERFLOW.
LA=DABS(L1) & MA=DABS(M1) & NA=DABS(N1)
MULT1=DMAX1(LA,MA,NA)
LA=DABS(L2) & MA=DABS(M2) & NA=DABS(N2)
MULT2=DMAX1(LA,MA,NA)
15 MULT=MULT1*MULT2
IF (MULT.LT.IORD) GO TO 1
PRINT 50
50 FORMAT(1X,10(1H*),28HDANGER OF OVERFLOW IN MULTCD,10(1H*))
1 IP=L1*M2+L2*M1
20 L=L1*L2*IL1+IP*IP1+M1*M2*IT1+L1*N2+L2*N1
M=L1*L2*IL2+IP*IP2+M1*M2*IT2+M1*N2+M2*N1
N=L1*L2*IL3+IP*IP3+M1*M2*IT3+N1*N2
LR=LN*LD
CALL ICF2(L,M,N,Y1)
25 Y2=Y1
CALL ICF2(Y2,Y1,LR,Z)
L=L/Z & M=M/Z & N=N/Z & LR=LR/Z
RETURN
END

```

```
DOUBLE FUNCTION DCI(X+Y+Z, RD)
C++ I DOUBLE X+Y+Z, RD(3)
DCI=X*RD(1)+Y*RD(2)+Z*RD(3)
RETURN
5      1 END
      2 IF (I .EQ. 2) GO TO 1
      3 DO WHILE
      4 RETURN
      5 END
```

```

      FUNCTION IOF(I)
      C...I IS LESS THAN 10**IOF(I).
      IFL=I
      IFL1=0
5      1 IFL1=IFL1+1
      IFL=IFL/10
      IF(IFL.GT.0) GO TO 1
      IOF=IFL1
10     RETURN
      END

```

THE PROGRAM CONT.

SUBROUTINES

IN THIS SECTION		IN PREVIOUS SECTION	
BASE	34	SUP	18
FACTOR	37	INVER	15
IDEAL	38	ICE	17
PRIN	40	ICE2	18
MULT2	45	MULT	19
CHANGE	46	DIVCD	21
TEST	47	ICE	24
INCON	50		
MODAL	52		
JALMOD	53		
DIVC	54		
MULTC	55		
DNORM	56		
FDET	57		



```

PROGRAM CONG(INPUT,OUTPUT)
C...THIS PROGRAM TESTS WHETHER ANY OF THE RESIDUE CLASSES MODULO THE
C...FACTORS OF E+1 AND E-1 CONTAIN ANY INTEGERS NOT CONGRUENT TO AN
C...INTEGER WITH NORM LESS IN ABSOLUTE VALUE THAN THE ABSOLUTE VALUE OF
5 C...THE NORM OF THE MODULUS UNDER CONSIDERATION. IF THIS IS SO THEN THE
C...FIELD BEING CONSIDERED IS NON-EUCLIDEAN. OTHERWISE NO CONCLUSION IS
C...REACHED.
COMMON/D20/MU1,MU2,MU3,NU1,NU2,NU3,IGU1
C
10 C...INITIALIZING
C
INTEGER A,B,C,D,E,F
INTEGER EUCLID,AA,BB,CC,DD,EE,FF
DOUBLE DNORM,DIM,DJM,DKM,DN
15 COMMON/D1/IA,IB,IC,INDEX,IDET
COMMON/D2/IL1,IL2,IL3,IT1,IT2,IT3,IP1,IP2,IP3
COMMON/D3/MF,MP,MAP(500,2)
COMMON/D4/IP(2000),IR(2000),IS(2000)
COMMON/D5/IM,JM,KM
20 COMMON/D6/EUCLID
COMMON/D11/A(3),B(3),C(3),D(3),E(3),F(3)
COMMON/D12/AA,BB,CC,DD,EE,FF,IPP,IRR,ISS
COMMON/D21/I,J,K
DIMENSION NF(100), NP(100)
25 DIMENSION IPW(100),IRW(100),ISW(100)
C
C...THE PROGRAM BEGINS
C
30 40 READ 101,IDET,INDEX,IA,IB,IC,I,J,K,L
IF(IDET.EQ.0) STOP
101 FORMAT(I6,I2,3I5,3I16,I3)
C...THE FIELD BEING CONSIDERED HAS DISCRIMINANT -IDET, AND IS K(X) WHERE
C...X**3-IA*X**2+IR*X-IC=0. THE FUNDAMENTAL UNIT OF THE FIELD IS
C...F=(I*X**2+J*X+K)/L.
35 IDETM=-IDET
PRINT 102,IDETM,INDEX,IA,IB,IC
102 FORMAT(1H1,1X,26HTHE FIELD HAS DISCRIMINANT,I8,7H, INDEX,I3,
129H, AND POLYNOMIAL COEFFICIENTS,3I6)
CALL BASE
40 C...AN INTEGRAL BASIS FOR THE FIELD IS (1,X,Y) WHERE
C... X**2=IT1*Y + IT2*X + IT3
C... Y**2=IL1*Y + IL2*X + IL3
C... X*Y =IP1*Y + IP2*X + IP3
C...AND SO Y=(X**2-IT2*X-IT3)/IT1, IT1=INDEX.
45 PRINT 103
103 FORMAT(/,1X,45HTHE FIELD IS K(X) AND HAS BASIS (Y,X,1) WHERE)
ASSIG1=4HX**2 E ASSIG2=4HY**2 E ASSIG3=4HX*Y
PRINT 104,ASSIG1,IT1,IT2,IT3
PRINT 104,ASSIG2,IL1,IL2,IL3
50 PRINT 104,ASSIG3,IP1,IP2,IP3

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104 FORMAT(12X,A4,2H =,I10,3HY +,I10,3HX +,I10)
      CALL CHANGE(I,J,K,L,I1,J1,K1)
C...E=(I*X**2+J*X+K)/L=I1*Y+J1*X+K1
      I=I1 & J=J1 & K=K1
55 C...E=I*Y+J*X+K
      108 FORMAT(1X,26HTHE UNIT OF THE FIELD IS (,I16,1H,,I16,1H,,I16,1H))
      PRINT 108,I,J,K
C...THE 'REDUCED' FORM OF THE UNIT IDEAL IS NOW CALCULATED.
      AA=BB=DD=1 & CC=EE=FF=0
60      CALL PRIN(1)
      EUCLID=0 & ICHECK=-1
      MAPC=1
C...EUCLID=0 WHEN THE EUCLIDEAN PROPERTY OF THE FIELD IS UNKNOWN AND IS
C...EQUAL TO 1 WHEN THE FIELD IS FOUND TO BE NON-EUCLIDEAN.
65 C...MAPC IS A POINTER TO THE NEXT 'EMPTY' ELEMENT OF (IP,IR,IS).
C...ICHECK=-1 WHEN CONSIDERING E-1 AND ICHECK=1 WHEN CONSIDERING E+1
C...WE NOW BEGIN A CONSIDERATION OF E-1.
      MAX=0
      IM=I & JM=J & KM=K-1
70 C...IM*Y+JM*X+KM=E+ICHECK.
      31 DIM=IM & DJM=JM & DKM=KM
      DN=DNORM(DIM,DJM,DKM)
      N=DN
      DN=DARS(DN)
75      IF (DN.LT.1.0D14) GO TO 956
      PRINT 957,ICHECK
      957 FORMAT(1X,10(1H*),10HNORM OF E+,I2,24HIS TOO LARGE FOR INTEGER,
      110(1H*))
      GO TO 20
80      956 PRINT 107,ICHECK,N
      107 FORMAT(/,1X,36HNOW CONSIDER THE FACTORIZATION OF E+,I2,15H WHICH H
      1AS NORM,I16)
      N=TARS(N)
      EN=FLOAT(N) & REN=SQRT(EN) & MAXR=IFIX(REN)+1 & MAXF=MAXR/2+1
85      MAXF=MIN0(100,MAXF)
C...MAXF IS AN UPPER BOUND FOR THE NUMBER OF DISTINCT PRIME FACTORS
C...OF N.
      CALL FACTOR(N,NF,NP,MAXF,NFC)
90 C...NFC IS THE EXACT NUMBER OF DISTINCT PRIME FACTORS OF N. NP(I) IS THE
C...POWER TO WHICH NF(I) IS A FACTOR OF N.
C...NOW FIND MAX WHICH IS THE LARGEST NUMBER OF
C...(F:F=R OR R**3 WHERE R IS A FACTOR OF N).
      IZFR0=0
95      IF (ICHECK.EQ.-1) GO TO 15
      IF (MAX.NE.MBX) GO TO 17
      15 MBX=1
      DO 2 I1=1,NFC
      IF (NP(I1).LT.3) GO TO 21
      MRX1=NF(I1)**3
100      IF (MRX1.LT.500) GO TO 22

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21 MBXX1=NF(II)
   IF(MBXX1.GT.500) GO TO 2
22 MBX=MAX0(MBX,MBXX1)
   2 CONTINUE
105   IF(MBX.GT.1) GO TO 1
      PRINT 105,ICHECK
105  FORMAT(1X,2HE+,I2,29H HAS NO FACTORS LESS THAN 500)
      MAX=MBX
      GO TO 20
110   1 IF(MBX.LE.MAX) GO TO 17
      MLX=MAX0(2,MAX+1)
      MAX=MBX
      PRINT 109,MLX,MAX
115  109 FORMAT(1X,49H ALL INTEGERS WITH NORM OF ABSOLUTE VALUE BETWEEN,IS,
      14H AND,I5,25H INCLUSIVE ARE CALCULATED)
      C...ALL INTEGERS WITH NORM BETWEEN MLX AND MAX ARE NOW CALCULATED.
      DO 3 INORM=MLX,MAX
         IF(INORM.EQ.2) GO TO 8
         FNORM=FLOAT(INORM) & RNORM=SQRT(FNORM) & NNORM=IFIX(RNORM)
120         DO 7 INM=2,NNORM
            IF(MOD(INORM,INM).EQ.0) GO TO 9
            7 CONTINUE
            GO TO 8
            9 MAP(INORM,1)=MAPC & NNORM=INORM/2
125 C.....HAVING REACHED THIS POINT INORM IS COMPOSITE SO FIND THE ALGEBRAIC
      C.....INTEGERS OF NORM INORM BY CONSIDERING THE PRIME FACTORS OF INORM.
      C.....MAP(INORM,1) IS THE FIRST POSITION IN THE ARRAYS (IP,IR,IS) IN
      C.....WHICH AN INTEGER OF NORM INORM APPEARS. NNORM IS THE GREATEST
      C.....POSSIBLE PROPER RATIONAL INTEGRAL FACTOR OF INORM.
130      IP(MAPC)=0 & IR(MAPC)=0 & IS(MAPC)=1
      C.....SET THIS FIRST POSITION INITIALLY TO 1 SO THAT AS PRIME FACTORS OF
      C.....INORM ARE FOUND, THE ALGEBRAIC INTEGERS OF NORM THESE PRIMES
      C.....RAISED TO THE POWER TO WHICH THEY DIVIDE INORM CAN BE MULTIPLIED
      C.....TOGETHER SO THAT ALL THE INTEGERS OF NORM INORM ARE FOUND.
135      MAPC=MAPC+1 & IINORM=INORM
      C...IINORM IS SET TO INORM HERE SO THAT AS PRIME FACTORS OF INORM ARE
      C...FOUND IINORM MAY BE DIVIDED BY A SUITABLE POWER OF THAT PRIME WHILE
      C...SAVING INORM.
      DO 50 INM=2,NNORM
140 C.....LOOK FOR A FACTOR OF IINORM I.E. A PRIME FACTOR OF INORM.
      IF(MOD(IINORM,INM).NE.0) GO TO 50
      C.....HAVING FOUND A PRIME FACTOR, FIND TO WHAT POWER, INPW, IT DIVIDES
      C.....IINORM.
      INPW=0
145      81 IF(MOD(IINORM,INM).NE.0) GO TO 80
         IINORM=IINORM/INM & INPW=INPW+1
         GO TO 81
      80 IF(MAP(INM,1).EQ.-1) GO TO 51
         ISM=MAP(INM,2)-MAP(INM,1)+1
150         IF(ISM.EQ.1) GO TO 52

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      IF(ISM.EQ.2) GO TO 53
C.....HAVING REACHED THIS POINT INM HAS THREE DISTINCT ALGEBRAIC FACTORS
C.....SO ALL ALGEBRAIC INTEGERS OF NORM INM**INPW ARE PUT IN THE ARRAYS
C.....(IPW,IRW,ISW).
155      IS1=MAP(INM,1) & IS2=IS1+1 & IS3=IS2+1
      IPS1=IP(IS1) & IRS1=IR(IS1) & ISS1=IS(IS1)
      IPS2=IP(IS2) & IRS2=IR(IS2) & ISS2=IS(IS2)
      IPS3=IP(IS3) & IRS3=IR(IS3) & ISS3=IS(IS3)
      INPWP=INPW+1 & IWC=0
160      DO 82 JMAP1P=1,INPWP
      JMAP1=JMAP1P-1
      LIMIT=INPWP-JMAP1
      DO 82 JMAP2P=1,LIMIT
      JMAP2=JMAP2P-1
165      JMAP3=INPW-(JMAP1+JMAP2) & IWC=IWC+1
      IF(IWC.GT.100) GO TO 97
      IPWP=0 & IRWP=0 & ISWP=1
      IF(JMAP1.EQ.0) GO TO 84
      DO 83 JPW=1,JMAP1
170      CALL MULTC(IPWP,IRWP,ISWP,IPS1,IRS1,ISS1,IPWW,IRWW,ISWW)
      IPWP=IPWW & IRWP=IRWW & ISWP=ISWW
      83 CONTINUE
      84 IF(JMAP2.EQ.0) GO TO 86
      DO 85 JPW=1,JMAP2
175      CALL MULTC(IPWP,IRWP,ISWP,IPS2,IRS2,ISS2,IPWW,IRWW,ISWW)
      IPWP=IPWW & IRWP=IRWW & ISWP=ISWW
      85 CONTINUE
      86 IF(JMAP3.EQ.0) GO TO 89
      DO 87 JPW=1,JMAP3
180      CALL MULTC(IPWP,IRWP,ISWP,IPS3,IRS3,ISS3,IPWW,IRWW,ISWW)
      IPWP=IPWW & IRWP=IRWW & ISWP=ISWW
      87 CONTINUE
      89 IPW(IWC)=IPWP & IRW(IWC)=IRWP & ISW(IWC)=ISWP
      82 CONTINUE
185      C...THERE ARE IWC INTEGERS OF NORM INM**INPW CONTAINED IN (IPW,IRW,ISW).
      GO TO 99
      53 IS1=MAP(INM,1) & IS2=IS1+1
      IPS1=IP(IS1) & IRS1=IR(IS1) & ISS1=IS(IS1)
      IPS2=IP(IS2) & IRS2=IR(IS2) & ISS2=IS(IS2)
190      C.....HAVING REACHED THIS POINT INM HAS TWO DISTINCT ALGEBRAIC FACTORS.
      INPWP=INPW+1 & IWC=0
      DO 92 JMAP1P=1,INPWP
      JMAP1=JMAP1P-1
      JMAP2=INPW-JMAP1 & IWC=IWC+1
195      IF(IWC.GT.100) GO TO 97
      IPWP=0 & IRWP=0 & ISWP=1
      IF(JMAP1.EQ.0) GO TO 94
      DO 93 JPW=1,JMAP1
      CALL MULTC(IPWP,IRWP,ISWP,IPS1,IRS1,ISS1,IPWW,IRWW,ISWW)
200      IPWP=IPWW & IRWP=IRWW & ISWP=ISWW

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93 CONTINUE
94 IF(JMAP2.EQ.0) GO TO 91
DO 95 JPW=1,JMAP2
CALL MULTC(IPWP,IRWP,ISWP,IPS2,IRS2,ISS2,IPWW,IRWW,ISWW)
205 IPWP=IPWW & IRWP=IRWW & ISWP=ISWW
95 CONTINUE
91 IPW(IWC)=IPWP & IRW(IWC)=IRWP & ISW(IWC)=ISWP
92 CONTINUE
C.....THERE ARE NOW IWC INTEGERS WITH NORM INM**INPW.
210 GO TO 99
52 IMAP=MAP(INM,1)
IPS1=IP(IMAP) & IRS1=IR(IMAP) & ISS1=IS(IMAP)
C.....HAVING REACHED THIS POINT INM HAS JUST ONE LINEAR FACTOR.
IF(INPW.EQ.1) GO TO 55
215 IF(MOD(IDET,INM).EQ.0) GO TO 55
C.....INM DIVIDES THE DISCRIMINANT OF THE FIELD IF AND ONLY IF INM HAS A
C.....SQUARE FACTOR, I.E. IN THIS INSTANCE IF INM IS THE CUBE OF AN
C.....ALGEBRAIC INTEGER, HENCE IF INM DOES NOT DIVIDE IDET THEN INM HAS
C.....ONE LINEAR AND ONE QUADRATIC FACTOR.
220 CALL MULTC(IPS1,IRS1,ISS1,IPS1,IRS1,ISS1,IPL2,IRL2,ISL2)
CALL DIVC(0,0,INM,IPS1,IRS1,ISS1,IPQ,IRQ,ISQ)
C.....(IPL2,IRL2,ISL2) AND (IPQ,IRQ,ISQ) ARE THE TWO DISTINCT ALGEBRAIC
C.....INTEGERS OF NORM INM**2.
IF(MOD(INPW,2).EQ.0) GO TO 56
225 INPW2=(INPW-1)/2 & INPW=1 & GO TO 57
56 INPW2=INPW/2 & INPW=0
57 INPW2P=INPW2+1 & IWC=0
DO 64 JMAP1P=1,INPW2P
JMAP1=JMAP1P-1
230 JMAP2=INPW2-JMAP1
IWC=IWC+1
IF(IWC.GT.100) GO TO 97
IPWP=0 & IRWP=0 & ISWP=1
IF(JMAP1.EQ.0) GO TO 65
235 DO 66 JPW=1,JMAP1
CALL MULTC(IPWP,IRWP,ISWP,IPL2,IRL2,ISL2,IPWW,IRWW,ISWW)
IPWP=IPWW & IRWP=IRWW & ISWP=ISWW
66 CONTINUE
65 IF(JMAP2.EQ.0) GO TO 63
240 DO 68 JPW=1,JMAP2
CALL MULTC(IPWP,IRWP,ISWP,IPQ,IRQ,ISQ,IPWW,IRWW,ISWW)
IPWP=IPWW & IRWP=IRWW & ISWP=ISWW
68 CONTINUE
63 IPW(IWC)=IPWP & IRW(IWC)=IRWP & ISW(IWC)=ISWP
245 64 CONTINUE
C.....IWC IS THE NUMBER OF INTEGERS WITH NORM INM**PW WHERE PW=INPW IF
C.....INPW WAS EVEN AND INPW-1 IF INPW WAS ODD.
IF(INPW.EQ.0) GO TO 99
250 GO TO 67
55 IWC=1 & IPW(IWC)=0 & IRW(IWC)=0 & ISW(IWC)=1

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C.....55 IS REACHED EITHER IF INPW WAS ORIGINALLY 1 OR IF INM HAS NO
C.....QUADRATIC FACTOR . IN EITHER OF THESE CASES IF PA IS THE ALGEBRAIC
C.....INTEGER OF NORM INM THEN PA**INPW IS THE ONLY INTEGER OF NORM
C.....INM**INPW.
255 DO 62 JWC=1,IWC
      IPWP=IPW(JWC) & IRWP=IRW(JWC) & ISWP=ISW(JWC)
      DO 69 JPW=1,INPW
        CALL MULTC(IPWP,IRWP,ISWP,IPS1,IRS1,ISS1,IPWW,IRWW,ISWW)
        IPWP=IPWW & IRWP=IRWW & ISWP=ISWW
260      69 CONTINUE
      IPW(JWC)=IPWP & IRW(JWC)=IRWP & ISW(JWC)=ISWP
      62 CONTINUE
C.....IWC IS THE NUMBER OF ALGEBRAIC INTEGERS OF NORM INM**INPW.
265 99 MPC=MAPC-1 & LMAP=MAP(INORM,1) & KMAP=MAPC-LMAP
      DO 88 JMAP=LMAP,MPC
        IPP=IP(JMAP) & IRR=IR(JMAP) & ISS=IS(JMAP)
        DO 88 JPW=1,IWC
          KMAP1=JMAP*(JPW-1)*KMAP
          IF(KMAP1.GT.2000) GO TO 90
270        CALL MULTC(IPP,IRR,ISS,IPW(JPW),IRW(JPW),ISW(JPW),IP(KMAP1),
          IIR(KMAP1),IS(KMAP1))
        88 CONTINUE
        MAPC=MAPC+(IWC-1)*KMAP
        IF(MAPC.LE.2000) GO TO 96
275      90 PRINT 900,MAX
      900 FORMAT(1X,68HMORE THAN 2000 INTEGERS WITH NORM AT MOST MAX IF MAX
      1 IS GREATER THAN,IS)
      GO TO 16
      96 IF(IINORM.FO.1) GO TO 61
280      GO TO 50
      51 IF(MOD(INPW,3).NE.0) GO TO 60
C.....THIS POINT IS REACHED IF INM HAS NO ALGEBRAIC FACTORS, THUS THE
C.....ONLY POSSIBLE INTEGERS WITH NORM DIVISIBLE BY INM ARE POWERS OF
C.....INM. THUS THERE ARE ONLY INTEGERS OF NORM INM**INPW AND SO OF
285 C.....INORM IF INPW IS A MULTIPLE OF 3.
      INWP=INPW/3 & MPC=MAPC-1 & LMAP=MAP(INORM,1)
      DO 58 JMAP=LMAP,MPC
        IPU=IP(JMAP) & IRU=IR(JMAP) & ISU=IS(JMAP)
        DO 59 JPW=1,INWP
290        IPP=IPU & IRR=IRU & ISS=ISU
        CALL MULTC(IPP,IRR,ISS,0,0,INM,IPU,IRU,ISU)
        59 CONTINUE
        IP(JMAP)=IPU & IR(JMAP)=IRU & IS(JMAP)=ISU
        58 CONTINUE
295      IF(MAPC.GT.2000) GO TO 90
      IF(IINORM.EQ.1) GO TO 61
      50 CONTINUE
      60 MAPC=MAP(INORM,1) & MAP(INORM,1)=-1 & GO TO 3
C...MAP(INORM,1)=-1 INDICATES THAT THERE ARE NO INTEGERS OF NORM INORM.
300      61 MAP(INORM,2)=MAPC-1 & GO TO 3

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C...MAP(INORM,2) IS THE LAST POSITION IN THE ARRAYS (IP,IR,IS) IN WHICH
C...AN INTEGER OF NORM INORM APPEARS.
97 PRINT 901,INORM
305 901 FORMAT(1X,7I)THERE ARE MORE THAN 100 INTEGERS OF NORM P**A WHERE P
1**A IS A FACTOR OF ,IS)
GO TO 16
R CALL IDEAL(INORM,ISC)
C...THIS POINT IS REACHED WHEN INORM IS A RATIONAL PRIME, ITS ALGEBRAIC
C...FACTORS ARE NOW CALCULATED. THERE ARE ISC SUCH FACTORS WHERE ISC IS
310 C...AT MOST 3.
IF(ISC.NE.0) GO TO 70
C...IF INORM IS ALSO AN ALGEBRAIC PRIME THEN THERE ARE NO INTEGERS OF
C...NORM INORM.
MAP(INORM,1)=-1 & GO TO 3
315 70 MAP(INORM,1)=MAPC
DO 10 IFILL=1,ISC
AA=A(IFILL) & BB=B(IFILL) & CC=C(IFILL)
DD=D(IFILL) & EE=E(IFILL) & FF=F(IFILL)
CALL PRIN(INORM)
320 IP(MAPC)=IPP & IR(MAPC)=IRR & IS(MAPC)=ISS
MAPC=MAPC+1
IF(MAPC.LE.2000) GO TO 10
IF(IFILL.EQ.ISC) IZERO=1
PRINT 900,MAX
325 GO TO 16
10 CONTINUE
MAP(INORM,2) =MAPC-1
3 CONTINUE
GO TO 17
330 16 MAX=INORM-1+IZERO
17 DO 6 II=1,NFC
C...RUN THROUGH THE POSSIBLE NORM VALUES FOR FACTORS OF E+ICHECK-1
MP=NP(II) & MF=NF(II)
IF(MF.LE.3) GO TO 23
335 IF(MF.GT.MAX) GO TO 6
IF(MAP(MF,1).EQ.-1) GO TO 23
CALL TEST
IF(EUCLID.EQ.1) GO TO 40
C...IF THE EUCLIDEAN PROPERTY OF THE FIELD IS STILL UNKNOWN FIND IF
340 C...MF**3 IS A FACTOR OF NORM(IM*Y+JM*X+KM) OTHERWISE STOP.
23 IF(MP.LT.3) GO TO 6
C...MF**3 DOES NOT DIVIDE NORM(IM*Y+JM*X+KM) GO ON TO NEXT POSSIBLE
C...PRIME NORM VALUE.
MF=MF**3 & MP=1
345 IF(MF.GT.MAX) GO TO 6
CALL TEST
IF(EUCLID.EQ.1) GO TO 40
6 CONTINUE
C...NOW AS POSSIBLE NORM VALUES OF THE MODULI TO BE CONSIDERED WE HAVE
350 C...GONE THROUGH ALL THE PRIME FACTORS OF IM*Y+JM*X+KM AND THEIR CUBES.

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C...HAVING REACHED NO CONCLUSION WE GO ON TO A CONSIDERATION OF THE  
 C...FACTORS OF E+1 IF  $IM*Y+JM*X+KM=E-1$  AND END OTHERWISE.

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20 IF(ICHECK.EQ.1) GO TO 30
355 ICHECK=1 E IM=I E JM=J E KM=K+1
GO TO 31
30 PRINT 106
106 FORMAT(1X,30HNO CONCLUSION HAS BEEN REACHED)
GO TO 40
END

10 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
31 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
32 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
33 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
34 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
35 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
36 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
37 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
38 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
39 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
40 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
41 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
42 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
43 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
44 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
45 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
46 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
47 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
48 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
49 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
50 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
51 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
52 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
53 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
54 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
55 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
56 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
57 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
58 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
59 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
60 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
61 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
62 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
63 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
64 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
65 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
66 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
67 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
68 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
69 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
70 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
71 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
72 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
73 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
74 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
75 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
76 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
77 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
78 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
79 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
80 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
81 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
82 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
83 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
84 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
85 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
86 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
87 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
88 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
89 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
90 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
91 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
92 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
93 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
94 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
95 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
96 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
97 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
98 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
99 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1
100 30=31*(E+1)
IM=I+1 E JM=J+1 KM=K+1

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SUBROUTINE BASE
DOURLE IURM,IORD
COMMON/D1/IA,IB,IC,INDEX,IDET
COMMON/D2/IL1,IL2,IL3,IT1,IT2,IT3,IP1,IP2,IP3
COMMON/D10/H,ZZ,HHH/D9/IO,INN,0,V2
COMMON/D15/IURN,I0B,I0BDS/D16/IURM,IORD
C...WHEN CONSIDERING THE FIELD K(X) SUPPOSE H=3*X-IA THEN H SATISFIES AN
C...EQUATION H**3-I0*H-INN=0. FIND H AND H**2-I0 FOR USE IN THE
C...SUBROUTINE PRIN.
10      I0=3*IA*IA-9*IB
      INN=27*IC-9*IA*IB+2*IA*IA*IA
      O=FLOAT(I0)
      V1=O/3.0
      V2=FLOAT(INN)
15      V3=SQRT(V2**2-4.0*V1**3)
      V4=(V2+V3)/2.0
      V5=(V2-V3)/2.0
      IF(V4.NE.0.0) GO TO 51
      V6=0.0 & GO TO 52
20      51 D=1.0
      IF(V4.GE.0.0) GO TO 53 TO OVERFLOW
      V4=-V4 & D=-1.0
      V6=EXP((ALOG(V4))/3.0)*D
25      52 IF(V5.NE.0.0) GO TO 54
      V7=0.0 & GO TO 55
      54 D=1.0
      IF(V5.GE.0.0) GO TO 56 FOR THE NEGATIVE OF D.
      V5=-V5 & D=-1.0
30      56 V7=EXP((ALOG(V5))/3.0)*D
      55 H=V6+V7 & HHH=H*H
      ZZ=HHH-0
C...NOW FIND IT AND IS SUCH THAT (1,X,Y), WHERE Y=(X*X+IT*X+IS)/INDEX,
C...IS A BASIS FOR THE FIELD.
35      IN=INDEX
      IN2=IN**2 & IN3=IN2*IN
      K1=IA*IA-2*IB & K2=IA*IC & K3=IB*IC
      K4=IB*IB-2*IA*IC & K5=IA*IB-3*IC & IC2=IC*IC
      IF(INDEX.NE.1) GO TO 1
      IT=0 & IS=0 & GO TO 2
40      1 DO 3 ITC=1,IN
      IT=ITC-1
      DO 4 ISC=1,IN
      IS=ISC-1
      JA=K1+IT*IA+3*IS
45      IF(MOD(JA,IN).NE.0) GO TO 4
      JB=K4+IT**2*IB+3*IS**2+IT*K5+2*IS*K1+2*IS*IT*IA
      IF(MOD(JB,IN2).NE.0) GO TO 4
      JC=IS**3+IT**3*IC+IC**2+IS**2*IT*IA+IS*K4+IS*IT**2*IB+IT*IB*IC
      1+IT**2*IA*IC+IS**2*K1+IS*IT*K5
50      IF(MOD(JC,IN3).NE.0) GO TO 4

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```

        GO TO 2
    4 CONTINUE
    3 CONTINUE
C...NOW FIND THE COEFFICIENTS OF Y,X AND 1 IN THE EXPRESSIONS FOR Y*Y,
55 C...X*X AND X*Y...
    C...  Y*Y=IL1*Y+IL2*X+IL3
    C...  X*X=IT1*Y+IT2*X+IT3
    C...  X*Y=IP1*Y+IP2*X+IP3
    2 IT1=INDEX
60     IT2=-IT
        IT3=-IS
        ML=2*IA*IT+IA*IA+2*IS*IT*IT-IB
        IL1=ML/IN
        IL2=(-IT*ML+2*IS*IT+IC-2*IB*IT-IA*IB)/IN2
65     IL3=(-IS*ML+2*IC*IT+IA*IC+IS*IS)/IN2
        MP=IA+IT
        IP1=MP
        IP2=(IS-IB-IT*MP)/IN
        IP3=(IC-IS*MP)/IN
70 C...FIND AN UPPER BOUND, A POWER OF 10, FOR THE ARGUMENTS OF DNORM
    C...WHICH WILL NOT GIVE RISE TO OVERFLOW.
        K1P=IABS(K1) & K4P=IABS(K4) & K5P=IABS(K5)
        NF=MAX0(K1P,K2,K3,K4P,K5P,IC2)
        IONF=IOF(NF)
75     IONS=(29-IONF)/3
        IURN=10**IONS
    C...NOW FIND AN UPPER BOUND FOR THE ARGUMENTS OF MULT2.
        IOA=IABS(IO) & INNA=IABS(INN)
        NON=MAX0(IOA,INNA)
80     ION=IOF(NON)
        IOMS=29-ION
        IURM=10.000**IOMS
    C...FIND SIMILAR BOUNDS FOR THE ARGUMENTS OF MULTC, DIVC AND DIVCD.
85     IL1A=IABS(IL1) & IL2A=IABS(IL2) & IL3A=IABS(IL3)
        IT1A=IABS(IT1) & IT2A=IABS(IT2) & IT3A=IABS(IT3)
        IP1A=IABS(IP1) & IP2A=IABS(IP2) & IP3A=IABS(IP3)
        IOBL=MAX0(IL1A,IL2A,IL3A)
        IOBT=MAX0(IT1A,IT2A,IT3A)
        IOBP=MAX0(IP1A,IP2A,IP3A)
90     IOBAS=MAX0(IOBL,IOBT,IOBP)
        IOBSS=IOF(IOBAS)
        IOBS=14-IOBSS
        IOB=10**IOBS
        IF (IOBL.GT.IOBP) GO TO 20
        IF (IOBP.GT.IOBT) GO TO 21
95     IOBAS=IOBP*IOBT
        GO TO 22
    21 IOBAS=IOBP*IOBP
        GO TO 22
100    20 IF (IOBP.GT.IOBT) GO TO 23

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SUBROUTINE FACTOR(NN,NF,NP,NL,I)
C...THE SUBROUTINE FACTOR FINDS THE FACTORS OF NN WHICH ARE LESS THAN
C...500 AND PUTS THEM IN THE ARRAY NF. NP(I) IS THE POWER TO WHICH NF(I)
C...IS A FACTOR OF NN. NL IS AN UPPER BOUND ON THE NUMBER OF POSSIBLE
5 C...DISTINCT PRIME FACTORS OF NN. FROM THE MAIN PROGRAM NL IS AT MOST
C...100.
DIMENSION NF(NL),NP(NL)
N=NN
EN=FLOAT(N) & RN=SQRT(EN) & NR=IFIX(RN)
10 C...NR IS THE GREATEST INTEGER LESS THAN SQRT(N) AND SO N CAN HAVE AT
C...MOST ONE FACTOR GREATER THAN NR.
I=0 & JF=1
3 JF=JF+1 & JP=0
IF(JF.GT.NR) GO TO 6
IF(JF.GT.500) GO TO 7
15 1 IF(MOD(N,JF).NE.0) GO TO 4
JP=JP+1 & N=N/JF
EN=FLOAT(N) & RN=SQRT(EN) & NR=IFIX(RN)
6 IF(N.NE.1) GO TO 1
20 GO TO 5
4 IF(JP.EQ.0) GO TO 3
I=I+1 & NF(I)=JF & NP(I)=JP
IF(I.GE.NL) GO TO 7
GO TO 3
25 6 JF=N & JP=1
5 I=I+1 & NF(I)=JF & NP(I)=JP
7 RETURN
END

```

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SUBROUTINE IDEAL(N,ISC)
INTEGER A,B,C,D,E,G
C...IDEAL FINDS THE IDEALS (A(I),B(I)*X+C(I),D(I)*Y+E(I)*X+G(I))
C...FOR I=1,ISC OF NORM N WHERE N IS A RATIONAL PRIME, HENCE
5 C...ISC=0,1,2, OR 3.
COMMON/D1/IA,IB,IC,L,IDET
COMMON/D2/KX,KY,KZ,IX,IY,IZ,JX,JY,JZ
COMMON/D11/A(3),B(3),C(3),D(3),E(3),G(3)
DIMENSION MY(6,3)
10 C...NORL(KR)=NORM(X**2+KT*X+KB)/L
NORL(KR)=(IC*IC+IC*KT*KT*KT+KB*KB*KB+IB*IC*KT+(IB*IB-2*IA*IC)*KB
1+(IA*IC+IB*KB)*KT*KT+((IA*IA-2*IB)+IA*KT)*KB*KB+(IA*IB-3*IC)*KT
2*KR)/(L*L*L)
C...NORN(J)=NORM(X+J)
15 NORN(J)=IC+J*(IB+J*(IA+J))
INDEX=L
KT=-IY & KS=-IZ

C**** SECTION 1 ****
20 ISC=0
C.....LOOP FOR VALUES OF J I.E. C(ISC).....
DO 4 JJ=1,N
J=JJ-1
25 C...N MUST DIVIDE NORM(X+J)
IF(MOD(NORN(J),N).NE.0) GO TO 4
C.....LOOP FOR VALUES OF K I.E. G(ISC).....
DO 6 KK=1,N
K=KK-1
30 KB=K*L+KS
C.....N MUST DIVIDE NORM(Y+K).....
IF(MOD(NORL(KR),N).NE.0) GO TO 6

C**** SECTION 2 ****
35 C..... MULTIPLES OF IDEAL MEMBERS BY INTEGERS MUST BE IDEAL MEMBERS ..
C.....MY(IT,1)*Y+MY(IT,2)*X+MY(IT,3) IS THE PRODUCT OF ONE OF (X,Y)
C.....WITH ONE OF THE BASIS ELEMENTS OF THE IDEAL. IT VARIES FROM 1 TO
C.....6 SO THAT EVERY CASE IS CONSIDERED.
40 MY(1,1)=0 & MY(1,2)=N & MY(1,3)=0 & MY(2,1)=N & MY(2,2)=0
MY(2,3)=0 & MY(3,1)=IX & MY(3,2)=J+IY & MY(3,3)=IZ & MY(4,1)=J+JX
MY(4,2)=JY & MY(4,3)=JZ & MY(5,1)=JX & MY(5,2)=JY+K & MY(5,3)=JZ
MY(6,1)=KX+K & MY(6,2)=KY & MY(6,3)=KZ
DO 8 IT=1,6
45 C...CHECKS THAT MY(IT,1)*Y+MY(IT,2)*X+MY(IT,3) BELONGS TO THE IDEAL.
NA=MY(IT,1) & NB=MY(IT,2) & NC=MY(IT,3)
NC=NC-NA*K-NB*J
IF(MOD(NC,N).NE.0) GO TO 6
8 CONTINUE
50 ISC=ISC+1

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SUBROUTINE PRIN(INK)
DOUBLE X,Y,Z,DA,DB,DC,DD,DE,DF,DG,DH,DI,Y2
DOUBLE IUBM,I0BD
DOUBLE DAXMUL,DXMUL1,DXMUL2
5  DOUBLE DDA,ddb,DDC
COMMON/D1/IA,IB,IC,INDEX,IDET
COMMON/D2/KX,KY,KZ,IX,IY,IZ,JX,JY,JZ
COMMON/D7/M1,M2,M3,N1,N2,N3,A,B,C/D8/MB(3)
COMMON/D10/H,ZZ,HHH/D9/IQ,IN,Q,RN
10 COMMON/D12/JA,JB,JC,JD,JE,JF,JP,JR,JS
COMMON/D16/IUBM,I0BD
COMMON/D20/MU1,MU2,MU3,NU1,NU2,NU3,IGU1
DIMENSION LI(8),IAN(3,99),JAN(3,99),NP(99),K(4,3)
C...RHO IS THE SQUARE OF THE DISTANCE OF THE POINT U+V*H+W*H**2 FROM
15 C...THE REAL AXIS.
RHO(U,V,W)=((U+W*Q)**2-V*(V*Q+W*RN))+(W*W*RN-V*U)*H
1+(V*V-W*(U+W*Q))*HHH
N=1
20 C**** SECTION 1 ****
LI(1)=JA & LI(4)=0 & LI(3)=JB*INDEX & LI(2)=JC*INDEX
LI(7)=JD & LI(6)=JF*INDEX-IY*LI(7)
LI(5)=JF*INDEX-IZ*LI(7) & LI(8)=INDEX
25 C...BASIS FOR LATTICE IS [1,(M1+M2*H+M3*H**2)/IG,(N1+N2*H+N3*H**2)/IG]
IG=9*LI(1)*LI(8) & M1=9*LI(2)+3*IA*LI(3)+LI(4)*IA*IA
N1=9*LI(5)+3*IA*LI(6)+LI(7)*IA*IA
M2=3*LI(3)+2*LI(4)*IA
N2=3*LI(6)+2*LI(7)*IA
30 M3=LI(4) & N3=LI(7)
C..... A,B,C ARE COEFFICIENTS OF THE ASSOCIATED BINARY QUADRATIC FORM
A=R=C=0.
C...PARTS (I),... (VIII) ARE THE STEPS OF VORONOI'S ALGORITHM AS
C...DESCRIBED IN DELONE AND FADDEEV.
35 C... ( I ) .....
C..... TEST THAT M2*N3-M3*N2 IS GREATER THAN 0 .....
10 IF(M2*N3-M3*N2.GT.0) GO TO 20
I1=M1 & I2=M2 & I3= M3
M1=N1 & M2=N2 & M3= N3
40 N1=I1 & N2=I2 & N3=I3
C... ( II ) .....
C..... PRODUCTION OF INITIAL VALUES OF A,B,C .....
20 R1=M1 & R2=M2 & R3=M3
S1=N1 & S2=N2 & S3=N3
G=IG
45 A=R2*R2+R2*R3*H+R3*R3*ZZ
B=R2*S2+(R2*S3+R3*S2)*H*0.5+R3*S3*ZZ
C=S2*S2+S2*S3*H+S3*S3*ZZ
C... ( III ) .....
50 C..... (III) AND (IV) FIND TWO BASIS ELEMENTS OF THE REDUCED HEXAGON OF

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C...ZELLING.
IF (R.GT.0.) GO TO 22
CALL SUB(0,1,-1,0)
55 C.... ( IV ) .....
22 IF (A-R.LT.0.) GO TO 32
IF (C-R.GT.0.) GO TO 50
32 IF (A-C.GT.0.) GO TO 40
ID=INT(B/A)
CALL SUB(1,-ID,0,1)
60 GO TO 22
40 ID=INT(B/C)
CALL SUB(1,0,-ID,1)
GO TO 22
C.... ( V ) .....
65 C...TO FIND THE TWO ZELLING HEXAGON BASIS ELEMENTS WHICH COVER THE
C...NEGATIVE 'XSI' AXIS.
50 R2=M2 & R3=M3 & S2=N2 & S3=N3
RR=R2-R3*H & RD=S2-S3*H
RUP=RD-RR
70 IF (RR.LT.0.) GO TO 60
IF (RD.LT.0.) GO TO 90
IF (RUP.GT.0.) GO TO 55
CALL SUB(0,-1,1,1)
GO TO 90
75 55 CALL SUB(-1,-1,1,0)
GO TO 90
60 IF (RD.LT.0.) GO TO 70
CALL SUB(-1,0,0,-1)
GO TO 90
80 70 IF (RD-RR.GT.0.) GO TO 80
CALL SUB(1,1,-1,0)
GO TO 90
80 CALL SUB(0,1,-1,-1)
C.... ( VI ) .....
85 C...FINDING THE PINHEADS CORRESPONDING TO (M1+M2*H+M3*H*H2)/IG AND
C...(N1+N2*H+N3*H*H2)/IG.
90 R1=M1 & R2=M2 & R3=M3
S1=N1 & S2=N2 & S3=N3
T=(R1+R2*H+R3*HHH)/G
90 I=IFIX(T)
IF (T.GT.0.) GO TO 4001
I=I-1
4001 M1=M1-I*IG
S=(S1+S2*H+S3*HHH)/G
95 I=IFIX(S)
IF (S.GT.0.) GO TO 4002
I=I-1
4002 N1=N1-I*IG
C.... ( VII ) .....
100 C...CHOOSING THE TWO ELEMENTS OF THE REDUCED BASIS OF THE LATTICE FROM

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C...THE SEVEN POSSIBILITIES BY FINDING THE TWO WITH MINIMUM RHO VALUE.
200 R1=M1 & S1=N1 & G2=2.*G
XA=(2.*(R1+R3*Q)-R2*H-R3*HHH)/G2
XC=(2.*(S1+S3*Q)-S2*H-S3*HHH)/G2
105 IF(XA-0.5.GT.0.) GO TO 201
K(1,1)=M1 & K(1,2)=M2 & K(1,3)=M3
GO TO 202
201 K(1,1)=I6-M1 & K(1,2)= -M2 & K(1,3)= -M3
202 IF(XC-0.5.GT.0.) GO TO 203
110 K(2,1)=N1 & K(2,2)=N2 & K(2,3)=N3
GO TO 204
203 K(2,1)=I6-N1 & K(2,2)= -N2 & K(2,3)= -N3
204 T=(R1+R2*H+R3*HHH)/G
S=(S1+S2*H+S3*HHH)/G
115 IF(T-S.LT.0.) GO TO 220
IF(XA-XC.GT.0.5) GO TO 210
K(3,1)=M1-N1 & K(3,2)=M2-N2 & K(3,3)=M3-N3
GO TO 240
210 K(3,1)=I6-M1+N1 & K(3,2)= -M2+N2 & K(3,3)= -M3+N3
120 GO TO 240
220 IF(XC-XA.GT.0.5) GO TO 230
K(3,1)=N1-M1 & K(3,2)=N2-M2 & K(3,3)=N3-M3
GO TO 240
230 K(3,1)=I6-N1+M1 & K(3,2)= -N2+M2 & K(3,3)= -N3+M3
125 240 R70=RHO(FLOAT(K(1,1)),FLOAT(K(1,2)),FLOAT(K(1,3)))
R71=RHO(FLOAT(K(2,1)),FLOAT(K(2,2)),FLOAT(K(2,3)))
R72=RHO(FLOAT(K(3,1)),FLOAT(K(3,2)),FLOAT(K(3,3)))
IF(R72.LT.R70) GO TO 280
IF(R72.LT.R71) GO TO 260
130 IF(T+S.GT.1.)GO TO 250
IF(XA+XC.GT.0.5) GO TO 250
K(4,1)=M1+N1 & K(4,2)=M2+N2 & K(4,3)=M3+N3
RZ3=RHO(FLOAT(K(4,1)),FLOAT(K(4,2)),FLOAT(K(4,3)))
IF(RZ3.GT.R70) GO TO 245
IF(RZ3.GT.RZ1) GO TO 244
135 M1=K(4,1) & M2=K(4,2) & M3=K(4,3)
IF(RZ0.GT.RZ1) GO TO 243
N1=K(1,1) & N2=K(1,2) & N3=K(1,3)
GO TO 300
140 243 N1=K(2,1) & N2=K(2,2) & N3=K(2,3)
GO TO 300
244 M1=K(2,1) & M2=K(2,2) & M3=K(2,3)
N1=K(4,1) & N2=K(4,2) & N3=K(4,3)
GO TO 300
145 245 IF(RZ3.GT.RZ1) GO TO 250
M1=K(1,1) & M2=K(1,2) & M3=K(1,3)
N1=K(4,1) & N2=K(4,2) & N3=K(4,3)
GO TO 300
150 250 IF(RZ0.GT.RZ1) GO TO 255
M1=K(1,1) & M2=K(1,2) & M3=K(1,3)

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155 N1=K(2,1) & N2=K(2,2) & N3=K(2,3)
GO TO 300
255 M1=K(2,1) & M2=K(2,2) & M3=K(2,3)
N1=K(1,1) & N2=K(1,2) & N3=K(1,3)
GO TO 300
160 260 M1=K(1,1) & M2=K(1,2) & M3=K(1,3)
N1=K(3,1) & N2=K(3,2) & N3=K(3,3)
GO TO 300
280 IF(RZ1.GT.RZ2) GO TO 290
M1=K(2,1) & M2=K(2,2) & M3=K(2,3)
N1=K(3,1) & N2=K(3,2) & N3=K(3,3)
GO TO 300
165 290 IF(RZ0.GT.RZ1) GO TO 295
M1=K(3,1) & M2=K(3,2) & M3=K(3,3)
N1=K(1,1) & N2=K(1,2) & N3=K(1,3)
GO TO 300
295 M1=K(3,1) & M2=K(3,2) & M3=K(3,3)
N1=K(2,1) & N2=K(2,2) & N3=K(2,3)
C..... ( VIII ) .....
170 300 JAN(1,N)=N1 & JAN(2,N)=N2 & JAN(3,N)=N3 & NP(N)=IG
IAN(1,N)=M1 & IAN(2,N)=M2 & IAN(3,N)=M3
IF(INK.EQ.1) GO TO 9
N00=N-1
C..... FIND THE INVERSE OF THE SECOND BASIS ELEMENT,THE FIRST RELATIVE
175 C...MINIMUM OF THE LATTICE.
8988 CALL INVER(M1,M2,M3,JDET)
I1=MB(1)*IG & I2=MB(2)*IG & I3=MB(3)*IG
CALL ICF(I1,I2,I3,IF1)
C..... DIVIDE LATTICE BY SECOND BASIS ELEMENT AND PRODUCE A NEW LATTICE
180 CALL MULT(MB(1),MB(2),MB(3),N1,N2,N3,J1,J2,J3)
CALL ICF(J1,J2,J3,IF2)
CALL ICF(IF1,IF2,JDET,IF)
M1=J1/IF & M2=J2/IF & M3=J3/IF
N1=I1/IF & N2=I2/IF & N3=I3/IF & IG=JDET/IF
185 IF(N.EQ.1) GO TO 501
C.....CHECK IF THE OLD LATTICE IS THE UNIT LATTICE.....
IF(IAN(1,N).NE.MU1) GO TO 501
IF(IAN(2,N).NE.MU2) GO TO 501
IF(IAN(3,N).NE.MU3) GO TO 501
190 IF(JAN(1,N).NE.NU1) GO TO 501
IF(JAN(2,N).NE.NU2) GO TO 501
IF(JAN(3,N).NE.NU3) GO TO 501
IF(NP(N).EQ.IGU1) GO TO 600
501 N=N+1
195 C.....IF N.GT.99 PRINT DIAGNOSTIC.....
IF(N.LT.99) GO TO 10
PRINT 1111
1111 FORMAT(38H LATTICE LOOP HAS MORE THAN 99 MEMBERS)
GO TO 7
200

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C*** SECTION 2 *****
C...SECTION 2 IS ENTERED WHEN THE LATTICE LOOP IS COMPLETE. IT
C...MULTIPLIES ALL THE RELATIVE MINIMA TOGETHER TO GIVE AN INTEGER OF
205 C...NORM INK.
      600 DB=DC=0 & DA=LI(1) & NPD=1
          DO 700 JIM=1,NQ0
              DD=IAN(1,JIM) & DE=IAN(2,JIM) & DF=IAN(3,JIM)
              DDA=DABS(DA) & DDB=DABS(DB) & DDC=DABS(DC)
210          DXMUL1=DMAX1(DDA,DDB,DDC)
              DDA=DABS(DD) & DDB=DABS(DE) & DDC=DABS(DF)
              DXMUL2=DMAX1(DDA,DDB,DDC)
              DAXMUL=DXMUL1*DXMUL2
215          IF(DAXMUL.LT.IUBM) GO TO 601
              PRINT 650,INK
      650 FORMAT(1X,10(1H*),40HDANGER OF OVERFLOW WITH INTEGERS OF NORM,15,
          110(1H*))
      601 CALL MULT2(DA,DB,DC,DD,DE,DF,DG,DH,DI)
          NPD=NPD*NP(JIM)
220          X=NPD
              CALL ICF2(X,DG,DH,Y)
              Y2=Y
              CALL ICF2(Y,Y2,DI,Z)
              DA=DG/Z & DB=DH/Z & DC=DI/Z
225          NPD=X/Z
      700 CONTINUE
          DD=DC*9.D0 & DE=3.D0*DB-DC*6.D0*IA
          DF=IA*IA*DC-IA*DB+DA & DG=NPD
          CALL ICF2(DG,DD,DE,DH)
230          CALL ICF2(DH,DF,DH,DI)
              DD=DD/DI & DE=DE/DI & DF=DF/DI & DG=DG/DI
              DDA=DABS(DD) & DDB=DABS(DE) & DDC=DABS(DF)
              IF(DDA.GE.1.0D14) GO TO 701
              IF(DDB.GE.1.0D14) GO TO 701
235          IF(DDC.LT.1.0D14) GO TO 702
      701 PRINT 650,INK
      702 JP1=DD & JR1=DE & JS1=DF & JD=DG
          CALL CHANGE(JP1,JR1,JS1,JD,JP,JR,JS)
          GO TO 7
240          9 MU1=M1 & MU2=M2 & MU3=M3 & NU1=N1 & NU2=N2 & NU3=N3 & IGU1=IG
          7 RETURN
          END

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SUBROUTINE MULT2(A,B,C,D,E,F,G,H,I)
C
C DOUBLE PRECISION VERSION OF MULT
C
5  DOUBLE A,B,C,D,E,F,G,H,I,DO,DN
COMMON/D9/IQ,IN,Q,RN
DQ=Q & DN=RN
G=A*D+DN*(B*F+C*E)
10 H=A*E+B*D+DQ*(B*F+C*E)+DN*C*F
I=A*F+B*E+C*D+DQ*C*F
RETURN
END
```



```

SUBROUTINE TEST
C...THIS SUBROUTINE TRIES CONGRUENCES TO MODULI OF NORM MF WHERE MF IS A
C...FACTOR OF NORM(E+ICHECK) AND SETS EUCLID TO 1 IF NOT ALL THE RESIDUE
C...CLASSES TO ONE OF THE MODULI OF NORM MF CONTAIN AN INTEGER OF NORM
5 C...LESS THAN MF.
INTEGER EUCLID
DOUBLE IM,JM,KM,JP,JR,JS
COMMON/D3/MF,MP,MAP(500,2)
COMMON/D4/IP(2000),IR(2000),IS(2000)
10 COMMON/D5/IIM,JJM,KKM/D6/EUCLID
COMMON/D21/I,J,K
COMMON/D30/ITHETA,ILAM,IPRIME
DIMENSION ICP(500),ICR(500),ICS(500),ICD(500),LIC(500)
IM=IIM & JM=JJM & KM=KKM
15 JCOUNT=-1
LMAP1=MAP(MF,1) & LMAP2=MAP(MF,2)
DO 7 JJ=LMAP1,LMAP2
C...NOW USE IN TURN EACH OF THE ALGEBRAIC INTEGERS OF NORM MF.
JP=IP(JJ) & JR=IR(JJ) & JS=IS(JJ)
20 JJP=JP & JJR=JR & JJS=JS
C...CHECK TO SEE IF JP*Y+JR*X+JS IS A FACTOR OF IM*Y+JM*X+KM IF IT IS
C...NOT GO ON TO THE NEXT NUMBER OF NORM MF.
IF(IALMOD(IM,JM,KM,JP,JR,JS),EQ,1) GO TO 7
ICL=MF-1
25 C...ICL IS THE NUMBER OF NON-ZERO RESIDUE CLASSES MODULO (JP,JR,JS).
CALL INCON(JP,JR,JS,ICL,ICP,ICR,ICS,LIC)
C...LIC(I) IS SET TO 1 INITIALLY AND TO 0 WHEN AN INTEGER OF NORM LESS
C...THAN MF IS FOUND TO BELONG TO ICP(I)*Y+ICR(I)*X+ICS(I). NOW RUN
C...THROUGH THE INTEGERS WITH NORM LESS THAN MF AND ELIMINATE THE
30 C...RESIDUE CLASSES TO WHICH THEY BELONG.
IPRIME=0
ITHETA=0 & ILAM=0
C...IPRIME=1 WHEN X AND Y ARE CONGRUENT TO RATIONAL INTEGERS AND
C...IPRIME=0 OTHERWISE.
35 IF(ICS(ICL),NE,ICL) GO TO 20
IPRIME=2
C...IPRIME=2 WHEN FINDING THE VALUES OF ILAM AND OF ITHETA.
CALL MODAL(1,0,0,ICP,ICR,ICS,LIC,ICL,JP,JR,JS)
DO 21 JJC=1,ICL
40 IF(LIC(JJC),NE,0) GO TO 21
ILAM=JJC
LIC(JJC)=1
GO TO 22
21 CONTINUE
45 22 CALL MODAL(0,1,0,ICP,ICR,ICS,LIC,ICL,JP,JR,JS)
DO 23 JJC=1,ICL
IF(LIC(JJC),NE,0) GO TO 23
ITHETA=JJC
LIC(JJC)=1
50 GO TO 24

```

```

23 CONTINUE
24 IPRIME=1
   PRINT 999,ITHETA,ILAM
999 FORMAT(26H MODULUS IS PRIME, ITHETA=,I5,6H ILAM=,I5)
55 C...X IS CONGRUENT TO ITHETA AND Y TO ILAM.
   20 CALL MODAL(I,J,K,ICP,ICR,ICS,LIC,ICL,JP,JR,JS)
      DO 8 JJC=2,ICL
         IF(MAP(JJC,1).EQ.-1) GO TO 8
         MMAP1=MAP(JJC,1) & MMAP2=MAP(JJC,2)
60       DO 9 JJD=MMAP1,MMAP2
          KP=IP(JJD) & KR=IR(JJD) & KS=IS(JJD)
          CALL MODAL(KP,KR,KS,ICP,ICR,ICS,LIC,ICL,JP,JR,JS)
          9 CONTINUE
            JCOUNT=0
65           DO 10 ICOUNT=1,ICL
              IF(LIC(ICOUNT).EQ.0) GO TO 10
              JCOUNT=JCOUNT+1
              ICD(JCOUNT)=ICOUNT
            10 CONTINUE
70 C.....JCOUNT IS THE NUMBER OF RESIDUE CLASSES STILL UNCOVERED. IF THE
C.....RESIDUE CLASSES ARE ALL COVERED GO ON TO CONGRUENCES TO ANOTHER
C.....MODULUS, OTHERWISE GO ON TO INTEGERS OF A GREATER NORM.
      IF(JCOUNT.EQ.0) GO TO 11
      8 CONTINUE
75 C...NOW THE ONLY REMAINING RESIDUE CLASSES WITH LIC=1 ARE THOSE WHICH
C...DO NOT CONTAIN AN INTEGER OF NORM LESS THAN MF.
C...IF JCOUNT=-1 THEN THERE ARE NO INTEGERS WITH NORM LESS THAN MF
C...OTHER THAN E AND SO WE GO ON TO FIND WHICH RESIDUE CLASSES ARE
C...UNCOVERED SINCE THIS HAS NOT BEEN DONE.
80     IF(JCOUNT.NE.-1) GO TO 13
        JCOUNT=0
        DO 12 ICT=1,ICL
           IF(LIC(ICT).EQ.0) GO TO 12
85         JCOUNT=JCOUNT+1
           ICD(JCOUNT)=ICT
        12 CONTINUE
           IF(JCOUNT.EQ.0) GO TO 11
90 C...IF THERE ARE STILL SOME UNCOVERED RESIDUE CLASSES THEN THE FIELD IS
C...NON-EUCLIDEAN SO THIS FACT IS PRINTED TOGETHER WITH THE MODULUS AND
C...UNCOVERED RESIDUE CLASSES.
        13 PRINT 102,JJP,JJR,JJS
        102 FORMAT(1X,67HTHE FIELD IS SHOWN TO BE NON-EUCLIDEAN WHEN CONSIDERIN
          G CONGRUENCES,/,1X,14HTO THE MODULUS,3I16,/,1X,39HWHEN THE UNCOVE
          2RED RESIDUE CLASSES ARE-)
95         DO 14 ICT=1,JCOUNT
            ID=ICD(ICT)
            PRINT 103,ICP(ID),ICR(ID),ICS(ID)
        103 FORMAT(15X,3I16)
        14 CONTINUE
100        EUCLID=1

```



```

SURROUTINE INCON(JP,JR,JS,M,ICP,ICR,ICS,LIC)
C...INCON FINDS THE RESIDUE CLASSES MODULO JP*Y+JR*X+JS EXCLUDING THE
C...ZERO RESIDUE CLASS. THESE ARE ICP(I)*Y+ICR(I)*X+ICS(I) FOR I=1,M,
C...THE CORRESPONDING LIC(I) IS SET TO 1. M=NORM(JP*Y+JR*X+JS)-1.
5   DOUBLE DJ,DJJ,JP,JR,JS,DARG1,DARG2
COMMON/D1/IA,IB,IC,INDEX,IDET
DIMENSION ICP(M),ICR(M),ICS(M),LIC(M) GO TO 2
C...NORN(JJ,J)=NORN(JJ*X+J).
NORN(JJ,J)=JJ*JJ*JJ*IC+J*(JJ*JJ*IB+J*(JJ*IA+J))
10  MM=M+1
    JJP=JP & JJR=JR & JJS=JS
    PRINT 160,JJP,JJR,JJS,MM
160  FORMAT(/,1X,35HCONGRUENCES ARE CONSIDERED MODULO (,I15,1H,,I15,1H,
15  1,I15,1H),/,1X,14HWHICH HAS NORM,I15)
    DO 1 I=1,M
      ICP(I)=ICR(I)=ICS(I)=0 & LIC(I)=1
      1 CONTINUE
      IP=IR=IS=0
      IPP=IRR=ISS=1
20  C...FIRST CALCULATE IS SUCH THAT 0,1,...IS ARE INCONGRUENT MODULO
C...JP*Y+JR*X+JS.
      DO 2 J=1,MM
        IF(MOD(J*J*J,MM).NE.0) GO TO 2
        JJP=JP & JJR=JR & JJS=JS
25        DJ=J
        DARG1=0.D0 & DARG2=0.D0
        IF(IALMOD(DARG1,DARG2,DJ,JP,JR,JS).EQ.1) GO TO 2
        IS=J-1 & GO TO 3
      2 CONTINUE
30      3 ISS=IS+1
        IF(ISS.EQ.MM) GO TO 7
        IS2=MIN0(ISS,MM/ISS)
C...NOW FIND IR SUCH THAT R*X+S FOR R=0,IR AND S=0,IS ARE INCONGRUENT
C...MODULO JP*Y+JR*X+JS.
35        DO 4 JJ=1,IS2
          DO 10 JL=1,ISS
            J=JL-1
            IF(MOD(NORN(JJ,J),MM).NE.0) GO TO 10
            DJJ=JJ & DJ=J
40            DARG1=0.D0
            IF(IALMOD(DARG1,DJJ,DJ,JP,JR,JS).EQ.1) GO TO 10
            IR=JJ-1 & GO TO 5
          10 CONTINUE
          4 CONTINUE
45          5 IRR=IR+1
            IF((ISS*IRR).EQ.MM) GO TO 7
            IPP=MM/(ISS*IRR)
            IPP=IPP-1
50  C...THE RESIDUE CLASSES ARE NOW (P*Y+R*X+S) P=0,IP, R=0,IR, S=0,IS.
      7 IJK=0

```



```

DO 8 II=1,IPP
I=II-1
DO 8 JJ=1,IRR
J=JJ-1
55 DO 8 KK=1,ISS
K=KK-1
IF((I.EQ.0).AND.(J.EQ.0).AND.(K.EQ.0)) GO TO 8
IJK=IJK+1
ICP(IJK)=I & ICR(IJK)=J & ICS(IJK)=K
60      8 CONTINUE
PRINT 161,IP,IR,IS
161 FORMAT(4X,46H THE RESIDUE CLASSES ARE (P*Y+R*X+S) WHERE P=0,,I4,
15H R=0,,I4,5H S=0,,I4,/,1X,18HP,R,S NOT ALL ZERO)
65 RETURN
END

```

```

SUBROUTINE MODAL(KP,KR,KS,ICP,ICR,ICS,LIC,M,JP,JR,JS)
C...MODAL FINDS WHICH ICP(I)*Y+ICR(I)*X+ICS(I) IS CONGRUENT TO
C...+OR-(KP*Y+KR*X+KS) MODULO JP*Y+JR*X+JS AND SETS THE CORRESPONDING
C...LIC(I) TO 0.
5   DOUBLE IP,IR,IS,JP,JR,JS
COMMON/D30/ITHETA,ILAM,IPRIME
DIMENSION ICP(M),ICR(M),ICS(M),LIC(M),LR,LS,LOI
IF(IPRIME.NE.1) GO TO 5
MM=M+1
10  KKP=MOD(KP,MM)
KKR=MOD(KR,MM)
KKS=MOD(KS,MM)
KPRS=KKP*ILAM+KKR*ITHETA+KKS
KPRS=MOD(KPRS,MM)
15  IF(KPRS.LT.0) KPRS=MM+KPRS
LIC(KPRS)=0
KPRS=MM-KPRS
LIC(KPRS)=0
RETURN
20  5 DO 1 I=1,M
IP=KP-ICP(I) & IR=KR-ICR(I) & IS=KS-ICS(I)
MODIAL =IALMOD(IP,IR,IS,JP,JR,JS)
IF(LIC(I).NE.0) LIC(I)=MODIAL
IF(MODIAL.EQ.0) GO TO 3
25  1 CONTINUE
3  IF(IPRIME.EQ.2) RETURN
DO 2 I=1,M
IP=KP+ICP(I) & IR=KR+ICR(I) & IS=KS+ICS(I)
MODIAL=IALMOD(IP,IR,IS,JP,JR,JS)
30  IF(LIC(I).NE.0) LIC(I)=MODIAL
IF(MODIAL.EQ.0) GO TO 4
2  CONTINUE
4  RETURN
END

```

```

FUNCTION IALMOD(IP,IR,IS,JP,JR,JS)
C... THE FUNCTION IALMOD TAKES THE VALUE 0 IF JP*Y+JR*X+JS DIVIDES IP*Y+IR*X+IS. OTHERWISE IT TAKES THE VALUE 1.
C... DOUBLE IP,IR,IS,JP,JR,JS,ID,JD,LP,LR,LS,LD AND SEE THIS COMMENT
5   IALMOD=0
    ID=1.D0 5 / JD=1.D0
    CALL DIVCD(IP,IR,IS,ID,JP,JR,JS,JD,LP,LR,LS,LD)
    IF (LD.NE.1.D0) IALMOD=1
    RETURN
10  END
END

```

```

SUBROUTINE DIVC(L1,M1,N1,L2,M2,N2,L,M,N)
C...THE SUBROUTINE DIVC FINDS THE QUOTIENT OF THE TWO ALGEBRAIC NUMBERS
C...L1*Y+M1*X+N1 AND L2*Y+M2*X+N2 IN THE CUBIC FIELD WITH INTEGRAL
C...BASIS (1,X,Y), WHERE X**3-IA*X**2+IB*X-IC=0, AND SETS THIS QUOTIENT
5 C...TO L*Y+M*X+N.
COMMON/D2/IL1,IL2,IL3,IT1,IT2,IT3,IP1,IP2,IP3
COMMON/D15/IURN,IOR,IORDS
KDET(I1,I2,I3,J1,J2,J3,K1,K2,K3)=I1*J2*K3+I2*J3*K1+I3*J1*K2
1-I1*J3*K2-I2*J1*K3-I3*J2*K1
10 C...FIRST TEST FOR THE POSSIBILITY OF OVERFLOW.
LA=IABS(L1) E MA=IABS(M1) E NA=IABS(N1)
IDIV1=MAX0(LA,MA,NA)
LA=IABS(L2) E MA=IABS(M2) E NA=IABS(N2)
IDIV2=MAX0(LA,MA,NA)
15 IF(IDIV1.GT.IDIV2) GO TO 1
IDIV=IDIV2*IDIV2*IDIV2
GO TO 2
1 IDIV=IDIV1*IDIV2*IDIV2
2 IF(IDIV.LT.IORDS) GO TO 3
20 PRINT 50
50 FORMAT(1X,10(1H*),26HDANGER OF OVERFLOW IN DIVC,10(1H*))
3 I1=L2*IL1+IP1*M2+N2
I2=L2*IL2+IP2*M2
25 I3=L2*IL3+IP3*M2
J1=IP1*L2+M2*IT1
J2=IP2*L2+M2*IT2+N2
J3=IP3*L2+M2*IT3
K1=L2 E K2=M2 E K3=N2
30 JDET=KDET(I1,I2,I3,J1,J2,J3,K1,K2,K3)
L=KDET(L1,M1,N1,J1,J2,J3,K1,K2,K3)/JDET
M=KDET(I1,I2,I3,L1,M1,N1,K1,K2,K3)/JDET
N=KDET(I1,I2,I3,J1,J2,J3,L1,M1,N1)/JDET
RETURN
END

```

```

SUBROUTINE MULTC(L1,M1,N1,L2,M2,N2,L,M,N)
C...THE SUBROUTINE MULTC FINDS THE PRODUCT OF THE TWO ALGEBRAIC NUMBERS
C...L1*Y+M1*X+N1 AND L2*Y+M2*X+N2 IN THE CUBIC FIELD WITH INTEGRAL
C...BASIS (1,X,Y), WHERE X**3-IA*X**2+IB*X-IC=0, AND SETS THIS PRODUCT
5 C...TO L*Y+M*X+N.
COMMON/D2/IL1,IL2,IL3,IT1,IT2,IT3,IP1,IP2,IP3
COMMON/D15/IURN,I0B,I0BDS
C...FIRST TEST FOR THE POSSIBILITY OF OVERFLOW.
10 LA=IABS(L1) & MA=IABS(M1) & NA=IABS(N1)
MULT1=MAX0(LA,MA,NA)
LA=IABS(L2) & MA=IABS(M2) & NA=IABS(N2)
MULT2=MAX0(LA,MA,NA)
MULT=MULT1*MULT2
15 IF (MULT.LT.I0B) GO TO 1
PRINT 50
50 FORMAT(1X,10(1H*),27HDANGER OF OVERFLOW IN MULTC,10(1H*))
1 IP=L1*M2+L2*M1
L=L1*L2*IL1+IP*IP1+M1*M2*IT1+L1*N2+N1*L2
20 M=L1*L2*IL2+IP*IP2+M1*M2*IT2+M1*N2+N1*M2
N=L1*L2*IL3+IP*IP3+M1*M2*IT3+N1*N2
1 RETURN
100 END

```

```

DOUBLE FUNCTION DNORM(IP,IR,IS)
C...DNORM(IP,IR,IS)=NORM(IP*Y+IR*X+IS).
DOUBLE IP,IR,IS,P,R,S,AD,BD,CD,DEX,DEX3,DIUBN,DT2,DT3
DOUBLE PA,RA,SA
5 COMMON/D1/IA,IR,IC,INDEX,IDET
COMMON/D2/IL1,IL2,IL3,IT1,IT2,IT3,IP1,IP2,IP3
COMMON/D15/IUBN,IOB,IOBDS
AD=IA & BD=IB & CD=IC & DEX=INDEX
DEX3=DEX*DEX*DEX
10 DIUBN=IUBN & DT2=IT2 & DT3=IT3
C...NOW FIND P,R,S SUCH THAT (P*X**2+R*X+S)/INDEX=IP*Y+IR*X+IS.
P=IP
R=IR*DEX-DT2*IP
S=IS*DEX-DT3*IP
15 C...IF ANY OF P,R,S ARE SO LARGE THAT THEY MAY CAUSE 'OVERFLOW' IN THE
C...CALCULATION OF DNORM THEN PRINT A DIAGNOSTIC.
PA=DABS(P) & RA=DABS(R) & SA=DABS(S)
IF(PA.GE.DIUBN) GO TO 1
IF(RA.GE.DIUBN) GO TO 1
20 IF(SA.LT.DIUBN) GO TO 2
1 PRINT 100
100 FORMAT(/,1X,10(1H*),48HDANGER OF OVERFLOW WHEN USING THE FUNCTION
1DNORM,10(1H*))
2 DNORM=(P*P*P*CD*CD+R*R*R*CD+S*S*S+P*P*(R*BD*CD+S*(BD*BD-2.0D0*AD*
25 1CD))+R*R*(P*AD*CD+S*BD)+S*S*(P*(AD*AD-2.0D0*BD)+R*AD)+P*R*S*(AD*BD
2-3.0D0*CD))/DEX3
3 RETURN
END

```



```

DOUBLE FUNCTION FDET(I1,I2,I3,J1,J2,J3,K1,K2,K3)
DOUBLE I1,I2,I3,J1,J2,J3,K1,K2,K3
FDET=I1*J2*K3+I2*J3*K1+I3*J1*K2-I1*J3*K2-I2*J1*K3-I3*J2*K1
RETURN
END

```

5

THE PROGRAMS CUBOID, FCUB AND CURX

CUBOID 58  
 FCUB 69  
 CURX 65

SUBROUTINES - ALL LISTED IN THIS SECTION

BASIS	68	COVER (CUBOID ONLY)	75
USE	70	C1	74
BOUND	72	COVER (FCUB AND CURX)	75



```

PROGRAM CUBOID(INPUT,OUTPUT,PUNCH)

COMMON/D1/IA,IB,IC,INDEX,IDET/D2/R(3),U(3),H(3)
COMMON/D3/PM(1000),PN(1000),PO(1000),KZ/D4/BK,CK
5 COMMON/D7/XL,XU,YL,YU,ZL,ZU,XE,YE,ZE
COMMON/D8/IBC
DIMENSION X(11),Y(11),Z(11),L(10,10,10)
GNOM(CX,CY,CZ,DX,DY,DZ)=CX*DX+CY*DY+CZ*DZ

10 C...READ IN THE DETAILS OF THE FIELD AND FIND A BASIS. THEN FIND THE
C...SMALLEST CUBOID CONTAINING THE FUNDAMENTAL REGION OF POINTS
C...(XC,YC,ZC) RELATIVE TO THE BASIS FOR WHICH -0.5<XC<0.5; -0.5<YC<0.5;
C...-0.5<ZC<0.5.
111 FORMAT(I6,I2,3I5)
15 READ 111, IDET, INDEX, IA, IB, IC
CALL BASIS
CALL ROUND(R,XL,XU)
XF=XU-XL
CALL ROUND(U,YL,YU)
20 YE=YU-YL
CALL ROUND(H,ZL,ZU)
ZF=ZU-ZL
XU=XE/2.0 & XL=-XU
25 YU=YE/2.0 & YL=-YU
ZU=ZF/2.0 & ZL=-ZU
116 FORMAT(/,2X,*THE SMALLEST CUBOID WITH EDGES PARALLEL TO THE CO-ORD
1 INATE AXES*,/1X,*WHICH CONTAINS A FUNDAMENTAL REGION CONTAINS THO
2 SE POINTS (X,Y,Z) WHERE*,/5X,*X IS BETWEEN*,F8.4,1X,*AND*,F8.4,/,
35X,*Y IS BETWEEN*,F8.4,1X,*AND*,F8.4,/,5X,*Z IS BETWEEN*,F8.4,1X,
30 4*AND*,F8.4)
PRINT 116,XL,XU,YL,YU,ZL,ZU
C...THE SMALLEST CUBOID CONTAINING THE FUNDAMENTAL REGION IS GIVEN BY
C...XL<X<XU, YL<Y<YU, ZL<Z<ZU.

35 C...READ IN THE DETAILS OF THE SUB-DIVISION.
C...FIND THE SUITABLE INTEGER POINTS WITH COEFFICIENTS LESS THAN IBC.
C...INITIALIZE THE MARKER ARRAY L, SETTING L(I,J,K) TO 2 FOR THOSE
C...SUB-REGIONS OF THE CUBOID WHICH DO NOT INTERSECT THE FUNDAMENTAL
C...REGION AND TO 1 OTHERWISE.
40 112 FORMAT(F7.4)
READ 112,RK
READ 112,AM
CTN=.1000
45 119 FORMAT(I5)
READ 119,IBC
CALL USE
70 READ 119,LIMIT
DO 12 I=1,10
80 DO 12 J=1,10
50 DO 12 K=6,10

```

```

12 L(I,J,K)=1
DO 13 I=1,11
FI=FLOAT(I-1) & RX=FI*CCIN
X(I)=XL+RX*XE & Y(I)=YL+RX*YE & Z(I)=ZL+RX*ZE
55 13 CONTINUE
BDET=R(1)*U(2)*H(3)+R(2)*U(3)*H(1)+R(3)*U(1)*H(2)-R(1)*U(3)*H(2)
1-R(2)*U(1)*H(3)-R(3)*U(2)*H(1)
X12=(U(1)*H(2)-U(2)*H(1))/BDET
X23=(U(2)*H(3)-U(3)*H(2))/BDET
60 X31=(U(3)*H(1)-U(1)*H(3))/BDET
Y12=(H(1)*R(2)-H(2)*R(1))/BDET
Y23=(H(2)*R(3)-H(3)*R(2))/BDET
Y31=(H(3)*R(1)-H(1)*R(3))/BDET
Z12=(R(1)*U(2)-R(2)*U(1))/BDET
65 Z23=(R(2)*U(3)-R(3)*U(2))/BDET
Z31=(R(3)*U(1)-R(1)*U(3))/BDET
DO 35 I=1,10
DO 25 J=1,10
DO 15 K=6,10
70 DO 36 I1=1,2
DO 26 J1=1,2
DO 16 K1=1,2
I2=I+I1-1 & J2=J+J1-1 & K2=K+K1-1
75 XC=CNOM(X(I2),Y(J2),Z(K2),X23,Y23,Z23)
YC=CNOM(X(I2),Y(J2),Z(K2),X31,Y31,Z31)
ZC=CNOM(X(I2),Y(J2),Z(K2),X12,Y12,Z12)
IF((ABS(XC).LE.0.5).AND.(ABS(YC).LE.0.5).AND.(ABS(ZC).LE.0.5))
160 TO 15
16 CONTINUE
80 26 CONTINUE
36 CONTINUE
L(I,J,K)=2
15 CONTINUE
25 CONTINUE
85 35 CONTINUE

C...NOW FIND WHICH REGIONS ARE LEFT UNCOVERED. BY ALTERATION OF BK.
C...IF NECESSARY, REDUCE THE NUMBER OF UNCOVERED REGIONS TO LESS THAN
C...LIMIT.
90 118 FORMAT(/,1X,*BK=*,F7.4)
8 PRINT 118,BK
CALL COVER(X,Y,Z,L,ICOUNT)
PRINT 113,ICOUNT
113 FORMAT(1X,*THERE ARE NOW*,I5,1X,*UNCOVERED REGIONS*)
95 IF(ICOUNT.GT.LIMIT) GO TO 19
IF(ICOUNT.EQ.0) GO TO 3
201 FORMAT(I6)
PUNCH 201,IDET
101 FORMAT(1X,*THESE HAVE VERTICES:*)
100 PRINT 101

```

```

DO 17 I=1,10 (INPUT,OUTPUT,PUNCH)
DO 17 J=1,10
DO 17 K=6,10
IF (L(I,J,K).NE.1) GO TO 17
105 102 FORMAT(2X,3F15.10)
PRINT 102,X(I),Y(J),Z(K)
202 FORMAT(3F15.10)
PUNCH 202,X(I),Y(J),Z(K)
17 CONTINUE
110 GO TO 3
19 IF (ICOUNT.LE.100) GO TO 9
AN=AM *E GO TO 11
9 AN=0.1*AM
11 BK=BK+CIN*AN *E GO TO 8
115 3 STOP
END BASIS
ALL POINTS IN X
SECT=ZL
ALL POINTS IN Y
SECT=ZL
CALL SUBROUTINE
SECT=ZL
END BASIS
...READ THE COORDINATES OF THE SUBDIVISION AND FIND THE SUITABLE
...INTEGER POINTS. THEN READ IN THE NUMBER OF REGIONS WHICH ARE
...TO BE CONSIDERED.
112 FORMAT(7,F15.10)
READ 112,N
READ 113,CIN
113 FORMAT(I5)
READ 114,INC
CALL ONE
READ 115,NR
READ 116,M
DO 3 I=1,M
...FOR WORK THROUGH EACH REGION IN TURN. FIRST INITIALIZE THE MARKER
...ARRAY LA SETTING LA(I,J,K) TO 2 FOR THOSE SUBREGIONS OF THE
...CUBOID WHICH DO NOT INTERSECT THE FUNDAMENTAL REGION AND TO 1 OTHERWISE.
120 FORMAT(5F15.10)
READ 120,K,M,N,Z0
121 FORMAT(2,F15.10)
READ 121,X1,X2
DO 13 I=1,I1
PUNCH(2X,I1-1) *E KAK+CIN
X11=X1+VXVX *E Y11=Y1+VYVY *E Z11=Z0+VZVZ
13 CONTINUE

```

```

PROGRAM FCUB(INPUT,OUTPUT,PUNCH)

COMMON/D1/IA,IB,IC,INDEX,IDET/D2/R(3),U(3),H(3)
COMMON/D3/PM(1000),PN(1000),PO(1000),KZ/D4/BK,CIN
COMMON/D7/XL,XU,YL,YU,ZL,ZU,XE,YE,ZE
COMMON/D8/IBC
DIMENSION X(11),Y(11),Z(11),L(10,10,10)
CNOM(CX,CY,CZ,DX,DY,DZ)=CX*DX+CY*DY+CZ*DZ

5
10 C...READ IN THE DETAILS OF THE FIELD AND FIND A BASIS. THEN FIND THE
C...SMALLEST CUBOID CONTAINING THE FUNDAMENTAL REGION OF POINTS
C...(XC,YC,ZC) RELATIVE TO THE BASIS FOR WHICH -0.5<XC<0.5
C...-0.5<YC<0.5: -0.5<ZC<0.5.
111 FORMAT(I6,I2,3I5)
15 READ 111,IDEI,INDEX,IA,IB,IC
CALL BASIS
CALL ROUND(R,XL,XU)
XE=XU-XL
CALL ROUND(U,YL,YU)
20 YE=YU-YL
CALL ROUND(H,ZL,ZU)
ZF=ZU-ZL
XU=XE/2.0 & XL=-XU
YU=YE/2.0 & YL=-YU
25 ZU=ZE/2.0 & ZL=-ZU

C...READ THE DETAILS OF THE SUBDIVISION AND FIND THE SUITABLE
C...INTEGER POINTS. THEN READ IN THE NUMBER OF REGIONS WHICH ARE
C...TO BE CONSIDERED.
30 112 FORMAT(F7.4)
READ 112,BK
READ 112,CIN
119 FORMAT(I5)
READ 119,IBC
35 CALL USE
READ 119,LIMIT
READ 119,M
DO 3 ICT=1,M
40 C...NOW WORK THROUGH EACH REGION IN TURN. FIRST INITIALIZE THE MARKER
C...ARRAY L, SETTING L(I,J,K) TO 2 FOR THOSE SUBREGIONS OF THE
C...CUBOID WHICH DO NOT INTERSECT THE FUNDAMENTAL REGION AND TO 1 OTHERWISE.
120 FORMAT(3F15,10)
READ 120,XB,YB,ZB
45 121 FORMAT(/,1X,*THE REGION BEING SUBDIVIDED HAS VERTEX*,/,1X,3F15,10)
PRINT 121,XB,YB,ZB
DO 13 I=1,11
FI=FLOAT(I-1) & RX=FI*CIN
X(I)=XB+RX*XE & Y(I)=YB+RX*YE & Z(I)=ZB+RX*ZE
50 13 CONTINUE

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DO 12 I=1,10
DO 12 J=1,10
DO 12 K=1,10
12 L(I,J,K)=1
55 BDET=R(1)*U(2)*H(3)+R(2)*U(3)*H(1)+R(3)*U(1)*H(2)-R(1)*U(3)*H(2)
1-R(2)*U(1)*H(3)-R(3)*U(2)*H(1)
X12=(U(1)*H(2)-U(2)*H(1))/BDET
X23=(U(2)*H(3)-U(3)*H(2))/BDET
X31=(U(3)*H(1)-U(1)*H(3))/BDET
60 Y12=(H(1)*R(2)-H(2)*R(1))/BDET
Y23=(H(2)*R(3)-H(3)*R(2))/BDET
Y31=(H(3)*R(1)-H(1)*R(3))/BDET
Z12=(R(1)*U(2)-R(2)*U(1))/BDET
Z23=(R(2)*U(3)-R(3)*U(2))/BDET
65 Z31=(R(3)*U(1)-R(1)*U(3))/BDET
DO 41 I=1,10
DO 31 J=1,10
DO 21 K=1,10
DO 42 I1=1,2
70 DO 32 J1=1,2
DO 22 K1=1,2
I2=I+I1-1 & J2=J+J1-1 & K2=K+K1-1
XC=CNOM(X(I2),Y(J2),Z(K2),X23,Y23,Z23)
YC=CNOM(X(I2),Y(J2),Z(K2),X31,Y31,Z31)
75 ZC=CNOM(X(I2),Y(J2),Z(K2),X12,Y12,Z12)
IF((ABS(XC).LE.0.5).AND.(ABS(YC).LE.0.5).AND.(ABS(ZC).LE.0.5))
160 TO 21
22 CONTINUE
32 CONTINUE
80 42 CONTINUE
L(I,J,K)=2
21 CONTINUE
31 CONTINUE
41 CONTINUE
85 C...NOW FIND WHICH SUB-REGIONS ARE LEFT UNCOVERED WITH THE GIVEN VALUE
C...OF RK. IF THIS NUMBER IS LESS THAN OR EQUAL TO LIMIT THEN PRINT
C...OUT THE VERTICES OF THE UNCOVERED REGIONS.
CALL COVER(X,Y,Z,L,ICOUNT)
90 113 FORMAT(1X,*IT CONTAINS*,I5,1X,*UNCOVERED REGIONS*)
PRINT 113,ICOUNT
IF(ICOUNT.EQ.0) GO TO 3
IF(ICOUNT.GT.LIMIT) GO TO 3
101 FORMAT(1X,*THESE HAVE VERTICES*)
95 PRINT 101
201 FORMAT(I6)
PUNCH 201,IDET
DO 17 I=1,10
DO 17 J=1,10
100 DO 17 K=1,10

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PROGRAM CUBX(INPUT,OUTPUT,PUNCH)

COMMON/D1/IA,IB,IC,INDEX,IDET/D2/R(3),U(3),H(3)
COMMON/D3/PM(1000),PN(1000),PO(1000),KZ/D4/BK,CIN
COMMON/D7/XL,XU,YL,YU,ZL,ZU,XE,YE,ZE
COMMON/D8/IBC
DIMENSION X(11),Y(11),Z(11),L(10,10,10)
CNOM(CX,CY,CZ,DX,DY,DZ)=CX*DX+CY*DY+CZ*DZ

5
10 C...READ IN THE DETAILS OF THE FIELD AND FIND A BASIS. THEN FIND THE
C...SMALLEST CUBOID CONTAINING THE FUNDAMENTAL REGION OF POINTS
C...(XC,YC,ZC) RELATIVE TO THE BASIS FOR WHICH -0.5<XC<0.5;
C...-0.5<YC<0.5; -0.5<ZC<0.5.
111 FORMAT(I6,I2,3I5)
15 READ 111, IDET, INDEX, IA, IB, IC
CALL BASIS
CALL ROUND(R, XL, XU)
XE=XU-XL
CALL ROUND(U, YL, YU)
20 YE=YU-YL
CALL ROUND(H, ZL, ZU)
ZF=ZU-ZL
XU=XE/2.0 & XL=-XU
YU=YE/2.0 & YL=-YU
25 ZU=ZE/2.0 & ZL=-ZU

C...READ IN THE DETAILS OF THE SUBDIVISION AND FIND THE SUITABLE
C...INTEGER POINTS. THEN READ IN THE NUMBER OF REGIONS WHICH ARE
C...TO BE CONSIDERED.
30 112 FORMAT(F7.4)
READ 112, BK
READ 112, CIN
119 FORMAT(I5)
35 READ 119, TBC
CALL USE
READ 119, M

DO 3 TCT=1, M
40 C...NOW WORK THROUGH EACH REGION IN TURN. FIRST INITIALIZE THE MARKER
C...ARRAY L, SETTING L(I,J,K) TO 2 FOR THOSE SUBREGIONS OF THE
C...CUBOID WHICH DO NOT INTERSECT THE FUNDAMENTAL REGION AND TO 1 OTHERWISE.
120 FORMAT(3F15.10)
READ 120, XB, YB, ZB
121 FORMAT(/, IX, *THE REGION BEING SUBDIVIDED HAS VERTEX*, /, IX, 3F15.10)
45 PRINT 121, XB, YB, ZB
DO 13 I=1, 11
FI=FLOAT(I-1) & RX=FI*CIN
X(I)=XB+RX*XE & Y(I)=YB+RX*YE & Z(I)=ZB+RX*ZE
13 CONTINUE
50 DO 12 I=1, 10

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DO 12 J=1,10
DO 12 K=1,10
12 L(I,J,K)=1
55 BDET=R(1)*U(2)*H(3)+R(2)*U(3)*H(1)+R(3)*U(1)*H(2)-R(1)*U(3)*H(2)
1-R(2)*U(1)*H(3)-R(3)*U(2)*H(1)
X12=(U(1)*H(2)-U(2)*H(1))/BDET
X23=(U(2)*H(3)-U(3)*H(2))/BDET
X31=(U(3)*H(1)-U(1)*H(3))/BDET
Y12=(H(1)*R(2)-H(2)*R(1))/BDET
60 Y23=(H(2)*R(3)-H(3)*R(2))/BDET
Y31=(H(3)*R(1)-H(1)*R(3))/BDET
Z12=(R(1)*U(2)-R(2)*U(1))/BDET
Z23=(R(2)*U(3)-R(3)*U(2))/BDET
Z31=(R(3)*U(1)-R(1)*U(3))/BDET
65 DO 41 I=1,10
DO 31 J=1,10
DO 21 K=1,10
DO 42 I1=1,2
DO 32 J1=1,2
70 DO 22 K1=1,2
I2=I+I1-1 & J2=J+J1-1 & K2=K+K1-1
XC=CNOM(X(I2),Y(J2),Z(K2),X23,Y23,Z23)
YC=CNOM(X(I2),Y(J2),Z(K2),X31,Y31,Z31)
ZC=CNOM(X(I2),Y(J2),Z(K2),X12,Y12,Z12)
75 IF (ABS(XC).GT.0.5) GO TO 22
IF (ABS(YC).GT.0.5) GO TO 22
IF (ABS(ZC).LE.0.5) GO TO 21
22 CONTINUE
32 CONTINUE
80 42 CONTINUE
L(I,J,K)=2
21 CONTINUE
31 CONTINUE
41 CONTINUE
85 C...NOW FIND WHICH SUB-REGIONS ARE LEFT UNCOVERED WITH THE GIVEN VALUE
C...OF RK. IF THIS NUMBER IS LESS THAN OR EQUAL TO LIMIT THEN PRINT
C...BOUNDS FOR THE REGION CONTAINING THE UNCOVERED REGIONS.
CALL COVER(X,Y,Z,L,ICOUNT)
90 113 FORMAT(1X,*IT CONTAINS*,I5,1X,*UNCOVERED REGIONS*)
PRINT 113,ICOUNT
IF(ICOUNT.EQ.0) GO TO 3
101 FORMAT(1X,*THESE HAVE VERTICES BETWEEN*)
PRINT 101
95 201 FORMAT(I6)
PUNCH 201,IDET
XLW=10.0 & XUP=-10.0
YLW=10.0 & YUP=-10.0
ZLW=10.0 & ZUP=-10.0
100 DO 17 I=1,10

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DO 17 J=1,10
DO 17 K=1,10
IF (L(I,J,K).NE.1) GO TO 17
IF (X(I).LT.XLW) XLW=X(I)
IF (X(I).GT.XUP) XUP=X(I)
IF (Y(J).LT.YLW) YLW=Y(J)
IF (Y(J).GT.YUP) YUP=Y(J)
IF (Z(K).LT.ZLW) ZLW=Z(K)
IF (Z(K).GT.ZUP) ZUP=Z(K)
17 CONTINUE
901 FORMAT(/,5X,A2,2F15.10)
AX=2HX: E AY=2HY: E AZ=2HZ:
PRINT 901,AX,XLW,XUP
PRINT 901,AY,YLW,YUP
PRINT 901,AZ,ZLW,ZUP
XUP=XUP+CIN*XE E YUP=YUP+CIN*YE E ZUP=ZUP+CIN*ZE
904 FORMAT(/,* GIVING LOWER AND UPPER BOUNDS RESPECTIVELY ON X,Y,Z:*)
PRINT 904
902 FORMAT(5X,3F15.10)
PRINT 902,XLW,YLW,ZLW
PRINT 902,XUP,YUP,ZUP
903 FORMAT(3F15.10)
PUNCH 903,XLW,YLW,ZLW
PUNCH 903,XUP,YUP,ZUP
3 CONTINUE
STOP
END

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SUBROUTINE BASIS
COMMON/D1/IA,IB,IC,INDEX,IDEF/D2/R(3),U(3),H(3)
A=FLOAT(IA) & B=FLOAT(IB) & C=FLOAT(IC) & IN=INDEX
5 C...FIND H(2) WHERE THE FIELD IS K(H(2)) AND ALSO R(2) AND U(2) WHERE
C...R(2)+I,U(2) AND R(2)-I,U(2) ARE THE CONJUGATES OF H(2).
V1=(B-(A**2)/3.0)/3.0
V2=- (A*B/3.0-2.0*A**3/27.0-C)
10 V3=SQRT(V2**2-4.0*V1**3)
V4=(V2+V3)/2.0
V5=(V2-V3)/2.0
IF (V4.NE.0.0) GO TO 51
V6=0.0 & GO TO 52
51 D=1.0
IF (V4.GE.0.0) GO TO 53
15 V4=-V4 & D=-1.0
53 V6=EXP((ALOG(V4))/3.0)*D
52 IF (V5.NE.0.0) GO TO 54
V7=0.0 & GO TO 55
20 54 D=1.0
IF (V5.GE.0.0) GO TO 56
V5=-V5 & D=-1.0
56 V7=EXP((ALOG(V5))/3.0)*D
55 H(2)=V6+V7+A/3.0
25 R(2)=(A-H(2))/2.0
U(2)=SQRT(C/H(2)-R(2)**2)
H(1)=1.0 & R(1)=1.0 & U(1)=0.0
C... (R(1),U(1),H(1))=(1,0,1) IS THE LATTICE POINT CORRESPONDING TO UNITY
30 C...FIND T AND S WHERE Y=(H(2)+2*T.H(2)+S)/INDEX IS SUCH THAT
C...(1,H(2),Y) IS A BASIS OF THE FIELD.
IF (INDEX.NE.1) GO TO 57
T=0.0 & S=0.0 & GO TO 58
35 57 K1=IA**2-2*IB & K4=IB**2-2*IA*IC & K5=IA*IB-3*IC
IN2=IN**2 & IN3=IN2*IN
DO 59 ITC=1,IN
IT=ITC-1
DO 60 ISC=1,IN
IS=ISC-1
40 JA=K1+IT*IA+3*IS
IF (MOD(JA,IN).NE.0) GO TO 60
JB=K4+IT**2*IB+3*IS**2+IT*K5+2*IS*K1+2*IS*IT*IA
IF (MOD(JB,IN2).NE.0) GO TO 60
45 JC=IS**3+IT**3*IC+IC**2+IS**2*IT*IA+IS*K4+IS*IT**2*IB+IT*IB*IC
1+IT**2*IA*IC+IS**2*K1+IS*IT*K5
IF (MOD(JC,IN3).NE.0) GO TO 60
T=FLOAT(IT) & S=FLOAT(IS) & GO TO 58
60 CONTINUE
59 CONTINUE
50 58 XIN=FLOAT(INDEX)

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H(3)=(H(2)**2+T*H(2)+S)/XIN
R(3)=(R(2)**2-U(2)**2+T*R(2)+S)/XIN
U(3)=(2.0*R(2)*U(2)+T*U(2))/XIN
C...R(3)+I.U(3) AND R(3)-I.U(3) ARE THE CONJUGATES OF H(3).
55 IDETM=-IDET
PRINT 501,IDETM,INDEX,IA,IR,IC,H(2),T,S,H(3)
501 FORMAT(1H),1X,*THE FIELD HAS DISCRIMINANT*,I7,* INDEX*,I4,/,1X,*AN
1D POLYNOMIAL COEFFICIENTS A=*,I4,* B=*,I4,* C=*,I4,/,2X,*A BASIS
60 2FOF THIS FIELD IS (1,H,L), L=(H,H+T.H+S)/INDEX,*,/1X,*WHERE H=*,
3F7.4,* T=*,F6.1,* S=*,F6.1,* AND L=*,F7.4)

C...FROM THE BASIS FOUND ABOVE FIND ANOTHER SUCH THAT THE FUNDAMENTAL
C...REGION IS 'ALMOST RECTANGULAR'.
ID=1
65 IF ((R(2)+H(2)).LT.0.0) ID=-1
IAK=IFIX(R(2)+H(2))
IF (MOD(IAK,2).NE.0) IAK=IAK+ID
AK=FLOAT(IAK)/2.0
70 A1=R(3)+H(3) & B1=R(2)+H(2) & C1=2.0
A2=R(3)*(R(2)-AK)+U(2)*U(3)+H(3)*(H(2)-AK)
B2=R(2)*(R(2)-AK)+U(2)*U(2)+H(2)*(H(2)-AK)
C2=R(2)+H(2)-2.0*AK
DET=B1*C2-B2*C1
P=(A1*C2-A2*C1)/DET
75 Q=(B1*A2-B2*A1)/DET
DP=DQ=1.0
IF (P.LT.0.0) DP=-1.0
IF (Q.LT.0.0) DQ=-1.0
IP=IFIX(P) & AP=FLOAT(IP)
80 IF (ABS(P-AP).GT.0.5) AP=AP+DP
IQ=IFIX(Q) & AQ=FLOAT(IQ)
IF (ABS(Q-AQ).GT.0.5) AQ=AQ+DQ
R(3)=R(3)-AP*R(2)-AQ
85 U(3)=U(3)-AP*U(2)
H(3)=H(3)-AP*H(2)-AQ
R(2)=R(2)-AK
H(2)=H(2)-AK
502 FORMAT(/,2X,*THE NEW BASIS IS 1,H=AK,L=AP.H-AQ WHERE AK=*,F6.1,
1* AP=*,F6.1,* AQ=*,F6.1)
90 PRINT 502,AK,AP,AQ
503 FORMAT(1X,*THUS GIVING THE BASIS FOR THE LATTICE OF INTEGERS*)
PRINT 503
504 FORMAT(1X,3F15.10)
95 DO 61 I=1,3
61 PRINT 504,R(I),U(I),H(I)
RETURN
END

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SUBROUTINE USE
COMMON/D2/R(3),U(3),H(3)/D3/PM(1000),PN(1000),PO(1000),KZ
COMMON/D4/BK,CIN/D7/XL,XU,YL,YU,ZL,ZU,XE,YE,ZE
COMMON/D8/IBC
5 C...FIND THE BOUNDS ON THE CARTESIAN CO-ORDINATES OF 'SUITABLE'
C...INTEGER POINTS.
XYM=AMAX1(XE,YE)
SQT1=SQRT(2.0*BK/(CIN*ZE))
10 SQT2=4.0*BK/((CIN*XYM)**2)
A1=XL-SQT1 & A2=XU+SQT1
B1=YL-SQT1 & B2=YU+SQT1
C1=ZL-SQT2 & C2=ZU+SQT2

15 C...WORK THROUGH ALL THE INTEGER POINTS WITH COEFFICIENTS LESS THAN IBC
C...IN ABSOLUTE VALUE. STARTING WITH THOSE 'NEAREST' TO THE
C...FUNDAMENTAL REGION.
KZ=0
IBCT=3*IBC+1 & IBCP=IBC+1
20 DO 35 ISIG=1,IBCT
IMU=MIN0(ISIG,IBCP)
DO 31 IM2=1,IMU
IM2=IM2-1
25 INU=MIN0((ISIG-IM2),IBCP)
DO 30 IN2=1,INU
IN2=IN2-1
IO2=ISIG-(IM2+IN2+1)
IF(IO2.GT.IBC) GO TO 30
DO 41 IMT=1,2
30 IMM=1
IF(IMT.EQ.2) IMM=-1
DO 42 INT=1,2
INN=1
IF(INT.EQ.2) INN=-1
35 DO 37 IOT=1,2
IOO=1
IF(IOT.EQ.2) IOO=-1
IM1=IMM*IM2 & IN1=INN*IN2 & IO1=IOO*IO2
RM=FLOAT(IM1) & RN=FLOAT(IN1) & RO=FLOAT(IO1)
40 IF(KZ.EQ.0) GO TO 36
C...TEST TO SEE THAT THIS INTEGER IS NOT ALREADY IN THE LIST, WHICH
C...MAY HAPPEN WHEN ONE OF THE COEFFICIENTS IS ZERO.
IZZ=MAX0(1,(KZ-8))
DO 39 JZ=IZZ,KZ
45 IF((RM.EQ.PM(JZ)).AND.(RN.EQ.PN(JZ)).AND.(RO.EQ.PO(JZ))) GO TO 37
39 CONTINUE
C...TEST TO FIND WHETHER THE INTEGER IS 'SUITABLE'
36 RJ=CI(RM,RN,RO,R) & UJ=CI(RM,RN,RO,U) & HJ=CI(RM,RN,RO,H)
IF(RJ.LT.(XL-0.1)) GO TO 34 & IF(RJ.GT.(XU+0.1)) GO TO 34
50 32 IF(UJ.LT.(YL-0.1)) GO TO 34 & IF(UJ.GT.(YU+0.1)) GO TO 34

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      IF(HJ.LT.C1) GO TO 34 & IF(HJ.LT.C2) GO TO 33
55  34 IF(HJ.LT.(ZL-0.1)) GO TO 37 & IF(HJ.GT.(ZU+0.1)) GO TO 37
      IF(RJ.LT.A1) GO TO 37 & IF(RJ.GT.A2) GO TO 37 AND-AU IS TH
      IF(UJ.LT.B1) GO TO 37 & IF(UJ.GT.B2) GO TO 37
      33 KZ=KZ+1 & PM(KZ)=RM & PN(KZ)=RN & PO(KZ)=RO
      C...TEST WHETHER THE ARRAYS HOLDING THE INTEGERS ARE 'FULL'.
      IF(KZ.GE.1000) GO TO 40
      37 CONTINUE
      42 CONTINUE
60  41 CONTINUE
      30 CONTINUE
      31 CONTINUE
      35 CONTINUE
65  302 FORMAT(/,2X,*FOR DISSECTING THE FUNDAMENTAL REGION INTO (1/CIN)*3
      1SUR-REGIONS WHERE CIN=*,F7.4,/,1X,*AND TAKING THE VALUE OF BK TO B
      2E*,F7.4,*, THEN*,I5,* INTEGER POINTS ARE USED*)
      40 PRINT 302,CIN,BK,KZ
      RETURN
      END

```

SUBROUTINE BOUND(R,AL,AU) (CONT)

C...AL IS THE MINIMUM OF THE EIGHT 'R' CO-ORDINATES AND AU IS THE  
C...MAXIMUM.

5 DIMENSION R(3) (1,1,1,1,1,1,1,1)  
ZERO=0.0  
C...THE FOLLOWING ARE THE 8 CORNERS OF THE REGION WHICH IT COVERS

10 ARG1=R(1)+R(2) & ARG2=R(2)+R(3) & ARG3=R(3)+R(1)  
ARG4=R(1)+R(2)+R(3)  
AL=AMIN1(ZERO,R(1),R(2),R(3),ARG1,ARG2,ARG3,ARG4)  
AU=AMAX1(ZERO,R(1),R(2),R(3),ARG1,ARG2,ARG3,ARG4)  
RETURN

END

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SUBROUTINE COVER(X,Y,Z,L,ICOUNT)
COMMON/D2/R(3),U(3),H(3)/D3/PM(1000),PN(1000),PO(1000),KZ
COMMON/D4/BK,CIN
DIMENSION X(11),Y(11),Z(11),L(10,10,10)
5 C...FOR EACH INTEGER SET THE L MARKERS OF THE REGIONS WHICH IT COVERS
C...TO ZERO AND COUNT THE NUMBER OF REMAINING UNCOVERED REGIONS.
DO 13 IE=1,KZ
10 PX=CI(PM(IE),PN(IE),PO(IE),R)
PY=CI(PM(IE),PN(IE),PO(IE),U)
PZ=CI(PM(IE),PN(IE),PO(IE),H)
ICOUNT=0
DO 20 I=1,10
DO 19 J=1,10
15 DO 18 K=6,10
IF(L(I,J,K).NE.1) GO TO 18
XD=AMAX1(ABS(X(I)-PX),ABS(X(I+1)-PX))
YD=AMAX1(ABS(Y(J)-PY),ABS(Y(J+1)-PY))
20 ZD=AMAX1(ABS(Z(K)-PZ),ABS(Z(K+1)-PZ))
TD=ZD*(XD**2+YD**2)
IF(TD.LT.BK) GO TO 12
ICOUNT=ICOUNT+1 & GO TO 18
12 L(I,J,K)=0
18 CONTINUE
25 19 CONTINUE
20 CONTINUE
C...IF ALL THE REGIONS ARE NOW COVERED RETURN TO THE MAIN PROGRAM.
IF(ICOUNT.EQ.0) GO TO 15
30 13 CONTINUE
15 RETURN
END

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FUNCTION CI(X,Y,Z,R) = Z\*L\*ICOUNT

C...THIS GIVES THE 'R' CO-ORDINATE OF THE POINT WITH COEFFICIENTS (X,Y,Z)

DIMENSION R(3)

CI=X\*R(1)+Y\*R(2)+Z\*R(3) / (L\*(L+10\*10))

C...IF INTEGER SET THE L MARKERS OF THE REGIONS WHICH IT COVERS

END=0 AND COUNT THE NUMBER OF REMAINING UNCOVERED REGIONS.

DO 12 I=1,82

IF (ABS(XI-PII)+ABS(YI-POI)+R)

IF (ABS(XI-PII)+ABS(YI-POI)+R)

IF (ABS(XI-PII)+ABS(YI-POI)+R)

ICOUNT=0

DO 20 I=1,10

DO 10 J=1,10

DO 10 K=1,10

IF (L(I,J,K) .NE. 1) GO TO 18

I=ABS(XI-PII)+ABS(YI-POI)+R

J=ABS(XI-PII)+ABS(YI-POI)+R

K=ABS(XI-PII)+ABS(YI-POI)+R

I=ZDMX(XI,YI,ZI)

IF (I .GT. R) GO TO 12

ICOUNT=ICOUNT+1 & GO TO 18

12 11, I, J, K=0

16 CONTINUE

18 CONTINUE

20 CONTINUE

C...IF ALL THE REGIONS ARE NOW COVERED RETURN TO THE MAIN PROGRAM.

IF (ICOUNT .EQ. 0) GO TO 15

15 CONTINUE

16 RETURN

END



```

SUBROUTINE COVER(X,Y,Z,L,ICOUNT)

COMMON/D2/R(3),U(3),H(3)/D3/PM(1000),PN(1000),PO(1000),KZ
COMMON/D4/BK,CIN
5 DIMENSION X(11),Y(11),Z(11),L(10,10,10)
C...FOR EACH INTEGER SET THE L MARKERS OF THE REGIONS WHICH IT COVERS
C...TO ZERO AND COUNT THE NUMBER OF REMAINING UNCOVERED REGIONS.
DO 13 IE=1,KZ
PX=CI(PM(IE),PN(IE),PO(IE),R)
10 PY=CI(PM(IE),PN(IE),PO(IE),U)
PZ=CI(PM(IE),PN(IE),PO(IE),H)
ICOUNT=0
DO 20 I=1,10
DO 19 J=1,10
15 DO 18 K=1,10
IF(L(I,J,K).NE.1) GO TO 18
XD=AMAX1(ABS(X(I)-PX),ABS(X(I+1)-PX))
YD=AMAX1(ABS(Y(J)-PY),ABS(Y(J+1)-PY))
ZD=AMAX1(ABS(Z(K)-PZ),ABS(Z(K+1)-PZ))
20 TD=ZD*(XD**2+YD**2)
IF(TD.LT.BK) GO TO 12
ICOUNT=ICOUNT+1 & GO TO 18
12 L(I,J,K)=0
18 CONTINUE
25 19 CONTINUE
20 CONTINUE
C...IF ALL THE REGIONS ARE NOW COVERED RETURN TO THE MAIN PROGRAM.
IF(ICOUNT.EQ.0) GO TO 15
13 CONTINUE
30 15 RETURN
END

```



```

PROGRAM TRANS(INPUT,OUTPUT)
C...THIS PROGRAM TRANSFORMS THE N UNCOVERED REGIONS AND DETERMINES THE
C...INTERSECTION OF THE TRANSFORMS WITH THE ORIGINAL REGIONS.
5     COMMON/D1/IA,IB,IC,INDEX,IDET/D2/R(3),U(3),H(3)
       DIMENSION X(2,100),Y(2,100),Z(2,100)
C...(R(I),U(I),H(I)) I=1,3 IS A BASIS FOR THE LATTICE OF INTEGERS.
C...A MAXIMUM OF 100 UNCOVERED REGIONS MAY BE CONSIDERED WHERE THE ITH
C...REGION HAS LOWER LIMITS ON ITS CARTESIAN CO-ORDINATES
10    C...(X(1,I),Y(1,I),Z(1,I)) AND UPPER LIMITS (X(2,I),Y(2,I),Z(2,I)).
       CNOM(CX,CY,CZ,DX,DY,DZ)=CX*DX+CY*DY+CZ*DZ
C...THIS FUNCTION IS USED TO CALCULATE THE CO-ORDINATES OF A POINT
C...RELATIVE TO THE BASE WHEN ITS CARTESIAN CO-ORDINATES ARE KNOWN.
C...READ THE DETAILS OF THE FIELD.
15    READ 100,IDET,INDEX,IA,IB,IC
       100 FORMAT(I6,I2,3I5)
       CALL BASIS
C...THE COEFFICIENTS OF THE FUNDAMENTAL UNIT RELATIVE TO THE BASIS
C...(H(1),H(2),H(3)) FOR THE FIELD ARE READ.
20    READ 101,EXC,EYC,EZC
       101 FORMAT(3F15,10)
C...THE CARTESIAN THEN THE CYLINDRICAL POLAR CO-ORDINATES OF THE
C...FUNDAMENTAL UNIT ARE NOW CALCULATED.
       EX=CI(EXC,EYC,EZC,R)
       EY=CI(EXC,EYC,EZC,U)
25    EZ=CI(EXC,EYC,EZC,H)
       ER=SQRT(EX*EX+EY*EY)
       EALP=ATAN2(EY,EX)
       112 FORMAT(/,47H THE UNIT HAS CO-ORDINATES RELATIVE TO THE BASE,
30    13F15,10,/,23H CARTESIAN CO-ORDINATES,3F15,10,/,35H AND CYLINDRICAL
       2 POLAR CO-ORDINATES,3F15,10)
       PRINT 112,EXC,EYC,EZC,EX,EY,EZ,ER,EALP,EZ
C...READ N, THE NUMBER OF REGIONS TO BE CONSIDERED, THEN THE BOUNDS ON
C...THEIR CARTESIAN CO-ORDINATES.
35    READ 102,N
       102 FORMAT(15)
       DO 50 I=1,N
       DO 51 J=1,2
       READ 101,X(J,I),Y(J,I),Z(J,I)
       X(J,I)=DCORR(X(J,I),J)
40    Y(J,I)=DCORR(Y(J,I),J)
       Z(J,I)=DCORR(Z(J,I),J)
       51 CONTINUE
       50 CONTINUE
       DET=FLOAT(IDET)
45    BDET=SQRT(DET)/2.0
C...BDET IS THE DETERMINANT OF THE REAL LATTICE. THE FOLLOWING NINE
C...VALUES ARE FOR USE IN CALCULATING COEFFICIENTS RELATIVE TO THE BASE
C...FROM CARTESIAN CO-ORDINATES.
50    X12=ADJ(U,H,1,2)/BDET
       X23=ADJ(U,H,2,3)/BDET

```

```

X31=ADJ(U,H,3,1)/BDET
Y12=ADJ(H,R,1,2)/BDET
Y23=ADJ(H,R,2,3)/BDET
Y31=ADJ(H,R,3,1)/BDET
55 Z12=ADJ(R,U,1,2)/BDET
Z23=ADJ(R,U,2,3)/BDET
Z31=ADJ(R,U,3,1)/BDET
C...READ M1,M2 WHERE REGIONS M1 UP TO M2 ARE TO BE CONSIDERED IN THIS
C...RUN.
60 READ 102,M1
READ 102,M2
104 FORMAT(3H1 R,I3)
DO 2 I=M1,M2
C...WORK THROUGH THE REGIONS ONE BY ONE.
65 PRINT 104,I
C...FIRST FIND THE CYLINDRICAL POLAR CO-ORDINATES OF THE VERTICES OF THE
C...ITH REGION AND THEN THE CO-ORDINATES OF THE VERTICES OF THE TRANSFORM
DO 4 II=1,2
DO 5 JJ=1,2
70 RA=SQRT(X(II,I)*X(II,I)+Y(JJ,I)*Y(JJ,I))
ALP=ATAN2(Y(JJ,I),X(II,I))
TRA=RA*ER
TALP=ALP+EALP
TX=TRA*COS(TALP)
75 TY=TRA*SIN(TALP)
DO 6 KK=1,2
TZ=Z(KK,I)*EZ
C...NOW FIND THE CO-ORDINATES RELATIVE TO THE BASE OF THIS VERTEX OF
C...THE TRANSFORM
80 TXC=CNOM(TX,TY,TZ,X23,Y23,Z23)
TYC=CNOM(TX,TY,TZ,X31,Y31,Z31)
TZC=CNOM(TX,TY,TZ,X12,Y12,Z12)
PRINT 113,X(II,I),Y(JJ,I),Z(KK,I),RA,ALP,Z(KK,I),TRA,TALP,TZ,TX,
85 1TY,TZ,TXC,TYC,TZC
113 FORMAT(/,38H THE POINT WITH CARTESIAN CO-ORDINATES,3F15.10,/,
135H HAS CYLINDRICAL POLAR CO-ORDINATES,3F15.10,/,64H IT IS TRANSFO
2RMED INTO THAT WITH CYLINDRICAL POLAR CO-ORDINATES,3F15.10,/,
333H WHICH HAS CARTESIAN CO-ORDINATES,3F15.10,/,38H AND CO-ORDINATE
4S RELATIVE TO THE BASE,3F15.10)
90 C...XL,YL,ZL ARE TO BE THE LOWER LIMITS ON THE CARTESIAN CO-ORDINATES OF
C...REGION I AND XU,YU,ZU THE UPPER LIMITS. XLC,YLC,ZLC AND XUC,YUC,ZUC
C...ARE TO BE THE BOUNDS ON THE CO-ORDINATES RELATIVE TO THE BASE.
IF(II.NE.1) GO TO 7
IF(JJ.NE.1) GO TO 7
95 IF(KK.NE.1) GO TO 7
XL=XU=TX & YL=YU=TY & ZL=TZ
XLC=XUC=TXC & YLC=YUC=TYC & ZLC=ZUC=TZC
GO TO 6
7 IF(KK.EQ.2) GO TO 9
100 CALL UPLW(XL,XU,TX)

```

```

CALL UPLW(YL,YU,TY)
9 IF (II.NE.1) GO TO 8
IF (JJ.NE.1) GO TO 8
ZU=TZ
105 8 CALL UPLW(XLC,XUC,TXC)
CALL UPLW(YLC,YUC,TYC)
CALL UPLW(ZLC,ZUC,TZC)
6 CONTINUE
5 CONTINUE
110 4 CONTINUE
C...NOW FIND THE LIMITS ON THE CO-ORDINATES RELATIVE TO THE BASE OF
C...THE INTEGERS TO BE USED IN TRANSLATING THIS REGION BACK ON TO THE
C...FUNDAMENTAL REGION.
115 CALL ROUNDI(XLC,IXL)
CALL ROUNDI(XUC,IXU)
CALL ROUNDI(YLC,IYL)
CALL ROUNDI(YUC,IYU)
CALL ROUNDI(ZLC,IZL)
CALL ROUNDI(ZUC,IZU)
120 PRINT 114,XL,YL,ZL,XU,YU,ZU,IXL,IYL,IZL,IXU,IYU,IZU
114 FORMAT(//,5X,1HL,3F20.10,/,5X,1HU,3F20.10,/,13H INTEGERS: L,3I6,/,
1,12X,1HU,3I6)
IFLAG=0
C...IFLAG=0 WHEN NO REGION HAS BEEN FOUND WHICH INTERSECTS WITH A
C...TRANSLATE OF THE TRANSFORM OF REGION I, IFLAG=1 WHEN AT LEAST ONE
C...SUCH REGION HAS BEEN FOUND. FOR EACH POSSIBLE TRANSLATING INTEGER
C...NOW FIND WHICH REGIONS INTERSECT.
130 LEXTT=IXU-IXL+1 & MEXTT=IYU-IYL+1 & NEXTT=IZU-IZL+1
RLAST=FLOAT(IXL)-1.0 & RMAST=FLOAT(IYL)-1.0 & RNAST=FLOAT(IZL)-1.0
RL=RLAST
DO 10 JL=1,LEXTT
RL=RL+1.0
IL=IFIX(RL)
RM=RMAST
135 DO 11 JM=1,MEXTT
RM=RM+1.0
IM=IFIX(RM)
RN=RNAST
DO 12 JN=1,NEXTT
RN=RN+1.0
IN=IFIX(RN)
ZC=CI(RL,RM,RN,H)
XC=CI(RL,RM,RN,R)
YC=CI(RL,RM,RN,U)
145 C...FOR THIS INTEGER RUN THROUGH THE ORIGINAL UNCOVERED REGIONS.
DO 15 J=1,N
C...CALCULATE THE BOUNDS ON THE TRANSLATE OF ORIGINAL REGION J BY
C...THE INTEGER (L,M,N)
150 ZUC=Z(2,J)+ZC & ZLC=Z(1,J)+ZC
XUC=X(2,J)+XC & XLC=X(1,J)+XC

```

```

      YUC=Y(2,J)+YC & YLC=Y(1,J)+YC
      MJ=J
C...MJ IS SET TO J WHEN CONSIDERING THE JTH REGION AND TO -J WHEN
C...CONSIDERING ITS 'NEGATIVE'.
155      14 IF(ZL.GT.ZUC) GO TO 17
          IF(ZU.LT.ZLC) GO TO 17
          IF(XL.GT.XUC) GO TO 17
          IF(XU.LT.XLC) GO TO 17
          IF(YL.GT.YUC) GO TO 17
160      IF(YU.LT.YLC) GO TO 17
          PRINT 105,MJ,IL,IM,IN
          105 FORMAT(/,42H THIS REGION'S TRANSFORM INTERSECTS WITH R,I3,4H * [,
                13I5,14))
          GO TO 16
165      17 IF(MJ.LT.0) GO TO 15
          ZLC=-Z(2,J)+ZC & ZUC=-Z(1,J)+ZC
          XLC=-X(2,J)+XC & XUC=-X(1,J)+XC
          YLC=-Y(2,J)+YC & YUC=-Y(1,J)+YC
          MJ=-J
170      GO TO 14
          16 IFLAG=1
          15 CONTINUE
          12 CONTINUE
          11 CONTINUE
175      10 CONTINUE
          IF(IFLAG.EQ.1) GO TO 2
          PRINT 106
          106 FORMAT(/,68H THIS REGION'S TRANSFORM DOES NOT INTERSECT WITH ANY 0
                1RIGINAL REGION)
180      2 CONTINUE
          STOP
          END

```



```
FUNCTION ADJ(R,U,I,J)
DIMENSION R(3),U(3)
ADJ=R(I)*U(J)-R(J)*U(I)
RETURN
END
```

5

```
5 SUBROUTINE UPLW(RL,RU,TR)
  IF(RL.LE.TR) GO TO 10
  RL=TR & GO TO 11
10 IF(RU.LT.TR) RU=TR
11 RETURN
END
```

```
SUBROUTINE ROUNDI(RC,IR)
```

```
  ID=1
```

```
  IF(RC.LT.0.0) ID=-1
```

```
  IR=IFIX(RC)
```

```
  RIR=FLOAT(IR)
```

```
  IF(ABS(RC-RIR).GT.0.5) IR=IR+ID
```

```
  RETURN
```

```
END
```

```
DCORR=RC-IR
```

```
RTURN
```

```
END
```

```

FUNCTION DCORR(X,J)
C...THIS FUNCTION ALLOWS FOR ROUNDING ERROR ON THE BOUNDS ON THE
C...CARTESIAN CO-ORDINATES.
CORR=1.0E-10
5 DCORR=0.0
IF (X.EQ.0.0) RETURN
IF (J.EQ.2) GO TO 1
DCORR=X-CORR
RETURN
10 1 DCORR=X+CORR
RETURN
END

```

THE PROGRAM EXCEPT

SUBROUTINES

IN THIS SECTION		IN PREVIOUS SECTIONS	
BASES	91	CHANGE	48
BULFD	93	DNORM	88
LCWUP	94	D-T	23
LIMITL	95	IOF	34
LIMITU	95		
DAJ	97		



```

PROGRAM EXCEP(INPUT,OUTPUT)
C
C...THIS PROGRAM FINDS THE MINIMUM OF THE ALGEBRAIC NUMBER A.
C
5  DOUBLE P,U,H,AX,AY,AZ,BX,BY,BZ,AN,BE,BP,ABP,CK,ANV,ANW,BEV,BEW,ARD
DOUBLE ABV,ABW,SB,TL1,TU1,TL,TU,VL,VU,WL,WU,PT,PV,PW,QT,QV,QW
DOUBLE PL,PU,QL,QU,BDET,P,Q,S,PUA,QLT,OUT,PR,SLT,SUT,QUA
DOUBLE PLT,PUT,OR
DOUBLE DCI,DAJ,PQ,X,Y,Z,BNORM,APNORM,DNORM
10  COMMON/D1/IDET,INDEX,IA,IB,IC/D2/R(3),U(3),H(3)
COMMON/D3/TL,VL,WL,TU,VU,WU/D4/IL1,IL2,IL3,IT1,IT2,IT3,IP1,IP2,IP3
COMMON/D5/HDET
C... H(3)*H(3)=IT1*H(3)+IT2*H(2)+IT3 ...
C... H(2)*H(2)=IT1*H(3)+IT2*H(2)+IT3
15  C... H(3)*H(2)=IP1*H(3)+IP2*H(2)+IP3
C...READ IN THE DETAILS OF THE FIELD.
READ 100,IDET,INDEX,IA,IB,IC
100 FORMAT(16,I2,3I5)
MDET=-IDET
20  PRINT 200,MDET,INDEX,IA,IB,IC
200 FORMAT(1H1,2BH THIS FIELD HAS DISCRIMINANT,I7,6H INDEX,I2,2BH AND
IPOLYNOMIAL COEFFICIENTS,3I6)
C...CALCULATE THE ZEROS OF THE DEFINING POLYNOMIAL AND A BASIS FOR THE
C...PEAL LATTICE FOR THIS FIELD USING BASES.
25  CALL BASES
C...H(2)*R(2)+I,U(2) AND R(2)-I,U(2) ARE THE ZEROS. (H(1),H(2),H(3))
C...IS A BASIS FOR THE FIELD.
C
C... (IZ*H(2)**2+IY*H(2)+IX)/ID IS THE NUMERATOR AND
30  C... (JZ*H(2)**2+JY*H(2)+JZ)/JD IS THE DENOMINATOR OF THE ALGEBRAIC
C...NUMBER A. NOW READ THEIR COEFFICIENTS AND CALCULATE THEIR VALUES.
READ 101,IZ,IY,IX,ID,JZ,JY,JX,JD
101 FORMAT(8I10)
CALL CHANGE(IZ,IY,IX,ID,IIZ,IY,IIX)
35  CALL CHANGE(JZ,JY,JX,JD,JIJ,JIY,JIJX)
AX=FLOAT(IIX) & AY=FLOAT(IY) & AZ=FLOAT(IIZ)
BX=FLOAT(JIX) & BY=FLOAT(JIY) & BZ=FLOAT(JIJ)
AN=DCI(AX,AY,AZ,H)
RE=DCI(BX,BY,BZ,H)
40  C...THE TRANSFORMATION ON THE LATTICE WHICH TAKES P INTO ITSELF IS
C...Q = (Q-AN)/(1-BE). SO SAVE THE VALUES OF 1-BE AND OF -AN/(1-BE).
BP=1.0D0-RE
ABP=-AN/BP
C...THE TRANSFORMATION IS NOW Q=Q/BP+ABP.
45  C...READ THE LIMITING VALUE OF ABS(NORM(P-Q)).
READ 102,SK
102 FORMAT(F10,1)
CK=SK
PRINT 201,IIZ,IY,IIX,JIJ,JIY,JIJX,SK
50  201 FORMAT(/,27H FOR COVERING THE POINT A=(,3I10,3H)/(,3I10,12H) WITH

```



```

1800ND=F5.1)
ANV=DCI(AX,AY,AZ,R)
ANW=DCI(AX,AY,AZ,U)
REVD=DCI(BX,RY,RZ,R)
55 REW=DCI(BX,RY,RZ,U)
ARD=RFV*REV+REW*BEW
ABV=(ANV*REV+ANW*BEW)/ABD
ARW=(-ANV*BEW+BEV*ANW)/ABD
60 BNDORM=DNORM(BZ,BY,BX)
C...READ THE LOWER BOUND ON THE INTEGERS T WHICH ARE TO BE CONSIDERED
C...AND CALCULATE THE CORRESPONDING UPPER BOUND.
  READ 102,SL1
  TL1=SL1
  PRINT 205,SL1
65 205 FORMAT(/,4H TL=,F10.4)
  TU1=TL1/BP+ABP
  TL=TL1
  TU=TL+100.0D0
  TU=DMIN1(TU,TU1)
70 C...CALCULATE THE LIMITS ON V AND W WHERE V+I.W AND V-I.W ARE THE
C...CONJUGATES OF T.
  RMINIM=SK
C...IF T=P*H(3)+Q*H(2)+S*H(1) AND T HAS CONJUGATES V+I.W AND V-I.W, NOW
C...FIND THE COEFFICIENTS OF T,V,W IN THE FOLLOWING EXPRESSIONS FOR P,Q
75 C...AND S.
C... P=PT*T+PV*V+PW*W
C... Q=QT*T+QV*V+QW*W
C... S=ST*T+SV*V+SW*W
  RDET=FLOAT(IDET)
  BDET=DSORT(RDET)/2.0
  PT=DAJ(R,U,1,2)
  PV=DAJ(U,H,1,2)
  PW=DAJ(H,R,1,2)
  QT=DAJ(R,U,3,1)
  QV=DAJ(U,H,3,1)
  QW=DAJ(H,R,3,1)
  NEXT=T
80 C...NEXT=1 FOR THE FIRST BRANCH OF VALUES OF T AND 2 FOR THE SECOND.
  4 SR=DSQRT(CK/(TL-(AN/BE)))
  90 3 VL=ABV-SB
  VU=ABV+SB
  WL=ABW-SB
  WU=ABW+SB
C...NOW CALCULATE THE BOUNDS ON P AND Q.
  95 CALL LOWUP(PT,PV,PW,PL,PU)
  CALL LOWUP(QT,QV,QW,QL,QU)
C...NOW FIND WHAT THE INTEGER LIMITS ON P AND Q ARE AND FIND WHICH HAS
C...THE SMALLER RANGE.
  CALL LIMITL(PL,IPL)
100 CALL LIMITL(QL,IQL)

```

```

CALL LIMITU(PU,IPU)
CALL LIMITU(QU,IOU)
IF (IPL.LE.IPU) GO TO 70
71 PRINT 202,IPL,IPU,IQL,IOU
105 202 FORMAT(/,19H THE RANGE FOR P IS,15,3H TO,15,13H AND FOR Q IS,15,
13H TO,15)
STOP
70 IF (IQL.GT.IOU) GO TO 71
IF ((IPU-IPL).GT.(IOU-IQL)) GO TO 1
110 C...NOW P HAS THE SMALLER RANGE SO USING THE FACT THAT U(1)=0 AND SO
C... W=Q*U(2)+P*U(3)
C... Q=(W-P*U(3))/U(2).
C...THUS PROVIDED U(2)>0,WHICH IS SO, THE BOUNDS ON W AND EACH POSSIBLE
C...VALUE OF P Q MAY BE FURTHER RESTRICTED.
115 C...ALSO USING R(1)=1 AND SO
C... S=V-Q*R(2)-P*R(3).
C...THE BOUNDS ON V AND EACH POSSIBLE PAIR OF VALUES OF P AND Q THE
C...POSSIBLE VALUES OF S MAY BE FOUND. RUN THROUGH THE VALUES P FROM
C...IPL TO IPU.
120 IPCOUNT=IPU-IPL+1
P=FLOAT(IPL)-1.0
DO 12 IPC=1,IPCOUNT
P=P+1.000
C...FIND BOUNDS ON Q FOR THIS VALUE OF P
125 PQA=P*U(3)/U(2)
QLT=W/U(2)-PUA
QUT=WU/U(2)-PUA
C...AS SUITABLE VALUES OF Q ARE FOUND FOR EACH VALUE OF P: FOR EACH
C...PAIR P,Q FIND POSSIBLE S USING
130 C...V=S+Q*R(2)+P*R(3) .....SINCE R(1)=1
C... SO S=T-Q*R(2)-P*R(3)
CALL LIMITL(OLT,IOLT)
CALL LIMITU(QUT,IOUT)
IF (IOUT.LT.IOLT) GO TO 12
135 30 IOCCOUNT=IOUT-IOLT+1
O=FLOAT(IOLT)-1.0
PR=P*R(3)
DO 13 IQC=1,IOCCOUNT
O=O+1.000
140 PQ=-Q*R(2)-PR
SLT=VL+PQ
SUT=VU+PQ
CALL LIMITL(SLT,ISLT)
CALL LIMITU(SUT,ISUT)
145 IF (ISUT.LT.ISLT) GO TO 13
31 ISCOUNT=ISUT-ISLT+1
S=FLOAT(ISLT)-1.0
DO 14 ISC=1,ISCOUNT
S=S+1.000
150 IP=P & IO=Q & IS=S

```

```

CALL MULTD(P,Q,S,BZ,BY,BX,Z,Y,X)
X=X-AX E Y=Y-AY E Z=Z-AZ
APNORM=DNORM(Z,Y,X)
PRINT 203,IP,IQ,IS,APNORM,BNORM
155 203 FORMAT(/,8H WITH P=,I10,3H Q=,I10,7H AND S=,I10,/,27H THEN NORM(P-
1A) IS EQUAL TO,D24,1B,1H/,D24,1B)
RMNM=DABS(APNORM/BNORM)
IF (RMNM,LT,RMINIM)RMINIM=RMNM
14 CONTINUE
160 13 CONTINUE
12 CONTINUE
GO TO 2
1 IQCOUNT=IQ-IQL+1
C...PEACHING HERE INDICATES THAT Q HAS THE SMALLER RANGE SO USING THE
165 C...EXPRESSION FOR W P MAY BE RESTRICTED ASSUMING U(3)>0, WHICH IS SO.
C...THEN USING THE EXPRESSION FOR V, S MAY BE SIMILARLY RESTRICTED.
Q=FLOAT(IQL)-1.0
DO 15 IQC=1,IQCOUNT
Q=Q+1.0D0
170 C...FIND ROUNDS ON P FOR THIS Q.
QUA=Q*U(2)/U(3)
PLT=WL/U(3)-QUA
PUT=WU/U(3)-QUA
C...NOW FIND SUITABLE S.
175 CALL LIMITL(PLT,IPLT)
CALL LIMITU(PUT,IPUT)
IF (IPUT,LT,IPLT) GO TO 15
32 IPCOUNT=IPUT-IPLT+1
P=FLOAT(IPUT)-1.0
OR=Q*R(2)
180 DO 16 IPC=1,IPCOUNT
P=P+1.0D0
PO=-QR-P*R(3)
SLT=VL+PO
185 SUT=VU+PO
CALL LIMITL(SLT,ISLT)
CALL LIMITU(SUT,ISUT)
IF (ISUT,LT,ISLT) GO TO 16
33 ISCOUNT=ISUT-ISLT+1
190 S=FLOAT(ISLT)-1.0
DO 17 ISC=1,ISCOUNT
S=S+1.0D0
IS=S E IQ=Q E IP=P
CALL MULTD(P,Q,S,BZ,BY,BX,Z,Y,X)
195 X=X-AX E Y=Y-AY E Z=Z-AZ
APNORM=DNORM(Z,Y,X)
PRINT 203,IP,IQ,IS,APNORM,BNORM
RMNM=DABS(APNORM/BNORM)
IF (RMNM,LT,RMINIM)RMINIM=RMNM
200 17 CONTINUE

```

```

16 CONTINUE
15 CONTINUE
C...THE NEXT SECTION OF THIS BRANCH OF VALUES OF T IS NOW CONSIDERED.
205 2 IF(NEXT.EQ.2) GO TO 4
    IF(TU.GE.TU1) GO TO 5
C...WHEN TU REACHES THE VALUE TU1 THEN THIS BRANCH OF VALUES IS
C...COMPLETED SO THE OTHER IS NOW CONSIDERED.
    TL=TU
    TU=TU+100.000
210 TU=DMIN1(TU,TU1)
    GO TO 6
    5 NEXT=NEXT+1
    READ 102,SU1
    TU1=SU1
215 PRINT 206,SU1
    206 FORMAT(/,4H TU=,F10.1)
C...NOW READ THE UPPER BOUND ON THE LOWER BRANCH OF INTEGERS AND
C...CALCULATE THE LOWER BOUND
220 TL1=TU1/BB+ABP
    TU=TU1
    TL=TU-100.000
    TL=DMAX1(TL,TL1)
    9 SB=DSORT(CK/((ANZBE)-TU))
225 C...NOW REPEAT THE ABOVE WITH THESE LIMITS ON T AND THIS VALUE OF SB
    GO TO 3
    4 IF(TL.LE.TL1) GO TO 7
C...WHEN TL REACHES THE VALUE TL1 THEN THIS BRANCH OF VALUES IS
C...COMPLETED SO THE MINIMUM OF THIS POINT, RMINIM, IS PRINTED.
230 TU=TL
    TL=TL-100.000
    TL=DMAX1(TL,TL1)
    GO TO 8
    7 PRINT 204,RMINIM
235 204 FORMAT(/,30H THE MINIMUM FOR THIS POINT IS,F20.10)
    STOP
    END

```

```

SUBROUTINE BASES
C
DOUBLE A,B,C,R,U,H,V1,V2,V3,V4,V5,V6,V7,T,S,D,XIN
DOUBLE IORD
5 COMMON/D1/IDET,INDEX,IA,IB,IC/D2/R(3),U(3),H(3)
COMMON/D4/IL1,IL2,IL3,IT1,IT2,IT3,IP1,IP2,IP3
COMMON/D15/IURN/D16/IORD
A=FLOAT(IA) & B=FLOAT(IB) & C=FLOAT(IC) & IN=INDEX
C...FIND H(2) WHERE THE FIELD IS K(H(2)) AND ALSO R(2) AND U(2) WHERE
10 C...R(2)+T,U(2) AND R(2)-I,U(2) ARE THE CONJUGATES OF H(2).
V1=- (R-(A**2)/3.000)/3.000
V2=- (A*B/3.000-2.000*A**3/27.000-C)
V3=DSORT(V2**2-4.000*V1**3)
V4=(V2+V3)/2.000
15 V5=(V2-V3)/2.000
IF (V4.NE.0.000) GO TO 51
V6=0.000 & GO TO 52
51 D=1.000
IF (V4.GE.0.000) GO TO 53
V4=-V4 & D=-1.000
20 53 V6=DEXP((DLOG(V4))/3.000)*D
52 IF (V5.NE.0.000) GO TO 54
V7=0.000 & GO TO 55
54 D=1.000
IF (V5.GE.0.000) GO TO 56
V5=-V5 & D=-1.000
25 56 V7=DEXP((DLOG(V5))/3.000)*D
55 H(2)=V6+V7+A/3.000
R(2)=(A-H(2))/2.000
30 U(2)=DSQRT(C/H(2)-R(2)**2)
H(1)=1.000 & R(1)=1.000 & U(1)=0.000
C...(R(1),U(1),H(1))=(1,0,1) IS THE LATTICE POINT CORRESPONDING TO UNITY
C
35 C...FIND T AND S WHERE Y=(H(2)**2+T*H(2)+S)/INDEX.
C...(1,H(2),Y) IS A BASIS OF THE FIELD.
IN2=IN*IN & IN3=IN2*IN
K1=IA*IA-2*IB & K2=IA*IC & K3=IB*IC
K4=IB*IB-2*IA*IC & K5=IA*IB-3*IC & IC2=IC*IC
40 IF (INDEX.NE.1) GO TO 57
IT=0 & IS=0
T=0.0 & S=0.0 & GO TO 58
57 DO 59 ITC=1,IN
IT=ITC-1
DO 60 ISC=1,IN
IS=ISC-1
45 JA=K1+IT*IA+3*IS
IF (MOD(JA,IN).NE.0) GO TO 60
JB=K4+IT**2*IB+3*IS**2+IT*K5+2*IS*K1+2*IS*IT*IA
IF (MOD(JB,IN2).NE.0) GO TO 60
50 JC=IS**3+IT**3*IC+IC**2+IS**2*IT*IA+IS*K4+IS*IT**2*IB+IT*IB*IC

```

```

1+IT**2*IA*IC+IS**2*K1+IS*IT*K5
IF(MOD(JC,IN3).NE.0) GO TO 60
T=FLOAT(IT) & S=FLOAT(IS) & GO TO 58
55 60 CONTINUE
59 CONTINUE
58 XIN=FLOAT(INDEX)
H(3)=(H(2)**2+T*H(2)+S)/XIN
R(3)=(R(2)**2-U(2)**2+T*R(2)+S)/XIN
U(3)=(2.000*R(2)*U(2)+T*U(2))/XIN
60 IT1=INDEX
IT2=-IT
IT3=-IS
IML=2*IA*IT+IA*IA+2*IS+IT*IT-IB
IL1=IML/INDEX
65 IL2=(-IT*IML+2*IS*IT+IC-2*IB*IT-IA*IB)/IN2
IL3=(-IS*IML+2*IC*IT+IA*IC+IS*IS)/IN2
IP1=IA+IT
IP2=(IS-IB-IT*IP1)/INDEX
IP3=(IC-IS*IP1)/INDEX
70 PRINT 100,((R(I),U(I),H(I)),I=1,3)
100 FORMAT(/,27H A BASIS FOR THE LATTICE IS,/,15X,1HR,19X,1HU,19X,1HH,
13(/,5X,3D20.10))
I=2 & J=3
75 PRINT 101,I,I,IT1,IT2,IT3
PRINT 101,J,J,IL1,IL2,IL3
PRINT 101,I,J,IP1,IP2,IP3
101 FORMAT(/,10X,2HH(,I1,4H)*H(,I1,2H)=,I10,6HH(3) +,I10,6HH(2) +,I10,
14HH(1))
C...FIND AN UPPER BOUND, A POWER OF 10, FOR THE ARGUMENTS OF DNORM WHICH
80 C...WILL NOT GIVE RISE TO OVERFLOW.
K1P=IABS(K1) & K4P=IABS(K4) & K5P=IABS(K5)
NF=MAX0(K1P,K2,K3,K4P,K5P,IC2)
IONF=IOF(NF)
IONS=(29-IONF)/3
85 IURN=10**IONS
C...FIND BOUNDS FOR THE RELEVANT PRODUCTS OF MULTD SUCH THAT OVERFLOW
C...WILL NOT OCCUR.
IL1A=IABS(IL1) & IL2A=IABS(IL2) & IL3A=IABS(IL3)
IT1A=IABS(IT1) & IT2A=IABS(IT2) & IT3A=IABS(IT3)
90 IP1A=IABS(IP1) & IP2A=IABS(IP2) & IP3A=IABS(IP3)
ILTP=MAX0(IL1A,IL2A,IL3A,IT1A,IT2A,IT3A,IP1A,IP2A,IP3A)
ILTP1=IOF(ILTP)
IOBS=29-ILTP1
IOBD=10.000**IOBS
95 RETURN
END

```



```

SUBROUTINE MULTD(I1,J1,K1,I2,J2,K2,I,J,K)
C
C...MULTD FINDS THE PRODUCT OF I1*H(3)+J1*H(2)+K1*H(1) AND
C...I2*H(3)+J2*H(2)+K2*H(1). IT IS I*H(3)+J*H(2)+K*H(1)
5 C
      DOUBLE I1,J1,K1,I2,J2,K2,I,J,K,IP
      DOUBLE IOBD,IA,JA,KA,MULT1,MULT2,MULTP
      COMMON/D4/IL1,IL2,IL3,IT1,IT2,IT3,IP1,IP2,IP3
      COMMON/D16/IOBD
10      IA=DABS(I1) & JA=DABS(J1) & KA=DABS(K1)
      MULT1=DMAX1(IA,JA,KA)
      IA=DABS(I2) & JA=DABS(J2) & KA=DABS(K2)
      MULT2=DMAX1(IA,JA,KA)
      MULTP=MULT1*MULT2
15      IF (MULTP.LT.IOBD) GO TO 1
      PRINT 50
50      FORMAT(1X,10(1H*),27HDANGER OF OVERFLOW IN MULTD,10(1H*))
1      IP=I1*J2+I2*J1
      I=I1*I2*IL1+IP*IP1+J1*J2*IT1+I1*K2+I2*K1
20      J=I1*I2*IL2+IP*IP2+J1*J2*IT2+J1*K2+J2*K1
      K=I1*I2*IL3+IP*IP3+J1*J2*IT3+K1*K2
      RETURN
      END

```

R.H.C.  
LIBRARY

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SUBROUTINE LOWUP(XT,XV,XW,XL,XU)
C
C...THIS SUBROUTINE FINDS XL,XU SUCH THAT
C...      X=XT*Y(1)+XV*Y(2)+XW*Y(3)
5 C...IS BETWEEN XL AND XU, WHERE Y(I) IS BETWEEN YL(I) AND YU(I),I=1,3.
C
DOUBLE XT,XV,XW,XL,XU,X(3),YL,YU
COMMON/D3/YL(3),YU(3)
10 X(1)=XT & X(2)=XV & X(3)=XW
XL=XU=0.000
DO 1 I=1,3
IF (X(I).LT.0.000) GO TO 2
XL=XL+YL(I)*X(I)
15 XU=XU+YU(I)*X(I)
GO TO 1
2 XL=XL+YU(I)*X(I)
XU=XU+YL(I)*X(I)
1 CONTINUE
RETURN
20 END

```

```

      SUBROUTINE LIMITL(XXL,IXL)
C
C...IXL IS THE SMALLEST INTEGER GREATER THAN XXL.
C
5      DOUBLE XXL,XXLA
      XXLA=DABS(XXL)
      IF(XXLA.LT.1.0014) GO TO 1
      PRINT 50
10     50 FORMAT(1X,10(1H*),35HMORE THAN 14 DIGITS IN THIS INTEGER,10(1H*))
      1 XL=SNGL(XXL)
      IXL=IFIX(XL)
      IF(FLOAT(IXL)-XL)4,5,5
      4 IXL=IXL+1
      5 RETURN
15     END

```

```

SUBROUTINE LIMITU(XXU,IXU)
C
C...IXU IS THE GREATEST INTEGER LESS THAN XXU.
C
5   DOUBLE XXU,XXUA
   XXUA=DABS(XXU)
   IF(XXUA.LT.1.0D14) GO TO 1
   PRINT 50
10  50 FORMAT(1X,10(1H*),35HMORE THAN 14 DIGITS IN THIS INTEGER,10(1H*))
   1  XU=SINGL(XXU)
   IXU=IFIX(XU)
   IF(FLOAT(IXU)-XU)4,4,5
5   IXU=IXU-1
15  4 RETURN
   END

```

```
DOUBLE FUNCTION DAJ(R,U,I,J)
DOUBLE BDET,R(3),U(3)
COMMON/D5/BDET
DAJ=(R(I)*U(J)-R(J)*U(I))/BDET
RETURN
END
```

5